

FACTORY: $(444) 449-5879$
1422
$\therefore$ ante: the COMPLCCinch uses 270 S type EPROMS on its 24 K ROW -ards. If I desire to atomize my unit, I can do so by means + meriting properly = seamed 270,5's an the st card.
For that, Ind need a decent 80.80 development system. There is no doubt that the CCSciel has potential for a powerful system with proper development.

DESIRABLE ADDITIONS: (not necessarily all at once.)

1. D, C. S. \& EXTENDED BASIC WITH FILES.
2. RELOCATABLE MACRO -ASSEMBLER \& DISASSEMBLER.
3. variable precisicn/variable base ultan-calculatcr.
4. TEXT EDITOR/SELECTRIC I/O ROUTINES.
5. BASIC OR FORTRAN COMPILER. (VIGE INTERPETER.)
6. $270 \%$ PROGRAMMER /27I6 PROGRAMMER.
7. MORE POWERFUL PLOT SYSTEM INCLUDING 3-1) DIAGRAM CAPABILITIES, MAYBE HIDFEN-LINE/PLANES, AND SIRSTITUTION-BY-ELEMENT FROM TABLES.
 A FATE OF ? CPS...THIS COULD BE IMPROVED GREATLY BY INEREASNV
 these function are serviced b/ interrupts; cthermise, $N$. .)

# BASIC8001 



```
TABLE OF CONTENTS
```

PAGE
Introduction ..... l-2
Summary of Commands ..... 3-8
Error Messages ..... 9-11
BASIC 8001 Arithmetic ..... 12-17
BASIC 8001 Strings ..... 18-20
BASIC 8001 Immediate Mode ..... 21-23
BASIC 8001 Statements ..... 24-41
BASIC 8001 Arithmetic Functions ..... 42-50
BASIC 8001 User Define Functions ..... 51-53
BASIC 8001 String Functions ..... 54-55
BASIC 8001 Editing Commands ..... 56-60
Using Assembly Language Routineswith BASIC61

## BASIC 8001

INTRODUCTION

BASIC 8001 is a single-user, conversational programming language which uses simple English-type statements and familiar mathematical notations to perform an operation. BASIC 8001 is one of the simplest computer languages to learn and once learned has the facility of advanced techniques to perform more intricate manipulations or express a problem more efficiently.

BASIC 8001 is in incremental compiler which provides immediate translation and storage of user programs being input. This method decreases the response time of a RUN command and increases execution speed. BASIC 8001 has provision for alphanumeric character string, I/O and string variables, and allows user defined functions and assembly language subroutine calls from user BASIC 8001 programs.

BASIC 8001 can be run on any Intecolor 8001, Intecolor 8051 or Compucolor 8001 with a minimum of 8 K of user workspace.

LOADING AND RUNNING BASIC 8001
BASIC 8001 is provided in ROM and runs in ROM. BASIC 8001 is initiated by typing the ESC key, then the $W$ (BASIC) key. The dialogue described below is printed. This is a once-only dialogue and does not occur after an ESC key, and E key sequence. The READY message is printed after the ESC, E key sequence.

BASIC 8001 prints:
BASIC 8001 vl2/8/76 COPYRIGHT (C) 1976 BY CHARLES MUENCH
MAXIMUM RAM ADDRESS?
The user then types the highest RAM address that he has available or wants to use and then keys a carriage return.

MIDDLE OF FIRST RAM CAFD: $45055=65535-20480=$ AFFF HEX
One extra RAM card is 49151 $=65535-16384=$ BFFFhex.
Two extra RAM cards is $57343=65535-8192$ = DFFFhex.
Three extra RAM cards is $65529=65535-6=$ FFFq hex.

231 BYTES USED
FOR INITIALIZATION OF BASIC $8 Q Q 1$ STATUS. 247 BYTES USED AFTER DOING ANYTHING.

BASIC 8001 then prints the message,
READY
and waits for a command or program line to be typed.
If BASIC 8001 has been initialized as above but has returned to the CRT O.S. (by CPU Reset Key), then BASIC 8001 can be recalled without disturbing existing programs by typing the ESC key, then the E key. BASIC 8001 will then print the message READY.

If power fails, the CPU Reset key is hit or the unit is turned off, the unit returns to the CRT operating mode.

If the CPU Reset key or ESC delete keys are hit, the unit leaves BASIC 8001 and returns to the CRT operating mode. Any BASIC 8001 statement program is saved and can later be recalled if BASIC 8001 is re-entered by typing ESC, E.

BASIC 8001 has twenty-four (24) key word program statements, thirteen (13) editing and command statements, eighteen (18) mathematical functions, nine (9) string functions and eighteen (18) two-letter error messages. With these command and statement capabilities, BASIC 8001 is extremely simple to use and yet very versatile and powerful.

The next section provides an easy reference to BASIC 8001 capabilities. If the user is unfamiliar with BASIC 8001 language, then the remaining portion of this manual should be studied in sequence while having a terminal at your fingertips to run the example given. This manual should enable the user to become very proficient in BASIC 8001 when finished. Intelligent Systems Corporation and Compucolor Corporation have a number of BASIC 8001 programs on Floppy Tapes and are available at nominal prices. In addition, both companies will pay for BASIC 8001 programs that are provided on floppy tape when properly documented and accepted. Enjoy your self programming in BASIC 8001!

## 1. BASIC 8001 STATEMENTS

The following summary of BASIC statements defines the general format for the statement and gives a brief explanation of its use.

```
DATA value list
DEF function (argument) = expression
DIM variable (n), variable (n,m), variable \$(n), variable \$(n,m)
```

END

FOR variable=expressionl TO
expression2 STEP expression3
GOSUB line number

GOTO line number

THEN
IF expression GOTO line number

INPUT list

INPUT "string"; list
LET variable $=$ expression

NEXT variable

ON X GOSUB line number list

ON X GOTO line number list

Used in conjunction with READ to input data into an executing program.

Defines a user function to be used in the program.

Reserves space for lists and tables according to subscripts specified after variable name.

Placed at the physical end of the program to terminate program execution.

Sets up a loop to be executed the specified number of times.

Used to transfer control to the first line of a subroutine.

Used to unconditionally transfer control to other than the next sequential line in the program.

Used to conditionally transfer control to the specified line of the program.

Used to input data from the terminal keyboard, promps with "?".

Used to input data without promp character.
Used to assign a value to the specified variable(s).

Placed at the end of a FOR loop to return control to the $F O R$ statement.

Call the Xth line number subroutine after GOSUB.

Branch to the Xth line number after GOTO.

| OUT I, X | Causes the X BYTE to be output to port I. |
| :--- | :--- |
| PLOT expression | Sends the one BYTE result of the expression <br> to the 8001 CRT. The result must be |
| between 0 and 255 binary. |  |

## IN AMT PORTS:

$$
\begin{aligned}
& \{x>19210+1020 \div \\
& \text { 共251-LIGHT PEN TOUUH SWITCH }\langle 193>x>61 \text { TOuCH, NO DETECT } \\
& \text { \#252-X-GOOR NATE iNPUT }
\end{aligned}
$$

$$
\begin{aligned}
& \text { \#253-Y-CORRTNATE INPUT (INVERTED) }
\end{aligned}
$$

## 2. COMMANDS

The following key commands halt program execution, erase characters or delete lines.
Key
CTRL/J or Line Feed
CTRL/M or RETURN
CTRL/K or ERASE LINE
CTRL/L or ERASE PAGE
CTRL/Z or CURSOR LEFT

Explanation
Terminates program execution. BASIC 8001 prints READY.

Must be typed to end every line typed in or to indicate the end of an INPUT.

A colon is used to separate multiple statements per line.

Deletes the entire current line.

Erases the CRT screen, but does not change or disturb BASIC 8001 statements in any way.

Deletes the last character entered and echoes a cursor left.

The following commands list, load, save, erase and execute the program currently in core.

## Command

CLEAR
clbar x Required for String data!

LIST

LIST line number

LOAD I

LOAD ? I

RUN

RUN line number

SAVE I
SAVE I: LOAD? I (Very useful)

## Explanation

Sets the array and string buffers to nulls and zeroes.

Sets space for string variable to $X$ characters normally 50 characters.

Prints the user program currently in core on the list output device.

Prints the program from the line specified to the end.

Does a NEW and inputs the program on track \#I from the READER input device.

Does not do a NEW but inputs and compares the program on track \#I with what is existing in RAM Memory.

Executes the program in the buffer area.
Executes the program starting at line number specified.

Outputs the program in core to track \#I of the WRITE output device.

Erases the entire storage area.

CONT
Continues execution after CTRL/J is typed or after a STOP statement.

The following functions perform standard mathematical operations in BASIC 8001.

## Name

ABS (x)

ATN (x)

CALL (x) THIS JMP MUST BE POKED IN AT -24575 (LO BYTE) and - 24574 (Hi BYTE.)
$\cos (x)$
$\operatorname{EXP}(\mathrm{x})$
FRE (x) DOES NOT INCLUDE FRE $(x \$)$ BYTES! $\operatorname{INT}(\mathrm{x})$
$\operatorname{INP}(\mathrm{x})$

- LOG (x)

PEEK(x) Same as POKE locations.

POS (x)

RND (x) Repeats after 1995 numbers.
SGN (x)

SIN (x)
$\operatorname{SPC}(x)$ DISTRUCTIVE TAB(x)
SOR (x)

TAB (x)

TAN (x)

## Explanation

Returns the absolute value of $x$.
Returns the arctangent of x as an angle in radians in the range + or $-\mathrm{pi} / 2$.

Call the user machine language routine at location $0 A 000 \mathrm{HEX}$. A $\alpha \phi D=-24576=5 \mathrm{MP}$
$A Q D 1=-24575=L 0$
$A Q D 2=-24574=$ Returns the cosine of $x$ radians.

Returns the value of $e^{x}$ where $e=2.71828$. Returns number of free BYTEs not in use.

Returns the greatest integer less than or equal to x .

Returns a BYTE from input port $0 \leq x \leq 255$.
Returns the natural logarithm of $x$.
Returns a BYTE from memory address $0 \leq x \leq 32767$
or if $X$ is negative the memory address is 65536- x.

Returns a value 0 to 79 current cursor position.

Returns a random number between 0 and 1. Returns a value indicating the sign of $x$. Returns the sine of x radians.

Causes x spaces to be generated.
Returns the square root of x .
Causes the cursor to tab to column number $x$ when used in a print statement.

Returns the tangent of x radians.

Name

AGC ( x \$)

CHRS (x)

GRE ( $\mathrm{x} \$$ )

LEFT\$ (x\$,I)

LEN ( x \$)

MID\$ (x\$,I,J)

RIGHT\$ (x\$,I)

STRS (x)

VAL ( $x$ \$)

CLEAR X

## Explanation

Returns as a decimal number the seven-bit internal code for the first character of string (x\$).

Generates a one-chasacter string having the ASCII value of $x$.

Returns number of free string BYTES.
Returns left most I characters of string ( $\mathrm{x} \$$ ) .

Returns the number of characters in the string (x\$).

Returns J characters of string (x\$)
starting at position I.
Returns right most $I$ characters of string ( $\mathrm{x} \$$ ) .

Returns the string which represents the numeric value of $x$.

Returns the number represented by the string ( $\mathrm{x} \$$ ) .

Reserves $X$ bytes for string data. Default value is $5 \phi$ bytes. No single input can exceed 96 bytes.

## ERROR MESSAGES

After an error occurs, BASIC 8001 returns to command level and types READY. Variable values and the program text remain intact, but the program cannot be continued and all GOSUB and FOR context is lost.

When an error occurs in a direct statement, no line number is printed.
Format of error messages:
Direct Statement XX ERROR

Indirect Statement XX ERROR IN YYYYY

In both of the above examples, "XX" will be the error code. The "YYYYY" will be the line number where the error occurred for the indirect statement.

The following are the possible error codes and their meanings:

ERROR CODE

BS

DD

CF

MEANING

Bad Subscript. An attempt was made to reference a matrix element which is outside the dimension of the matrix. This error can occur if the wrong number of dimensions are used in a matrix reference; for instance, LET A $(1,1, l)=Z$ when $A$ has been dimensioned DIM A(2,2).

Double Dimension. After a matrix was dimensioned, another dimension statement for the same matrix was encountered. This error often occurs if a matrix has been given the default dimension 10 because a statement like $A(I)=3$ is encountered and then later in the program a DIM A(100) is found.

Call Function error. The parameter passed to a math or string function was out of range. $C F$ errors can occur due to:
a) a negative matrix subscript (LET $A(-1)=0)$
b) an unreasonably large matrix subscript (>32767)
c) LOG-negative or zero argument
d) SOR-negative argument
e) A B with A negative and B not an integer.
f) A CALL (X) before the address of the machine language subroutine has been patched in (see Pg.7)
g) calls to MID\$, LEFT\$, RIGHT\$, INP, OUT, WAIT, PEEK, POKE, TAB, SPC or ON...GOTO with an improper argument.

| ID | Illegal Direct. You cannot use an INPUT or DEF statement as a direct command. |
| :---: | :---: |
| NF | NEXT without FOR. The variable in a NEXT statement corresponds to no previously executed FOR statement. |
| OD | Out of Data. A READ statement was executed but all of the DATA statements in the program have already been read. The program tried to read too much data or insufficient data was included in the program. |
| OM | Out of Memory. Program too large, too many variables, too many FOR loops, too many GOSUB's, too complicated an expression or any combination of the above. |
| OV | Overflow. The result of a calculation was too large to be represented in BASIC's number format. If an underflow occurs, zero is given as the result and execution continues without any error message being printed. |
| SN | Syntax error. Missing parenthesis in an expression, illegal character in a line, incorrect punctuation, etc. |
| RG | RETURN without GOSUB. A RETURN statement was encountered without a previous GOSUB statement being executed. |
| US | Undefined Statement. An attempt was made to GOTO, GOSUB or THEN to a statement which does not exist. |
| 10 | Division by Zero. |
| CN | Continue error. Attempt to continue a program when none exists, an error occurred, or after a new line was typed into the program. |
| LS | Long String. Attempt was made by use of the concatenation operator to create a string more than 255 characters long. |
| OS | Out of String Space. Save your program on paper tape or cassette, reload BASIC and allocate more string space or use smaller strings or less string variables. ALLOCATE STRING SPACE WITH CLEAR $X$. SEEPg. 5 ! |
| ST | String Temporaries. A string expression was too complex. Break it into two or more shorter ones. |
| TM | Type Mismatch. The left hand side of an assignment statement was a numeric variable and the right hand side was a string, or vice versa; or, a function which expected a string argument was given a numeric one or vice versa. |

Undefined Function. Reference was made to a user defined function which had never been defined.

## I. NUMBERS

BASIC treats all numbers (real and integer) as decimal numbers--that is, it accepts any decimal number and assumes a decimal point after an integer. The advantage of treating all numbers as decimal numbers is that any number or symbol can be used in any mathematical expression without regard to its type. Numbers used must be in the approximate range $10^{-38} \leqslant \mathrm{~N}<10^{+38}$.

In addition to integer and real formats, a third format is recognized and accepted by BASIC 8001. This format is called exponential or E-type notation, and in this format, a number is expressed as a decimal number times some power of l0. The form is:

## xxEn

where E represents "times 10 to the power of"; thus the number is read "xx times 10 to the power of $n$ ". For example:

$$
23.4 \mathrm{E} 2=23.4 * 10^{2}=2340
$$

Data may be input in any one or all three of these forms. Results of computations are output as decimals if they are within the range .O1_n 999999; otherwise, they are output in E format. Numbers are store $\bar{d}$ up to 24 bits of significance. If a number with more than 24 bits is entered, it is truncated and stored as 24 bits. BASIC 8001 handles six significant digits in normal operation and prints 6 decimal digits as illustrated below:

| Value Typed In | Value Output by BASIC 8001 |
| :---: | :---: |
| .01 | .01 |
| .0099 | $9.90000 \mathrm{E}-03$ |
| 999999 | 999999 |
| 1000000 | $1.00000 \mathrm{E}+06$ |

BASIC automatically suppresses the printing of leading and trailing zeroes in integer and decimal numbers, and, as can be seen from the preceding examples, formats all exponential numbers in the form:
(sign) x.xxxxxE(+ or -)n
where $x$ represents the number carried to six decimal places, $E$ stands for "times 10 to the power of", and $n$ represents the exponential value. For example:

$$
\begin{aligned}
& -3.47021 \mathrm{E}+08 \text { is equal to }-347,021,000 \\
& 7.26000 \mathrm{E}-04 \text { is equal to } .00726
\end{aligned}
$$

Floating point format is used when storing and calculating most numbers.

## NOTE

Because core size limitations prohibit the storage of infinite binary numbers, some numbers cannot be expressed exactly. In BASIC 8001, accuracy is approximately 5-1/2 digits, and errors in the 6th digit can occur. For example, . 999998 as a result of some functions may be equal to l. Discrepancies of this type are magnified when such a number is used in mathematical operation.

## II. VARIABLES

A variable in BASIC 8001 is an algebraic symbol representing a number, and is formed by a single letter, a letter optionally followed by a single digit or by double letters. For example:
NOTE: Variables may be a string of characters many long
Acceptable Variables Unacceptable Variables
Long variables are
very useful for inherant
documentation in a
program.

ACCEPTABLE $=8001$

I
EXCLUDES ANYTHING
B3
AB WHICH RESEMBLES A BASIC COMMAND ONLY 2 LEFTMOST CHARACTERS ARE SIGNIFIGANT.
X

2C-a digit cannot begin a variable.
ll-numbers alone cannot form a variable.

DALE $=69$
Subscripted and string variables are described in later sections. The user may assign values to variables either by indicating the values in a LET statement, or by inputting the values as data; these operations are discussed in another chapter.

The value assigned to a variable does not change until the next time a statement is encountered that contains a new value for that variable. All variables are set equal to zero (O) when the RUN command is issued. It is only necessary to assign a value to a variable when an initial value other than zero is required. However, good programming practice would be to set variables equal to 0 wherever necessary. This ensures that later changes or additions will not misinterpret values.

## III. SUBSCRIPTED VARIABLES

In addition to the simple variables described in the preceding section, BASIC 8001 allows the use of subscripted variables. Subscripted variables provide additional computing capabilities for dealing with lists, tables, matrices, or any set of related variables. In BASIC 8001 variables are allowed from 1 to 31 subscripts.

The name of a subscripted variable is any acceptable BASIC 8001 variable name followed by one or more integer expressions in parentheses within the range 0-32767. For example, a list might be described as A(I) where I goes from 0 to 5 as shown below:

$$
A(0), A(1), A(2), A(3), A(4), A(5)
$$

This allows reference to each of the six elements in the list, and can be considered a one dimensional algebraic matrix as follows:

| $A(0)$ |
| :---: |
| $A(1)$ |
| $A(2)$ |
| $A(3)$ |
| $A(4)$ |
| $A(5)$ |

A two-dimensional matrix $B(I, J)$ can be defined in a similar manner:

$$
B(0,0), B(0,1), B(0,2), . \quad . \quad \text {. } B(O J), . \quad . \quad, B(I, J)
$$

and graphically illustrated as follows:

| $B(0,0)$ | $B(0,1)$ | $B(0,2)$ | $B(0,3)$ |
| :--- | :--- | :--- | :--- |
| $B(1,0)$ | $B(1,1)$ | $B(1,2)$ | $B(1,3)$ |
| $B(2,0)$ | $B(2,1)$ | $B(2,2)$ | $B(2,3)$ |
| $B(3,0)$ | $B(3,1)$ | $B(3,2)$ | $B(3,3)$ |


| $B(0, J)$ |
| :--- |
| $B(1, J)$ |
| $B(2, J)$ |
| $B(3, J)$ |

- 
- 

.

| $\mathrm{B}(\mathrm{I}, 0)$ | $\mathrm{B}(\mathrm{I}, 1) \quad \mathrm{B}(\mathrm{I}, 2) \quad \mathrm{B}(I, 3) \mathrm{B}(I, J)$ |
| :--- | :--- | :--- |

Subscripts used with subscripted variables throughout a program can be explicitly stated or be any legal expression. If the value of the expression is non-integer, the value is truncated so that the subscript is an integer.

It is possible to use the same variable name as both a subscripted and unsubscripted variable. Both A and A(I) are valid variables and can be used in the same program. The variable A has no relationship to any element of the matrix A(I). BASIC 8001 will accept the same variable name as both a singly and a doubly subscripted variable name in the same program.

Character strings may also be subscripted variable arrays, and may have the same variable name i.e., A\$(I).

A Dimension (DIM) statement is used with subscripted variables to define the maximum number of elements in a matrix. ("Matrix" is the subscripted variable.) The DIM statement is discussed in a later paragraph.

If a subscripted variable is used without appearing in a DIM statement, it is assumed to be dimensioned to length 10 in each dimension (that is, having eleven elements in each dimension, 0 through l0). However, all matrices should be correctly dimensioned in a program.

## IV. EXPRESSIONS

An expression is a group of symbols which can be evaluated by BASIC 8001. Expressions are composed of numbers, variables, functions, or a combination of the preceding separated by arithmetic or relational operators.

The following are examples of expressions acceptable to BASIC 8001:
Arithmetic Expressions $\quad \underline{\text { String Expressions }}$

```
4
A$+B$+"ABC"
A7*(B^2+1)
```

Not all kinds of expressions can be used in all statements, as is explained in the sections describing the individual statements.

## V. ARITHMETIC OPERATIONS

BASIC 8001 performs addition, subtraction, multiplication, division and exponentiation. Formulas to be evaluated are represented in a format similar to standard mathematical notation. The five operators used in writing most formulas are:

Symbol
Operator

| OR |  |
| :--- | :--- |
| AND |  |
| NOT | A $+B$ |
| + | A $-B$ |
| - | A * B |
| * | A $/ B$ |
| / | A $\wedge$ B |

Meaning
Logical and bitwise "OR"
Logical and kitwise "AND" Logical and bitwise "NOT"
Add B to A
Subtract B from A
Multiply A by B
Divide A by B
Exponentiation (Raise A to the Bth power)

Unary plus and minus are also allowed, e.g., the - in -A+B or the + in $+X-Y$. Unary plus is ignored. Unary minus is treated as a zero minus the variable, e.g., $-A+B$ would be handled as $0-A+B$.
VI. PRIORITY OF ARITHMETIC OPERATIONS

When more than one operation is to be performed in a single formula, as is most often the case, rules are observed as to the precedence of the operators.

In any given mathematical formula, BASIC 8001 performs the arithmetic operations in the following order of evaluation:

1. Parentheses receive top priority. Any expression within parentheses is evaluated before an unparenthesized expression.
2. In the absence of parentheses, the order of priority is:
3. Exponentiation (proceeds from left to right).
b. Unary minus.
c. Multiplication and Division (of equal priority).
d. Addition and Subtraction (of equal priority).
e. Logical operators in the order NOT, AND, then OR.
4. If either 1 or 2 above does not clearly designate the order of priority, then the evaluation of expressions proceeds from left to right.

The expression A』BAC is evaluated from left to right as follows:

1. $\mathrm{A} \wedge \mathrm{B}$
$=$ step 1
2. (result of step l) $\boldsymbol{\wedge} \mathbf{C}=$ answer

The expression $A / B^{*} C$ is also evaluated from left to right since multiplication and division are of equal priority:

1. $A / B \quad=$ step 1
2. (result of step l)* $\mathrm{C}=$ answer

The expression $A+B^{*} C \wedge D$ is evaluated as:

1. $\mathrm{C} \boldsymbol{\wedge} \mathrm{D}=$ step 1
2. (result of step l)*B $=$ step 2
3. (result of step 2) $+\mathrm{A}=$ answer

Parentheses may be nested, or enclosed by a second set (or more) of parentheses. In this case, the expression within the innermost parentheses is evaluated first, and then the next innermost, and so on, until all have been evaluated.

In the following example:

$$
A=7 *((B \wedge 2+4) / X)
$$

The order of priority is:

1. $\mathrm{Bn}_{\mathrm{n}} \mathrm{C}=$ step 1
2. (result of step 1$)+4=$ step 2
3. (result of step 2 )/X = step 3
4. (result of step 3 )*7 = A

Parentheses also prevent any confusion or doubt as to how the expression is evaluated. For example:

A*B^2/7+B/C*DA2
( ( $\mathrm{A} * \mathrm{~B} \boldsymbol{\wedge} 2) / 7+((\mathrm{B} / \mathrm{C}) * \mathrm{D} \boldsymbol{\wedge} 2)$
Both of these formulas are executed in the same way, but the second is easier to understand.

Spaces may be used in a similar manner. Since the BASIC 8001 interpreter ignores spaces (except when enclosed in quotation marks), the two statements:

```
1\varnothing LET B = D^2 + 1
```

1 しLETB=D^2+1
are identical, but spaces in the first statement provide ease in reading. When the statement is subsequently printed, extra spaces are ignored.
VII. RELATIONAL OPERATORS

Relational operators allow comparison of two values and are used to compare arithmetic expressions or strings in an IF. . . THEN statement. The relational operators are:

Mathematical BASIC 8001.
Symbol Symbol
$=$
$\ll$
$\leq \quad<=$ or $=<$
$\gg$
$\geq \quad>=$ or $=>$
$\neq \quad\langle \rangle$ or $><$
The symbols $=\langle, \quad\rangle,\langle$ are accepted by BASIC 8001 but are converted to $\langle=\rangle=$,, and $\rangle$ and are shown in that form in a listing.

# All string variables initialize with a $5 \phi$-byte potential length. coppers. I cannot change this in -any way, not by 

 masc aol strauss
## CLEAR X RESERVES $X$ BYTES FOR STRING DATA <br> I. STRINGS

The previous section described the manipulation of numerical information only; however, BASIC 8001 also processes information in the form of character strings. A string, in this context, is a sequence of characters treated as a unit. A string can be composed of alphabetic, numeric, or alphanumeric characters. (An alphanumeric string is one which contains letters, numbers, spaces or any combination of characters.) A character string can be 255 characters long. Strings cannot be typed on more than one terminal line since a carriage return terminates the command.

## II. STRING VARIABLES

Any variable name followed by a dollar sign (\$) character indicates a string variable. For example:
are simple string variables and can be used, for example, as follows:

```
LET A$="HELLO"
```

PRINT AS

Note that the string variable A\$ is separate and distinct from the variable A.

In BASIC 8001, all control characters above control code $F$ (or 6) are legal within Quotes (") except for the following:

> Control Code K or $l l$ or erase line
> Control Code L or 12 or erase page
> Control Code M or 13 or return
> Control Code Z or 26 or cursor left

## III. SUBSCRIPTED STRING VARIABLES

Any list of matrix variable name followed by the $\$$ character denotes the string form of that variable. For example:

$$
\begin{array}{ll}
\mathrm{V} \$(\mathrm{n}) & \mathrm{M} 2 \$(\mathrm{n}) \\
\mathrm{C} \$(\mathrm{~m}, \mathrm{n}) & \mathrm{G} 1 \$(\mathrm{~m}, \mathrm{n})
\end{array}
$$

where $m$ and $n$ indicate the position of the matrix element within the whole.

The same name can be used as a numeric variable and as a string variable in the same program with no restriction. A one- and a two-dimensional matrix can have the same name in the same program. For example:

| A | $A(n)$ | $A(m, n)$ |
| :--- | :--- | :--- |
| $A S$ | $A \$(m, n)$ | $A \$(m, n($ |

can all be used in the same program.
String lists and matrices are defined with the DIM statement as are numerical lists and matrices.
IV. STRING OPERATIONS

Concatenation
Concatenation puts one string after another without any intervening characters. It is specified by a plus sign (+) and works only with strings. The maximum length of a concantenated string is 255 charasters.

For example:

```
l\varnothing READ A$, B$, C$
2\emptyset DATA "ll", "33", "22"
3\emptyset LET D$ = A$+C$+B$
3 5 ~ P R I N T ~ D \$ ~
4\varnothing END
RUN
112233
```

V. RELATIONAL OPERATIONS

When applied to string operands, the relational operators indicate alphabetic sequence. The comparison is done on the basis of the ASCII value associated with each character in the strings being compared. For example:

## 55 IF A\$<B\$ THEN 100

When line 55 is executed, the first characters of each string (A\$ and $B \$$ ) are compared, then the second characters of each string and so on until the character in $A \$$ is less than the character in $B \$$. Then execution continues at line 100. Essentially, the strings are compared for alphabetic order. The next page contains a list of the relational operators and their string interpretations.

In any string comparison, trailing blanks are ignored (i.e., "ABC" is equivalent to "ABC ").

GRE $(x)=A V A I L A B L E ~ B Y T E S ~ O F ~ P R O G R A M ~ M E M O R Y . ~$
FRE $(x \$)=$ AVAILABLE BITES OF STRING MEMORY.
Total Available memory is Fri $(x)+$ GRE $(x \$) \ldots$ can adjust $\ddagger$ transfer by CLEAR $x$.

BASIC 8001

Relational Operators Used With String Variables

| Operator | Example | Meaning |
| :---: | :---: | :---: |
| $=$ | $A \$=B \$$ | The strings A\$ and B\$ are alphabetically equal. |
| $<$ | $A \$<\mathrm{B}$ \$ | The string A\$ alphabetically precedes B\$. |
| $>$ | $A \$>B \$$ | The string A\$ alphabetically follows B\$. |
| $<=$ or $=\leqslant$ | $A \$ く=B \$$ | The string A\$ is equivalent to or precedes $B \$$ in alphabetical sequence. |
| $\boldsymbol{>}=$ or $=>$ | $A \$>=B \$$ | The string $A \$$ is equivalent to or follows $B \$$ in alphabetical sequence. |
| $<>$ or $><$ | $A \$\rangle$ B | The strings A\$ and B\$ are not alphabetically equal. |

## I. USE OF IMMEDIATE MODE FOR STATEMENT EXECUTION

It is not necessary to write a complete program to use BASIC 8001. Most of the statements discussed in this manual can be included in a program for later execution or given on-line as commands, which are immediately executed by the 8080 CPU. This latter facility makes BASIC 8001 an extremely powerful calculator.

BASIC 8001 distinguishes between lines entered for later execution and those entered for immediate execution solely by the presence (or absence) of a line number. Statements which begin with line numbers are stored; statements without line numbers are executed immediately upon being entered to the system. Thus the line:
$1 \varnothing$ PRINT "THIS IS A COMPUCOLOR 8001"
produces no action at the console upon entry, while the statement:
PRINT "THIS IS A COMPUCOLOR 8001"
causes the immediate output:
THIS IS A COMPUCOLOR 8001

## II. PROGRAM DEBUGGING

Immediate mode operation is especially useful in two areas: program debugging and the performance of simple calculations in situations which do not occur with sufficient frequency or with sufficient complications to justify writing a program.

In order to facilitate debugging a program, STOP statements can be liberally placed throughout the program. Each STOP statement causes the program to halt, at which time the various data values can be examined and perhaps changed in immediate mode. The

GO TO xxxxx
command is used to continue program execution (where xxxxx is the number of the next program line to be executed). GOSUB and IF commands could also be used. The values assigned to variables when the RUN command was executed remain intact until a NEW, CLEAR or another RUN command is executed.

If the STOP occurs in the middle of a FOR loop, modifications cannot be made to the section of the program preceding the FOR.

When using immediate mode, nearly all the standard statements can be used to generate or print results.

If CTRL/J or linefeed is used to halt program execution, the GO TO XXXX or CONT command can be used to continue execution, since CTRL J or linefeed does print the number of the line where execution stopped. It is easy to know where to resume the program.
III. MULTIPLE STATEMENTS PER LINE

Multiple statements can be used on a single line in immediate mode. For example:

```
A=l:PRINT A
l
```

Program loops are allowed in immediate mode; thus a table of square roots can be produced as follows:

FOR I=1 TO $1 \varnothing$ : PRINT I, SOR (I):NEXT I
IV. RESTRICTIONS ON IMMEDIATE MODE

The INPUT statement cannot be used in immediate mode and such use results in the following error message:

ID ERROR
READY

Certain commands, while not illegal, make no logical sense when used in immediate mode. Commands in this category are DEF, DIM and DATA.

Also since user functions are not defined until the program is executed, function references in immediate mode cause an error unless the program containing the definition was previously executed.

Thus, the following dialogue might result if a function was defined in a user program and then referenced in immediate mode.

```
1\varnothing DEF FNA(X) = X^2 + 2*X:REM SAVED STATEMENT
```

PRINT FNA(l):REM IMMEDIATE MODE

UF ERROR READY

```
but if the sequence of statements is:
    10 DEF FNA(X) = X^2+2*X:REM SAVED STATEMENT
    RUN
    READY
    PRINT FNA(1)
    3
    READY
the immediate mode statement is executed.
```

A user program is composed of lines of statements containing instructions to BASIC 8001. Each line of the program begins with a line number that identifies that line as a statement and indicates the order of statement execution. Each statement starts with an English word specifying the type of operation to be performed. The statement lines are terminated with the RETURN key which is non-printing.

## I. STATEMENT NUMBERS

An integer number is placed at the beginning of each line in a BASIC 8001 program. BASIC 8001 executes the statements in a program in numerically consecutive order regardless of the order in which they were typed. Statement numbers must be within the range 0 to 65529. When first writing a program, it is advisable to number lines in increments of five or ten to allow insertion of forgotten or additional lines when debugging the program.

All BASIC 8001 statements and computations must be written on a single line; they cannot be continued onto a following line. However, more than one statement may be written on a single line when each statement after the first is preceded by a colon (:). For example:
$1 \varnothing$ INPUT $\mathrm{A}, \mathrm{B}, \mathrm{C}$
is a single statement line, whereas
$2 \emptyset$ LET $\mathrm{X}=11:$ PRINT $\mathrm{X}, \mathrm{Y}, \mathrm{Z}: \operatorname{IF} \mathrm{X}=\mathrm{A}$ THEN $1 \varnothing$
is a multiple statement line containing three statements: LET, PRINT, and IF. Most statements may be used anywhere in a multiple statement line; exceptions are noted in the discussion of each statement. Only the first statement on a line can (and must) have a line number.: It should be remembered that program control cannot be transferred to'a statement within a line, but only to the first statement of a line.

## II. REMARK STATEMENT

It is often desirable to insert notes and messages within a user program. Such data as the name and purpose of the program, how to use it, how certain parts of the program work, and expected results at various points are useful things to have present in the program for ready reference by anyone using that program.

The REMARK or REM statement is used to insert remarks or comments into a program without these comments affecting execution. Remarks do, however, use core area which may be needed by an exceptionally long program.

The REMARK statement must be preceded by a line number and may be used anywhere in a multiple statement line. The message itself can contain
any printing character on the keyboard. BASIC 8001 completely ignores anything on a line following the letters REM. (The line number of a REM statement can be used in a GOTO or GOSUB statement, see sections pertaining to destination of a jump in the program execution.) Typical REM statements are shown below:

## 10 REM- THIS PROGRAM COMPUTES THE

11 REM- ROOTS OF A QUADRATIC EQUATION
III. THE ASSIGNMENT STATEMENT - LET

The LET statement assigns a value to the specified variable(s). The general format of the LET statement is:

```
LET variable = expression
```

where variable is a numeric or string variable and expression is an arithmetic or string expression. All items in the statement must be either string or numeric; they cannot be mixed. The word LET is optional.

The LET statement does not indicate algebraic equality, but performs calculations within the expression (if any) and assigns the value to the variable.

The meaning of the equal (=) sign should be clarified. In algebraic notation, the formula $X=X+1$ is meaningless. However, in BASIC 8001 (and most computer languages), the equal sign designates replacement rather than equality. Thus, this formula is actually translated: "add one to the current value of $X$ and store the new result back in the same variable X". Whatever value has previously been assigned to $X$ will be combined with the value 1 . An expression such as $A=B+C$ instructs the computer to add the values of $B$ and $C$ and store the result in a third variable $A$. The variable $A$ is not being evaluated in terms of any previously assigned value, but only in terms of $B$ and $C$. Therefore, if $A$ has been assigned any value prior to its use in this statement, the old value is lost; it is instead replaced by the value $B+C$.
Example: $\quad x=y=z=98$ is not evaluated for $x=98, y=98, z=98 \ldots$ rather, it is
evaluated logically where we test if $y=z=98$ and assign result to $x$.
LET $X=2$
Assigns the value 2 to the variable $x$.
LET $X=X+1+Y$
Adds $l$ to the current value of $X$ then adds the value of $Y$ to the result and assigns that value to X .
IV. THE DIMENSION STATEMENT - DIM

The DIMension statement is used to define the maximum number of elements in a matrix. The DIM statement is of the form:

$$
\text { DIM variable (n), variable }(n, m), \quad \text { variable\$(n), variable\$(n,m) }
$$

where variables specified are indicated with their maximum subscript value(s).

For example:
$1 \varnothing \operatorname{DIM} X(5), Y(4,2), A(1 \varnothing, 1 \varnothing)$
12 DIM A4 (1 $\varnothing \varnothing$ ), A\$(25)

Only integer constants (such as 5 or 5070) can be used in DIM statements to define the size of a matrix. Variables cannot be used to specify the bounds of arrays. Any number of matrices can be defined in a single DIM statement as long as their representations are separated by commas.

The first element of every matrix is automatically assumed to have a ubscript of zero. Dimensioning $A(6,10)$ sets up room for a matrix with 7 rows and ll columns. This zero element is illustrated in the following program:

```
l\varnothing REM - MATRIX CHECK PROGRAM
2\emptyset DIM A (6,1\varnothing)
3\emptyset FOR I=\emptyset TO 6
4\varnothing LET A (I, })=
5\emptyset FOR J=\varnothing TO l\varnothing
6\emptyset LET A ( , J) = J
7\emptyset PRINT A(I,J);
8\emptyset NEXT J:PRINT:NEXT I
9\varnothing END
```

| RUN |  |  |  |  |  |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\varnothing$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | $1 \varnothing$ |
| 1 | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ |
| 2 | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ |
| 3 | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ |
| 4 | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ |
| 5 | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ |
| 6 | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ | $\varnothing$ |

READY

Notice that a variable has a value of zero until it is assigned another value.

Whenever an array is dimensioned $(n, m)$, the matrix is allocated $m+l, n+l$ elements. Core space can be conserved by using the Oth element of the matrix. For example, DIM $A(5,9)$ dimensions a 6 x 10 matrix which would then be referenced beginning with the $A(0,0)$ element.

The size and number of matrices which can be defined depend upon the amount of storage space available.

A DIM statement can be placed anywhere in a multiple statement line and can appear anywhere in the program. A matrix can only be dimensioned once. DIM statements need not appear prior to the first reference to an array, although DIM statements are generally among the first statements of a program to allow them to be easily found if any alterations are later required.

All arrays specified in DIM statements are allocated space when the RUN command is executed.

## V. PLOT STATEMENT

The PLOT Statement is used to output the 8 bit BYTE value of an expression to the CRT Screen. The general format of the PLOT Statement is:

10 PLOT expression
The expression can be any combination of variables which will evaluate to a positive value between 0 and 255.

The following example will plot a point on the CRT Screen at Location 80, 96 (X,Y) :
$10 \mathrm{X}=80: \mathrm{Y}=96$
20
PLOT 2
30
PLOT $\mathrm{X}:$ REMARK THE 8001 PLOT MODE CODE
40 PLOT 255 : REMARKS THE 8001 PLOT MODE ESCAPE CODE

As another example enter:

PLOT 65
PLOT 65: PLOT66: PLOT67: PLOT68: PLOT69: PLOT 70: PLOT71 ABCDEFG
A READY
READY
READY k BASIC DOES NOT (L.F.) (C.R.) BETWEEN PLOTS. $^{\text {PI }}$

It can be seen that (since 65 is the decimal ASCII value for A) PLOT 65 is the same as PRINT "A";
VI. PRINT STATEMENT

The PRINT statement is used to output data to the terminal. The general format of the PRINT statement is:

1) PRINT list

The list is optional and can contain expressions, text strings, or both.
When used without the list, the PRINT statement:

25 PRINT
DOES NOT ERASE, JUST SKIPS DOWN.
causes a blank line to be output on the 8001 CRT Screen (a carriage return/ line feed operation is performed).
2) PRINT Expression

PRINT statements can be used to perform calculations and print results. Any expression within the list is evaluated before a value is printed. For example:

```
1\varnothing LET A=1 : LET B=2: LET C=3+A
2\emptyset PRINT
3\emptyset PRINT A+B+C
RUN
7
READY
```

All numbers are printed with a preceding and following blank space.

The PRINT statement can be used anywhere in a multiple statement line.
For example:
$1 \varnothing A=1:$ PRINT $A: A=A+5:$ PRINT: PRINT A
prints the following on the terminal when executed:

1

6

READY

Notice that the terminal performs a carriage return/line feed at the end of each PRINT statement. Thus the first PRINT statement outputs a 1 and a carriage return/line feed; the second PRINT statement the blank line; and the third PRINT statement, a 6 and another carriage return/line feed.
3) PRINT Strings

The PRINT statement can be used to print a message or string of characters, either alone or together with the evaluation and printing of numeric values. Characters are indicated for printing by enclosing them in double quotation marks. For example:
$1 \varnothing$ PRINT "TIME'S UP"
$2 \emptyset$ PRINT "NEVERMORE"
RUN
TIME'S UP
NEVERMORE
READY

As another example, consider the following line:
$4 \varnothing$ PRINT "AVERAGE GRADE IS"; X
which prints the following (where X is equal to 83.4):
AVERAGE GRADE IS 83.4

When a character string is printed, only the characters between the quotes appear; no leading or trailing spaces are added. Leading and trailing spaces can be added within the quotation marks using the keyboard space bar; spaces appear in the printout exactly as they are typed within the quotation marks.

When a comma separates a text string from another PRINT list item, the item is printed at the beginning of the next available print zone. Semicolons separating text strings from other items are ignored. Thus, the previous example could be expressed as:
$4 \varnothing$ PRINT "AVERAGE GRADE IS" X
and the same printout would result. A comma or semicolon appearing as the last item of a PRINT list always suppresses the carriage return/line feed operation.

BASIC 8001 does an automatic carriage return/line feed if a string is printing past column 80.
4) Use of "," and ";"

BASIC 8001 considers the 8001 CRT Screen to be divided into ten zones of eight spaces each. When an item in a PRINT statement is followed by a comma, the next value to be printed appears in the next available print zone. For example:
$1 \varnothing$ LET $A=3:$ LET $B=2$
$2 \emptyset$ PRINT $A, B, A+B, A * B, A-B, B-A$
When the preceding lines are executed, the following is printed:
$\begin{array}{llllll}3 & 2 & 5 & 6 & 1 & -1\end{array}$
Notice each character is 8 spaces from the next character.
Two commas together in a PRINT statement cause a print zone to be skipped. For example:
$1 \emptyset$ LET $A=1 \neq$ LET $B=2$
$2 \emptyset$ PRINT $A, B, A+B$
RUN
12
3
READY
If the last item in a PRINT statement is followed by a comma, no carriage return/line feed is output, and the next value to be printed (by a later PRINT statement) appears in the next available print zone. For example:

```
l\varnothing A=1:B=2:C=3
2\emptyset PRINT A, :PRINT B: PRINT C
RUN
l 2
3
READY
```

If a tighter packing of printed values is desired, the semicolon character can be used in place of the comma. A semicolon causes no further spaces to be output other than the leading and trailing space automatically output with each number. A comma causes the print head to move at least one space to the next print zone or possibly perform a carriage return/line feed. The following example shows the effects of the semicolon and comma.

```
l\emptyset LET A=l\swarrow}\textrm{B}=2\not~\textrm{C}=
2\emptyset PRINT A;B;C;
3\emptyset PRINT A+l;B+l;C+l
4\emptyset PRINT A,B,C
RUN
1
READY
```

The following example demonstrates the use of the formatting characters , and ; with text strings:
$12 \emptyset$ PRINT "STUDENT NUMBER"X, "GRADE ="G;"AVE. ="A;
13ø PRINT "NO. IN CLASS ="N
could cause the following to be printed (assuming calculations were done prior to line 130):

STUDENT NUMBER ll9ø5ø GRADE $=87$ AVE. $=85.44 \mathrm{NO} . \operatorname{IN}$ CLASS $=26$
5) PRINT Statement - TAB Function

The TAB function is used in a PRINT statement to write spaces to the specified column on the output device. The columns on the output devices are numbered 1 to 80.

The form of the command is:
PRINT TAB (x)
where (x) is the column number in the range $0-255$. (If X exceeds 80 , however, every other consecutive line is tabbed until the number of spaces to be output is less than or equal to 80). If the column number specified is greater than 255 or negative, an error message is printed as follows:

CF ERROR
READY

If (x) is non-integer, only the integer portion of the number is used.

If the column number ( x ) specified is less than or equal to the current column number, the TAB function has no effect.

## VII. INPUT STATEMENT

The INPUT statement is used when data is to be input from the terminal keyboard during program execution. The form of the statement is:

1) INPUT list
where list is a list of variable names separated by commas.
For example:
$1 \varnothing$ INPUT A,B,C
causes the computer to pause during execution, print a question mark, and wait for input of three numeric values separated by commas. The values are input to the computer by typing the RETURN key.

If too few values are entered, BASIC 8001 prints another ? to indicate that more data is needed. If too many values are typed, the excess data on that line is ignored and the message below is printed but program still continues. The values entered in response to the INPUT statement cannot be continued on another line and are terminated by the RETURN key. Values must be separated by commas, if more than one value is input on the same line.

When there are several values to be entered via the INPUT statement, it is helpful to print a message explaining the data needed. For example:

```
l\varnothing PRINT "YOUR AGE IS";
2\varnothing INPUT A
2) INPUT "string"; list
```

The INPUT statement can also contain quoted strings. The above example could be written:

10 INPUT "YOUR AGE IS?";A
Note that when a quoted string is included in a INPUT statement, the normal ? is not printed as a prompt character, and if desired, must be included as shown within the quotes above.

This feature allows BASIC 8001 to be programmed to handle fill-in-theforms type of applications.

## VIII. DATA STATEMENT

The DATA statement is used in conjunction with the READ statement to enter data into an executing program. One statement is never used without the other. The form of the statement is:

DATA value list
where the value list contains the numbers or strings to be assigned to the variables listed in a READ statement. Individual items in the value list are separated by commas; strings must be enclosed in quotation marks.

For example:
$15 \emptyset$ DATA 4,7.2,3,"ABC"
$17 \emptyset$ DATA $1,34 \mathrm{E}-3,3.17311$
The location of DATA statements is arbitrary as long as they appear in the correct order; however, it is good practice to collect all DATA statements near the end of the program.

When the RUN command is executed, BASIC 8001 searches for the first DATA statement and saves a pointer to its location. Each time a READ statement is encountered in the program, the next value in the data statement is assigned to the designated variable. If there are no more values in that DATA statement, BASIC 8001 looks for the next DATA statement.
IX. READ STATEMENT

A READ statement is used to assign the values listed in a DATA statement to the specified variables. The READ statement is of the form:

READ variable list

The items in the variable list may be simple variable names or string variable names and are separated by commas. For example:

```
l\varnothing READ A, B$, C(l)
2\emptyset DATA 12, "l2",.l2E2
```

Since data must be read before it can be used in a program, READ statements generally occur near the beginning of the program. A READ statement can be placed anywhere in a multiple statement line.

If there is no data available in the data table for the READ to store, the out of data message below is printed:

OD. ERROR IN xxxxx
READY

Items in the data list in excess of those needed by the program's READ statements are ignored.

## X. RESTORE STATEMENT

The RESTORE statement causes the program to reuse the data from the first DATA statement and is of the form:

RESTORE

For example:
$3 \emptyset$ RESTORE
causes the next READ statement following line 30 to begin reading data from the first DATA statement in the program, regardless of where the last value was found.

A further example of the use of RESTORE follows:

```
15 READ B,C,D
.
•
55 RESTORE
6\emptyset READ E,F,G
.
•
8\emptyset DATA 6,3,4,7,9,2
•
.
1\varnothing\varnothing END
```

The READ statements in lines 15 and 60 both read the first three data values provided in line 80. (If the RESTORE statement had not been inserted before line 60, then the second READ would pick up data in line 80 starting with the fourth value.)

Since the values are being read as though for the first time, the same variable names may be used the second time through the data, if desired. To skip unwanted values, replacement, or dummy, variables may be inserted. For example:

```
l REM - PROGRAM TO ILLUSTRATE USE OF RESTORE
2\emptyset READ N
25 PRINT "VALUES OF X ARE:"
3\emptyset FOR I=l TO N
4\varnothing READ X
5\emptyset PRINT X,
6\emptyset NEXT I
7\varnothing RESTORE
185 PRINT
19\emptyset PRINT "SECOND LIST OF X VALUES"
2\emptyset\emptyset PRINT "FOLLOWING RESTORE STATEMENT:"
21\varnothing FOR I=l TO N
22\emptyset READ X
23\emptyset PRINT X,
24\emptyset NEXT I
```

25\emptyset DATA 4,1,2
251 DATA 3,4
3\emptyset\emptyset END
RUN
VALUES OF X ARE:
l 2 3
SECOND LIST OF X VALUES
FOLLOWING RESTORE STATEMENT:
4 l 2 3
READY

```

The second time the data values are read, the first \(X\) picks up the value originally assigned to N in line 20 , and as a result, BASIC prints:
\(\begin{array}{llll}4 & 1 & 2 & 3\end{array}\)

To circumvent this, a dummy variable coula be inserted to pick up and store the first value. This variable would not be represented in the PRINT statement, so the output would be the same each time through the list.
XI. GOTO STATEMENT

The GOTO statement is used when it is desired to unconditionally transfer to some line other than the next sequential line in the program. In other words, a GOTO statement causes an immediate jump to a specified line, out of the normal consecutive line number order of execution. The general format of the statement is as follows:

GOTO line number

The line number to which the program jumps can be either greater or less than the current line number. It is thus possible to jump forward or backward within a program.

For example,
\(1 \varnothing\) LET A=2
\(2 \emptyset\) GOTO \(5 \varnothing\)
\(3 \varnothing\) LET A=SQR (A+l4)
\(5 \emptyset\) PRINT A,A*A
RUN
causes the following to be printed:
2

When the program encounters line 20, control transfers to line 50; line 50 is executed, control then continues to the line following line 50. Line 30 is never executed. Any number of lines can be skipped in either direction.

When written as part of a multiple statement line, GOTO should always be the last statement on the line, since any statement following the GOTO on the same line is never executed. For example:
\(11 \varnothing\) LET A=ATN (B2):PRINT A:GOTO 5
XII. IF-THEN, IF-GOTO STATEMENTS

The IF-THEN statement is used to transfer conditionally from the normal consecutive order of statement numbers, depending upon the truth of some mathematical relation or relations. The basic format of the IF statement is as follows:

THEN
IF expression rel.op. expression line number
GOTO
where expression is an arithmetic or string expression.
Expressions cannot be mixed; both must be string or both must be numeric. Numeric comparisons are handled as described in the ARITHMETIC Section. String comparisons are performed on the ASCII values of the strings as described in the STRING Section.
rel. op. is one of the operators described in the ARITHMETIC Section.
line number is the line of the program to which control is conditionally passed.

If the value of the expression is true, control passes to the line number specified.

If the value of the expression is false, control passes to the next statement in sequence.

Examples:
```

1\varnothing IF A=B THEN 2\emptyset:PRINT "A B"
15 STOP
2\emptyset PRINT A+B
1\varnothing IF A}<>1\varnothing\mathrm{ GOTO 2ø :PRINT A
15 STOP
2\emptyset D=A+B*C
1\varnothing IF A$<B$ THEN 2\emptyset:STOP
2\emptyset PRINT A\$

```
XIII. FOR-NEXT STATEMENTS

FOR and NEXT statements define the beginning and end of a loop. (A loop is a set of instructions which are repeated over and over again, each time
being modified in some way until a terminal condition is reached.) The FOR statement is of the form:
```

FOR variable = expressionl TO expression2 STEP expression3

```
where
\begin{tabular}{ll} 
variable & must be a nonsubscripted numeric variable. \\
expression & \begin{tabular}{l} 
is an arithmetic expression which may be non- \\
integer.
\end{tabular}
\end{tabular}

The variable is the index; expressionl is the initial value; expression2, the terminal value and expression3, the increment value.

For example:
```

15 FOR K=2 TO 2\emptyset STEP 2

```
causes the program execution of the designated loop as long as K is less than or equal to 20. Each time through the loop, \(K\) is incremented by 2 , so the loop is executed a total of 10 times. When \(K=20\), program control passes to the line following the associated NEXT statement.

The index variable must be unsubscripted, although a common use of such loops is to deal with subscripted variables using the control variable as the subscript of a previously defined variable. The expressions in the FOR statement can be any acceptable BASIC 8001 expression.

The NEXT statement signals the end of the loop which began with the FOR statement. The NEXT statement is of the form:

NEXT variable
where the variable is the same variable specified in the FOR statement. Together the FOR and NEXT statements define the boundaries of the program loop. When execution encounters the NEXT statement, the computer adds the STEP expression value to the variable and checks to see if the variable is still less than or equal to the terminal expression value. When the variable exceeds the terminal expression value, control falls through the loop to the statement following the NEXT statement. Note the variable is not necessary since when a NEXT statement is encountered it is assumed it is for the appropriate FOR loop variable.

If the STEP expression and the word STEP are omitted from the FOR statement, +1 is the assumed value. Since +1 is a common STEP value, that portion of the statement is frequently omitted.

The expressions within the FOR statement are evaluated once upon initial entry to the loop. The test for completion of the loop is made after each execution of the loop. (If the test fails initially, the loop is still executed once.)

The index variable can be modified within the loop. When control falls through the loop, the index variable retains the value used to fall through the loop.

The following is a demonstration of a simple FOR-NEXT loop. The loop is executed 10 times; the value of \(I\) is 11 when control leaves the loop; and +1 is the assumed STEP value:
```

l\varnothing FOR I=l TO l\varnothing
2\emptyset PRINT I
3\emptyset NEXT I
4\emptyset PRINT I

```

The loop itself is lines 10 through 30. The numbers 1 through 10 are printed when the loop is executed. After \(\mathrm{I}=10\), control passes to line 40 which causes 11 to be printed. If line 10 had been:
\[
1 \varnothing \text { FOR } I=1 \varnothing \text { TO } 1 \text { STEP }-1
\]
the value printed by line 40 would be \(\varnothing\).
```

l\emptyset FOR I = 2 TO 44 STEP 2
2\emptyset LET I = 44
3\emptyset NEXT I

```

The above loop is only executed once since the value of \(I=44\) has been reached and the termination condition is satisfied.

If the initial value of the variable is greater than the terminal value, the loop is still executed once. The loop set up by the statement:
\(1 \varnothing\) FOR I \(=2 \emptyset\) TO 2 STEP 2
will be executed only once although a statement like the following will initialize execution of a loop properly:
\(1 \varnothing\) FOR I=2 \(\quad\) TO 2 STEP -2
For positive STEP values the loop is executed until the control variable is greater than its final value. For negative STEP values, the loop continues until the control variable is less than its final value.

FOR loops can be nested but not overlapped. The depth of nesting depends upon the amount of user storage space available (in other words, upon the size of the user program and the amount of RAM available). Nesting is a programming technique in which one or more loops are completely within another loop. The field of one loop (the numbered lines from the FOR statement to the corresponding NEXT statement, inclusive) must not cross the field of another loop.

Two Level Nesting
\(\left[\begin{array}{l}\text { FOR Il }=1 \text { TO } 10 \\ {\left[\begin{array}{l}\text { FOR I2 }=1 \\ \text { NEXT I2 }\end{array}\right.} \\ {\left[\begin{array}{lll}\text { FOR I3 }=1 & \text { TO } & 10 \\ \text { NEXT I3 }\end{array}\right.} \\ \begin{array}{l}\text { NEXT Il }\end{array}\end{array}\right.\)
\(\left[\begin{array}{l}\text { FOR Il }=1 \text { TO } 10 \\ {\left[\begin{array}{l}\text { FOR I2 }=1 \\ \text { NEXT I1 } \\ \text { NEXT I2 }\end{array}\right.}\end{array}\right.\)

\section*{Three Level Nesting}



An example of nested FOR-NEXT loops is shown below:
```

5 DIM X (5,1\emptyset)
l\emptyset FOR A=1 TO 5
2\emptyset FOR B=2 TO l\emptyset STEP 2
3\emptyset LET X (A,B) = A+B
4\varnothing NEXT B
5\emptyset NEXT A
55 PRINT X(5,1\emptyset)

```

When the above statements are executed, BASIC 8001 prints 15 when line 55 is processed.

It is possible to exit from a FOR-NEXT loop without the control variable reaching the termination value. A conditional or unconditional transfer can be used to leave a loop. Control can only transfer into a loop which had been left earlier without being completed, ensuring that termination and STEP values are assigned.

Both FOR and NEXT statements can appear anywhere in a multiple statement line. For example:
lø FOR I=1 TO lø STEP 5:NEXT I: PRINT "I="; I
causes:
\[
\mathrm{I}=11
\]
to be printed when executed.

A subroutine is a section of code performing some operation required at more than one point in the program. Sometimes a complicated I/O operation for a volume of data, a mathematical evaluation which is too complex for a user-defined function, or any number of other processes may be best performed in a subroutine.

More than one subroutine can be used in a single program, in which case they can be placed one after another at the end of the program (in line number sequence). A useful practice is to assign distinctive line numbers to subroutines; for example, if the main program uses line numbers up to 199, use 200 and 300 as the first numbers of two subroutines.

Subroutines are usually placed physically at the end of a program before DATA statements, if any. The program begins execution and continues until it encounters a GOSUB statement of the form:
1) GOSUB line number
where the line number following the word GOSUB is that of the first line of the subroutine. Control then transfers to that line of the subroutine. For example:
\(5 \emptyset\) GOSUB 2øø
Control is transferred to line 200 in the user program. The first line in the subroutine can be a remark or any executable statement.

Having reached the line containing a GOSUB statement, control transfers to the line indicated after GOSUB; the subroutine is processed until BASIC 8001 encounters a RETURN statement of the form:

\section*{2) RETURN}
which causes control to return to the statement following the original GOSUB statement. A subroutine must always be exited via a RETURN statement.

Before transferring to the subroutine, BASIC 8001 internally records the next sequential statement to be processed after the GOSUB statement; the RETURN statement is a signal to transfer control to this statement. In this way, no matter how many subroutines there are or how many times they are called, BASIC 8001 always knows where to transfer control next. The following program demonstrates the use of GOSUB and RETURN.
```

l REM - THIS PROGRAM ILLUSTRATES GOSUB AND RETURN
1\varnothing DEF FNA(X)= ABS(INT(X))
2\emptyset INPUT A,B,C
3\emptyset GOSUB 1\varnothing\varnothing
4\varnothing LET A=FNA (A)

```
```

5\emptyset LET B=FNA(B)
6\varnothing LET C=FNA(C)
7\emptyset PRINT
8\emptyset GOSUB l\varnothing\emptyset
9\emptyset STOP
l\emptyset\emptyset REM - THIS SUBROUTINE PRINTS OUT THE SOLUTIONS
ll\emptyset REM - OF THE EQUATION: AX^2 + BX + C = \varnothing
l2\emptyset PRINT "THE EQUATION IS "A "*X^2 + " B"*X + "C
13\emptyset LET D=B*B - 4*A*C
14\emptyset IF D<>O THEN 17\emptyset
15\emptyset PRINT "ONLY ONE SOLUTION... X "; -B/(2*A)
16\emptyset RETURN
17\emptyset IF D<\emptyset THEN 2\emptyset\emptyset
18\varnothing PRINT "TWO SOLUTIONS...X =";
185 PRINT (-B+SQR(D))/(2*A); ") AND ("; (-B-SQR(D))/(2*A)
19\emptyset RETURN
2\emptyset\emptyset PRINT "IMAGINARY SOLUTIONS ...X=(";
2\emptyset5 PRINT -B/(2*A) "," SQR(-D)/(2*A) ") AND (";
2\emptyset7 PRINT -B/(2*A) ","; -SQR(-D)/(2*A) ")"
2l\varnothing RETURN
9\varnothing\varnothing END

```

Subroutines can be nested; that is, one subroutine can call another subroutine. If the execution of a subroutine encounters a RETURN statement, it returns control to the line following the GOSUB which called that subroutine. Therefore, a subroutine can call another subroutine, even itself. Subroutines can be entered at any point and can have more than one RETURN statement. It is possible to transfer to the beginning or any part of a subroutine; multiple entry points and RETURN's make a subroutine more versatile. Up to 20 levels of GOSUB nesting are allowed.
XV. END STATEMENT

The END statement is the last statement in a BASIC program and is of the form:

END

The line number of the END statement must be the largest line number in a given program, since any lines having line numbers greater than that of the END statement are not executed (although they are saved with the SAVE command).

The END statement is optional. When an END statement is executed, program execution stops and the READY message is printed.

\section*{XVI. STOP STATEMENT}

The STOP statement can occur several times throughout a single program with conditional jumps determining the actual end of the program. The STOP statement is of the form:

90 STOP
and causes:

BREAK IN 90
READY
to be printed when executed.
This signals that the execution of a program has been terminated and BASIC 8001 is able to accept further input.

\section*{BASIC 8001 FUNCTIONS}

\section*{ARITHMETIC FUNCTIONS}

BASIC 8001 provides functions to perform certain standard mathematical operations such as square roots, logarithms, etc.

These functions have three or four letter call names followed by a parenthesized argument. They are pre-defined and may be used anywhere in a program.
\begin{tabular}{|c|c|}
\hline Call Name & Function \\
\hline ABS (x) & Returns the absolute value of x . \\
\hline \(\operatorname{ATN}(\mathrm{x})\) & Returns the arctangent of x as an angle in radians in range + or \(-\mathrm{pi} / 2\). \\
\hline CALL ( x ) &  \\
\hline \(\cos (\mathrm{x})\) & Returns the cosine of x radians. \\
\hline EXP (x) & Returns the value of \(\mathrm{e}^{\mathrm{x}}\) where \(\mathrm{e}=2.71828\). \\
\hline FRE (x) & Returns number of free BYTES not in use. \\
\hline INT ( x ) & Returns the greatest integer less than or equal to \(x,(\operatorname{INT}(-.5)=-1)\). \\
\hline INP ( x ) & Returns a BYTE from input port \(0 \leq x \leq 255\). \\
\hline LOG (x) & Returns the natural logarithm of x . \\
\hline PEEK (x) & Returns a BYTE from memory address \(0 \leq x \leq 32767\) or if \(x\) is negative the memory address is \(65536+x\). \\
\hline POS (x) & Returns a value of current cursor positions between 0 and 79. \\
\hline RND (x) NOT only & Returns a random number between 0 and 1. \\
\hline SGN (x) & Returns a value indicating the sign of x . \\
\hline SIN (x) & Returns the sine of x radians. \\
\hline SPC ( x ) & Causes x spaces to be generated. \\
\hline SQR ( x ) & Returns the square root of x . \\
\hline TAB ( x ) & Causes the 8001 CRT to space over to column number \(x\). Valid in PRINT statement only. \\
\hline TAN ( x ) & Returns the tangent of x radians. 42 \\
\hline
\end{tabular}

The argument \(x\) to the functions can be a constant, a variable, an expression, or another function. A square bracket cannot be used as the enclosing character for the argument \(x\), e.g. SIN [x] is illegal.

Function calls, consisting of the function name followed by a parenthesized argument, can be used as expressions or as elements of expressions anywhere that expressions are legal.

Values produced by the functions \(\operatorname{SIN}(x), \operatorname{COS}(x), \operatorname{ATN}(x), \operatorname{SOR}(x), \operatorname{EXP}(\) and LOG(x) have six significant digits.
I. Sine and Cosine Functions, SIN(x) and COS(x)

The sine and cosine functions require an argument angle expressed in radian measure. If the angle is stated in degrees, conversion to radians may be done using the identity:
\[
\langle r a d i a n s\rangle=\langle d e g r e e s) *(p i / l 80)
\]

In the following example program, 3.14159 is used as a nominal value for pi. \(P\) is set equal to this value at line 20. At line 40 the above relationship is used (in the expression within the LET statemer to convert the input value into radians.
\begin{tabular}{|c|c|c|c|}
\hline \multicolumn{4}{|l|}{\(1 \emptyset\) REM - CONVERT ANGLE (X) TO RADIANS, AND} \\
\hline \multicolumn{4}{|l|}{11 REM - FIND SIN AND COS} \\
\hline \multicolumn{4}{|l|}{\(2 \emptyset\) LET P = 3.14159} \\
\hline \multicolumn{4}{|l|}{25 PRINT "DEGREES", "RADIANS", "SINE", "COSINE"} \\
\hline \multicolumn{4}{|l|}{\(3 \emptyset\) INPUT X} \\
\hline \multicolumn{4}{|l|}{\(4 \varnothing\) LET Y \(=\mathrm{X}^{*} \mathrm{P} / 180\)} \\
\hline \multicolumn{4}{|l|}{\(6 \emptyset\) PRINT X, Y, SIN(Y), COS (Y)} \\
\hline \multicolumn{4}{|l|}{\(7 \emptyset\) GOTO 30} \\
\hline \multicolumn{4}{|l|}{RUN} \\
\hline DEGREES & RADIANS & SINE & COSINE \\
\hline \multicolumn{4}{|l|}{? \(\varnothing\)} \\
\hline \(\varnothing\) & \(\varnothing\) & \(\emptyset\) & 1 \\
\hline \multicolumn{4}{|l|}{? \(1 \varnothing\)} \\
\hline \(1 \varnothing\) & . 174533 & . 173648 & . \(9848 \emptyset 8\) \\
\hline \multicolumn{4}{|l|}{? \(2 \varnothing\)} \\
\hline \(2 \varnothing\) & . \(349 \varnothing 66\) & . \(342 \not 22\) & . 939693 \\
\hline \multicolumn{4}{|l|}{? \(3 \varnothing\)} \\
\hline \(3 \varnothing\) & . 523598 & . 5 & . \(866 \emptyset 26\) \\
\hline \multicolumn{4}{|l|}{? \(36 \varnothing\)} \\
\hline \(36 \varnothing\) & 6.28318 & -5.2431øE-Ø6 & 1 \\
\hline \multicolumn{4}{|l|}{? 45} \\
\hline 45 & . 785398 & \(.7 \varnothing 71 \varnothing 6\) & \(.7 \emptyset 71 \varnothing 7\) \\
\hline \multicolumn{4}{|l|}{? \(9 \varnothing\)} \\
\hline \(9 \varnothing\) & 1.5798 & 1 & \(1.12352 \mathrm{E}-\varnothing 6\) \\
\hline \multicolumn{4}{|l|}{?RETURN} \\
\hline READY & & & \\
\hline
\end{tabular}

\section*{II. Arctangent Function, ATN(x); 'Tangent Function, TAN(x)}

The arctangent function returns a value in radian measure, in the range \(+\mathrm{pi} / 2\) to \(-\mathrm{pi} / 2\) corresponding to the value of a tangent supplied as the argument (X).

In the following program, input is an angle in degrees. Degrees are then converted to radians at line 40.

At line 70 the tangent value, \(Z\), is supplied as argument to the ATN function to derive the value found in column 4 of the printout under the label ATN (X). Also in line 70 the radian value of the arctangent function is converted back to degrees and printed in the fifth column of the printout as a check against the input value shown in the first column.

III. Square Root Function, SQR(x)

This function derives the square root of any positive value as shown below.
```

l\emptyset INPUT X
2\emptyset LET X = SQR(X)
3\emptyset PRINT X
4\varnothing GOTO l\emptyset
RUN
?16
4
?l\varnothing\emptyset
1\varnothing
?l\emptyset\varnothing\varnothing

```
```

        31.6228
        ?123456789
        lllll.l
        ?l7
        4.12311
        ?25E2
        5\varnothing
        ?197\varnothing
        44.3847
        ?(RETURN)
        READY
    IV. Exponential Function, EXP(x)
The exponential function raises the number e to the power x. EXP is
the inverse of the LOG function. The relationship is
LOG(EXP(X)) = X
The following program prints the exponential equivalent of an input
value. Note that the output values derived below are used as input to
the LOG function.
l\emptyset INPUT X
2\emptyset PRINT EXP(X)
4\emptyset GOTO l\emptyset
RUN
?4
54.5981
?1\varnothing
22ø26.5
?9.421\varnothing\varnothing6
l2345
?4.60517
1\emptyset\emptyset
?25
7.2\emptyset\emptyset49E+l\emptyset
?(RETURN)
READY
V. Logarithm Function, LOG(x)
The LOG function derives the logarithm to the base e of a given value. In the following program at line 20, the LOG function is used to convert an input value to its logarithmic equivalent.
$1 \varnothing$ INPUT X
$2 \emptyset$ PRINT LOG (X)
$3 \varnothing$ GOTO 1ø
RUN
?54.59815
4
? $22 \varnothing 26.47$
$1 \varnothing$

```
?12345
    9.421ø1
?1\varnothing\varnothing
    4.6Ø517
?.72\emptyset\emptyset49Ell
    25
?(RETURN)
READY
```

Logarithms to the base e may easily be converted to any other base using the following formula:

$$
\log _{a} N=\frac{\log _{e} N}{\log _{e} a}
$$

where a represents the desired base. The following program illustrates conversion to the base 10.

1 REM - CONVERT BASE E LOG TO BASE $1 \varnothing$ LOG.
5 PRINT "VALUE","BASE E LOG","BASE $1 \varnothing$ LOG"
15 INPUT X
17 PRINT X,
$2 \emptyset$ PRINT LOG(X),
$4 \varnothing$ PRINT LOG (X)/LOG (1 $\varnothing$ )
5甲 GOTO 15
$6 \varnothing$ END
RUN
VALUE BASE E LOG BASE $1 \varnothing$ LOG
?4
4 1.38629 .6ø2Ø6
? $25 \varnothing$
$25 \varnothing$ 5.52146 2.39794
? 5
5 1.6ø944 . 69897
?6甲
$6 \varnothing$
$4 . \varnothing 9434 \quad 1.77815$
? $1 \varnothing \varnothing$
$1 \varnothing \emptyset \quad 4.6 \emptyset 517 \quad 2$
? (RETURN)
READY

An attempt to do a LOG(0) or LOG of a negative number causes the CF error message.
VI. Absolute Function, ABS (x)

The ABS function returns an absolute value for any argument value. Absolute value is always positive. In the following program, various input values are converted to their absolute values and printed.

```
        l\varnothing INPUT X
        2\emptyset LET X = ABS (X)
        3\emptyset PRINT X
        4\emptyset GOTO l\emptyset
        RUN
    ?-35.7
        35.7
        ?2
        2
        ?25El\varnothing
        2.5\emptyset\emptyset\emptyset\emptysetE+ll
        ?1\varnothing5555567
        1. \emptyset5556E+\emptyset8
        ?1\varnothing.12345
        10.1234
        ?-44.555566668899
        44.5556
        ?(RETURN)
        READY
VII. Integer Function, INT(x)
    The integer function returns the value of the greatest integer not
    greater than x. For example:
        PRINT INT(34.67)
        34
        PRINT INT(-5.1)
        -6
        The INT of a negative number is a negative number with the same or
        larger absolute value, i.e., the same or smaller algebraic value.
        For example:
            PRINT INT(-23.45)
            -24
            PRINT INT(-14.39)
            -15
            PRINT INT(-ll)
            -ll
The INT function can be used to round numbers to the nearest integer, using INT(X+.5). For example:
PRINT INT (34.67+.5)
PRINT INT(-5.1+.5)
    -5
```

INT(x) can also be used to round to any given decimal place or integral power of 10 , by using the following expression as an argument:

```
(X*l0\uparrowD+.5)/l0 D
```

where $D$ is an integer supplied by the user.

```
1\varnothing REM - INT FUNCTION EXAMPLE
15 PRINT
2\emptyset PRINT "NUMBER TO BE ROUNDED:"
25 INPUT A
4\varnothing PRINT "NO. OF DECIMAL PLACES:"
4 5 ~ I N P U T ~ D ~
6\emptyset LET B = INT (A*lOAD + .5)/l\emptysetAD
7\emptyset PRINT "A ROUNDED = " B
8\emptyset GOTO 15
9\varnothing END
RUN
```

NUMBER TO BE ROUNDED:
? 55.65842
NO. OF DECIMAL PLACES:
? 2
A ROUNDED $=55.66$
NUMBER TO BE ROUNDED:
? 78.375
NO. OF DECIMAL PLACES:
?-2
A ROUNDED $=1 \varnothing \varnothing$
NUMBER TO BE ROUNDED:
?67. 38
NO. OF DECIMAL PLACES:
?-1
A ROUNDED $=7 \varnothing$
NUMBER TO BE ROUNDED:
? (RETURN)
READY
VIII. Random Number Function, RND (x) Repeats after 1995 numbers.

The random number function produces a random number, or random number set, between 0 and 1. The numbers are reproducible in the same order after ESC, E key if $x>0$ for later checking of a program. The argument (x) is not used and can be any number (it cannot be a string expression) ; it serves only to standardize all BASIC 8001 function representations. The form RND is not legal. For example:

```
1\varnothing REM - RANDOM NUMBER EXAMPLE.
25 PRINT "RANDOM NUMBERS:
3\emptyset FOR I = 1 TO 15
4\varnothing PRINT RND(l);
5\emptyset NEXT I
6\varnothing END
RUN
RANDOM NUMBERS:
.1\emptyset\emptyset25\emptyset.5\emptyset438 .964813.\varnothing267824 .886627.388\emptyset94 .636444.569123 .839\emptyset19.72\emptyset\emptyset21
.3\emptyset6121.2\emptyset9\emptyset46 . 285553.599886 .958221.744\emptyset55 .179351.46\emptyset434 .452117.433291
.985412.27376 .522186.7\emptyset1146 . 246246.59\emptyset584 . 7778\emptyset1.457448 .45\emptyset592.3\emptyset797
```

READY

To obtain random digits from 0 to 9 , change line 40 to read:
$4 \varnothing$ PRINT INT(10*RND(1)),
and run the program again. This time the results will be printed as follows:

RUN
RANDOM NUMBERS:

| 8 | 9 | 8 | 9 | 5 | 5 | 5 | 9 | 8 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 5 | 4 | 4 | 1 | 5 |  |  |  |  |  |

READY
It is possible to generate random numbers over a given range. If the open range $(A, B)$ is desired, use the expression:

$$
(B-A) * \operatorname{RND}(1)+A
$$

to produce a random number in the range $A<n<B$.
The following program produces a random number set in the open range 4,6 (the extremes, 4 and 6, are never reached).
$1 \varnothing$ REM - RANDOM NUMBER SET IN OPEN RANGE 4,6.
$2 \emptyset$ FOR B $=1$ TO 15
$3 \varnothing$ LET A $=(6-4) * \operatorname{RND}(1)+4$
$4 \varnothing$ PRINT A,
$5 \varnothing$ NEXT B
$6 \varnothing$ END
RUN

| $4.2 \emptyset \emptyset 54.59266$ | $5.929624 .2 \not 9985$ | $5.773255 .54 \varnothing 26$ | 5.272884 .76248 | $5.678 \emptyset 45.25946$ |
| :---: | :---: | :---: | :---: | :---: |
| $4.612245 .33 \emptyset 46$ | $4.5711 \varnothing 4.26695$ | 5.916445 .69965 | $4.3587 \emptyset 5.54721$ | $4.9 \emptyset 4235.65 \emptyset 21$ |
| $4.197 \emptyset 85 . \emptyset 9 \emptyset 34$ | $5 . \emptyset 44374.82533$ | $4.492495 .614 \emptyset 8$ | $5.5556 \emptyset 4.41632$ | $4.9 \emptyset 1185 . \emptyset 15 \emptyset 8$ |

READY

NOTE: Negative arguments, i.e., RND(-x) will start a new random number sequence. While RND ( $\varnothing$ ) will always generate the last random number. RND $(-x)$ will always restart the same random number sequence!
IX. Sign Function, SGN(x)

The sign function returns the value 1 if x is a positive value, $\varnothing$ if $x$ is 0 and -1 if $x$ is negative. For example:

PRINT SGN(3.42)
1
PRINT SGN(-42)
-l
PRINT SGN(23-23)
$\varnothing$
The following example program illustrates the use of the SGN function.
$1 \varnothing$ REM-SGN FUNCTION EXAMPLE.
$2 \varnothing$ READ A,B,C
25 PRINT "A = "A,"B = "B,"C = "C
$3 \varnothing$ PRINT "SGN(A) ="SGN(A), "SGN(B) ="SGN(B),
$4 \varnothing$ PRINT "SGN(C) ="SGN(C)
$5 \varnothing$ DATA $-7.32, .44, \varnothing$
$6 \varnothing$ END
RUN
$\mathrm{A}=-7.32 \quad \mathrm{~B}=.44 \quad \mathrm{C}=\varnothing$
$\operatorname{SGN}(\mathrm{A})=-1 \quad \mathrm{SGN}(\mathrm{B})=1 \quad \mathrm{SGN}(\mathrm{C})=\varnothing$
READY
X. Call Statement

The CALL statement can be inserted anywhere in the BASIC 8001 program and has the form:

CALL (expression)
Where
expression
is the argument to the assembly language routine. The argument may be an expression. This may include values passed to the user routine.

The CALL statement causes a jump to location A000 HEX, which, unless modified by the user, contains a jump to the CF ERROR routine. The user must modify these three locations to go to his routines.

USER DEFINED FUNCTIONS

In some programs it may be necessary to execute the same sequence of statements or mathematical formulas in several different places. BASIC 8001 allows definition of unique operations or expressions and the calling of these functions in the same way as the square root or trig functions.

These user-defined functions consist of a function name: the first two letters of which are FN followed by a third or a fourth letter. For example:

| legal |  |
| :--- | :--- |
|  |  |
| FNA | FNAS |
| FNAA | FN2 |
| FNAI |  |

Each function is defined once and the definition may appear anywhere in the program. The defining or DEF statement is formed as follows:

DEF FNa (argument) = expression (argument)
where a is a variable name. The argument may consist of a dummy variable and the number of arguments is limited to one variable. The expression may contain other program variables not among the argument variable. For example:

10 DEF FNA(S) $=$ SA2
causes a later statement:
$2 \emptyset \operatorname{LET} R=$ FNA (4) +1
to be evaluated as $R=17$. As another example:
$5 \varnothing \mathrm{DEF} \mathrm{FNB}(\mathrm{A})=\mathrm{A}+\mathrm{X} \boldsymbol{\mathrm { A }} 2$
$6 \emptyset Y=F N B(14)$
causes the function to be evaluated with the current value of the variable $X$ within the program.

The two following programs
Program \#l:

```
l\emptyset DEF FNS(A) = A^A
```

$2 \emptyset$ FOR I=1 TO 5
$3 \varnothing$ PRINT I, FNS (I)
$4 \varnothing$ NEXT I
$5 \varnothing$ END

Program \#2:

```
l\varnothing DEF FNS(X) = X^X
2\emptyset FOR I=l TO 5
3\emptyset PRINT I, FNS(I)
4\varnothing NEXT I
5\emptyset END
```

cause the same output:
RUN
$1 \quad 1$
$2 \quad 4$
$3 \quad 27$
$4 \quad 256$
$5 \quad 3125$

READY

The argument in the DEF statement can be seen to have no significance; it is strictly a dummy variable. (A DEF statement with no arguments is illegal.) The function itself can be defined in the DEF statement in terms of numbers, variables, other functions, or mathematical expressions. For example:

```
l\emptyset DEF FNA(X) = X^2+3* X+4
2\emptyset DEF FNB(X) = FNA(X)/2 + FNA(X)
3\emptyset DEF FNC (X) = SOR (X+4)+l
```

The statement in which the user-defined function appears can have that function combined with numbers, variables, other functions, or mathematical expressions. For example:
$4 \varnothing$ LET $R=F N A(X+Y+Z) * N /(Y \wedge 2+D)$
A user-defined function cannot have several arguments, as shown below:
$25 \operatorname{DEF} \operatorname{FNL}(\mathrm{X}, \mathrm{Y}, \mathrm{Z})=\operatorname{SQR}(\mathrm{X} \boldsymbol{\wedge} 2+\mathrm{Y} \boldsymbol{\wedge} 2+\mathrm{Z} \boldsymbol{\wedge} 2)$
will cause an error

SN ERROR IN 25.
READY

When calling a user-defined function, the parenthesized arguments can be any legal expressions. The value of each expression is substituted for the corresponding function variable. For example:
$1 \varnothing$ DEF $\operatorname{FNZ}(X)=X \wedge 2$
$2 \emptyset$ LET A=2
$3 \varnothing$ PRINT FNZ (2+A)
line 30 causes 16 to be printed.

If the same function name is defined more than once, then the last definition will be used. The program below
$1 \varnothing$ DEF FNX (X) $=\mathrm{X} \boldsymbol{\wedge} 2$
$2 \emptyset$ DEF FNX $(X)=X+X$
$3 \emptyset$ LET A=5
$4 \emptyset$ PRINT FNX(A)
will cause 10 to be printed.
The function variable need not appear in the function expression as shown below:

```
10 DEF FNA (X) = 4 +2
2O LET R = FNA(lO)+l
30 PRINT R
40 END
RUN
    7
```


## STRING FUNCTIONS

Like the intrinsic mathematical functions (e.g., SIN, LOG), BASIC 8001 contains various functions for use with character strings. These functions allow the program to concatenate two strings, access part of a string, determine the number of characters in a string, generate a character string corresponding to a given number or vice versa, search for a substring within a larger string, and perform other useful operations. The various functions available are summarized in the following table.

## String Functions

Function code Meaning

| ASC (x\$) | Returns the seven-bit internal code for the one-character string (x\$) as a decimal number. If the argument contains more than one character, then the first character in the string is returned. |
| :---: | :---: |
| CHRS ( x ) | Generates a one-character string having the ASCII value of x where x is a number greater than or equal to 0 and less than or equal to 255. For example: CHR\$(65) is equivalent to "A". Only one character can be generated. |
| FRE ( x \$ ) | Returns number of free string BY: , |
| LEFTS (xS, I) | Returns left most I characters of string (x\$) . |
| LEN ( x \$ ) | Returns the number of characters in the string $x \$$ (including trailing blanks). For example: $\begin{aligned} & \text { PRINT LEN(A\$) } \\ & 26 \end{aligned}$ |
| $\operatorname{MID}(\mathrm{x}$ ( $, I, J)$ | Returns the string of characters in position I through J in $x \$$. |
| RIGHT\$ (x\$, I) | Returns right most I characters of string (x\$). |
| STR\$ ( x ) | Returns the string which represents the numeric value of $x$ as it would be printed by a PRINT statement but without a leading or trailing blank. |

Returns the number represented by the string $x \$$. If $x \$$ does not represent a number, then $\varnothing$ value is returned.

In the above examples, $x \$$ and $y \$ r e p r e s e n t ~ a n y ~ l e g a l ~ s t r i n g ~ e x p r e s s i o n s, ~$ and $I$ and $J$ represent any legal arithmetic expressions.

User-Defined String Functions
Character string functions cannot be written in the same way as numeric functions.

BASIC 8001 provides several key commands which can be used to halt program execution, erase characters or delete lines. The below table provides an explanation of each of the key commands.

## Key Commands

| Key |
| :--- | :--- |
| CTRL/J |
| or LINEFEED |
| or $\downarrow$ |$\quad$| Interrupts execution of a command or program. |
| :--- |
| BASIC 800l prints the message |

If the RETURN key has already been typed, a program line can be corrected by typing the appropriate line number and retyping the line correctly.

The line can be deleted by typing the RETURN key immediately after the line number; removing both the line number and line from the program.

If the line number of a line not needing correction is accidentally typed, the cursor left key (CTRL Z) may be used to delete the number (s); then the correct number can be typed. Assume the line:
$1 \varnothing$ IF A>5 GO TO 23ø
is correct. A line 15 is to be inserted, but:
$1 \varnothing$ LET
is typed by mistake. The correction is made as follows:
$1 \varnothing$ LET \&\&\& \&5 LET $X=X-3$

Line 10 remains unchanged, and line 15 is entered.

Following an attempt to run a program, error messages may be output on the terminal indicating illegal characters or formats, or other user errors in the program. Most errors can be corrected by typing the line number(s) and the correction(s) and then rerunning the program. As many changes or corrections as desired may be made before runs.

The following editing commands are entered in immediate mode and terminated by the RETURN key. These commands are used to erase a program in RAM, and list, punch or run a program.

## I. NEW COMMAND

The NEW command clears current contents of the storage area set up by BASIC 800l. This deletes any commands, programs, arrays, strings or symbols currently stored by BASIC 8001.

NEW should be used before entering a new program from the terminal keyboard to be sure no old program lines will be mixed into the new program and to clear out the symbol table area.

Example:
NEW
READY
$1 \varnothing$ READ A
-
-
-
clears the storage area and inserts the program being input at the keyboard.

## II. LIST COMMAND

The LIST command prints the user program currently in core on the terminal.

A part of a program may be listed by typing LIST followed by a line number. This causes that line and all following lines in the program to be listed.

Type CTRL/J or linefeed key to halt the listing. BASIC 8001 returns to the READY message when the current line is finished.

The lines listed may differ slightly from those entered because:

1. Certain characters while acceptable to BASIC 8001 are stored in a standard manner.

| Character <br> Typed |
| :---: |
| $=<$ |
| $=>$ |


| Character Stored |
| :---: |
| < |
| $\rangle$ |

2. Literals are stored to 24 bits of accuracy. Those with more than 24 bits are truncated to 24 bits.
3. Although literal storage is 24 bits, output is truncated to 6 decimal digits.
4. Literals are output in standard BASIC 8001 format, regardless of how they were input; for example,
$1 \emptyset$ LET X=3. $\varnothing+1 . \emptyset \emptyset \emptyset \emptyset \emptyset \emptyset 1$
$2 \emptyset$ PRINT X-E7
LIST
$1 \emptyset$ LET X=3+1
$2 \emptyset$ PRINT X-1. $\varnothing \emptyset \emptyset \emptyset \emptyset E+\varnothing 7$
5. Spaces in the input program are ignored, except within strings and REM statements. The LIST command prints the program with a space inserted to separate the key word and the line number. The listed program is therefore easier to read.

Example:

## LIST $1 \varnothing \varnothing$

Lists line 100 and all remaining lines in the program.
III. .SAVE COMMAND

The SAVE command outputs the program in RAM to the specified device. The form of the command is:

SAVE A

The format of the program output by the SAVE command is exactly the same as that stored in RAM memory. It may be recalled by the same file name using the LOAD command.

## IV. RUN COMMAND

After the user program is entered into RAM, it can be executed by typing the command

RUN
and the RETURN key.
The program is scanned; arrays are created in core and then the program is executed. Any appropriate error messages are printed and when the END or STOP statement is encountered, execution halts and a message is printed.

After execution, the variables used in a program remain accessible for use in immediate mode until a NEW, CLEAR or another RUN command is executed.

## V. CLEAR COMMAND

The CLEAR command clears the contents of the user array and string buffers. This command is generally used when a program has been executed and then edited. Before it is rerun, the array and string buffers are set to zeros and nulls by the CLEAR command to provide more core.

These buffers will be filled again when the RUN command is executed.

Example:
$1 \varnothing A=1 \varnothing$
$2 \emptyset$ PRINT A
CLEAR

READY

RUN
$1 \varnothing$

READY

## VI. CLEAR X COMMAND

The CLEAR $X$ performs the same function as CLEAR without the argument, but the Argument $X$ reserves $X$ locations for string variables which are required in string calculations. Normally this is 50 locations unless changed by CLEAR X command.

## VII. CONTINUE COMMAND

Continues program execution after a Control J or line feed is typed or a STOP statement is executed. You cannot continue after any error, after modifying your program or before your program has been run.

One of the main purposes of CONT is debugging. Suppose at some point after running your program, nothing is printed. This may be because your program is performing some time-consuming calculation, but it may be because you have fallen into an "infinite loop". An infinite loop is a series of BASIC 8001 statements from which there is no escape. The BASIC 8001 will keep executing a series of statements over and over until you intervene or until power to the unit is cut off. If you suspect your program is in an infinite loop, type in a Control J or line feed. The line number of the statement BASIC 8001 was executing will be typed out.

After BASIC 8001 has typed out READY, you can use PRINT to type out some of the values of your variables. After examining these values, you may become satisfied that your program is functioning correctly. You should then type in CONT to continue executing your program where it left off, or type a direct GOTO statement to resume execution of the program at a different line.

You could also use assignment (LET) statements to set some of your variables to different values. Remember, if you line feed or Control J your program and expect to continue it later, you must not get any errors or type in any program lines. If you do, you won't be able to continue, and get a "CN" (continue not) error. It is impossible to continue a direct command. CONT always resumes execution at the next statement to be executed in your program when Control J or line feed was typed.

## VIII. LOAD I COMMAND

LOADS the program named I from the 8001 CPU operating system Reader Input port specified by the I/O BYTE at location 9 F 90 HEX , see the CPU O.S. Manual. A new command is automatically done before the LOAD I command is executed. When finished loading the READY command will appear as usual. If the unit can't find the file on the floppy tape, then an error message should appear.

## IX. LOAD?I COMMAND

Does same as LOAD I except that a NEW command is not performed and BASIC 8001 does a word-by-word comparison of file I with the program already existing in RAM memory. If they are the same, then READY appears, else

## VERIFY FAILURE READY

will appear.
This should always be used after saving a program with the SAVE I command to ensure that it was saved correctly and can be reloaded without error.

ROUTINES WITH BASIC


#### Abstract

BASIC 8001 has a facility which allows experienced 8080 assembly language programmers to interface their own assembly language routines to BASIC 8001. This facility permits the user to add functions to BASIC 8001 which can operate directly on special purpose peripheral devices. This section describes in some detail the internal characteristics of BASIC 8001 during the execution of a BASIC 8001 program, and is intended to serve as a programming guide for the creation of such user-coded assembly language functions. This material assumes the user is familiar with 8080 assembly language. For additional information on this subject, refer to an assembly language programming manual on the 8080 CPU .


The CALL statement is used to reference these assembly language routines from the BASIC 8001 program.

Example: To call Assembly Language program from BAsic and pass arguements.
Basic Program: (Multiply x by 2, by shifting to the left) [No speed advantage, by test.]
10 Input X
$20 \quad A=\operatorname{CALL}(X)$
30 PRINT XI
40 SOTO ID
5 PUKE - 24575, \%: POKE -24574, 176 Assembly lanculiae program:

CALL 25A2 H * Et X in DEREy.gster
MON $A_{9} E$

NOTE. COLLIN AT LOCATION $25 A$ LHEX W WII $^{2}$ get the value of $X$ into the $D, E$ register.

NoTe: This will return the contents of the $A, B$ register to the variable " $A$ " in line $\vec{Z} \phi$, and also return to BASIC $\mathbf{8 0 0 1}$.

REC
MON B, A
KRA $A$ *To return value must be in $A, B$ register. (connote: then why clear $A$ with XRA $A$ ?)
$\operatorname{JMP} 2 C 53_{\mathrm{H}}$ *Then jump to location $2 C .53_{\text {HEX }}$.
When Assembly Program is entered at a location, say $B \phi \phi \phi_{\text {HEX }}$, then must place a jump to $B \Phi O \dot{q}_{\text {Lex }}$ at location $A \Phi \phi \phi_{\text {HEx }}$ as BASIC SOO1 will jump to AOOO whenever it encounters a call to Assembly Language (This jump Mist be POKE $\mathbb{N}$ via BASLC)
ColeNote: these are informal "hints" from the COMPUCOLOR people, and they are not too accurate...take it for what it is worth. I corrected some obvious errors, but not all. Cannot yet verify their complete data structure. ( $D, E \notin A, B$ is incomplete) ColeNote: BASic 8QQ1 initializes A $\phi \phi \phi$ through $A \phi B 3 \ldots$ BEWARE? ColeNote: POKE $32687, x$ through POKE $\frac{61}{32} 767, X$ is stack area...gend scratch?

Additional comments:
My test program:
10 POKE - $24575, \phi$ : POKE - 24574,176
20 INPIAT $\times$
301 FOR $I=1$ TO 1000: $A=\operatorname{CALL}(x)$ : NEXT: PRINT X:A
H屮FGRI=1 TO IQ: $A=2^{*}(x): N E X T: P R I N T X ; A$
$5 \ddot{4}$ GOTO $2(1)$
99 END
Some programs

$$
\begin{array}{ll|l}
B Q O Q & \text { CALL } 25 A 2 \text { HEX } & \text { CD AR } 25 \\
B O Q 3 & \text { MON } B, E & 43 \\
B Q 04 & \text { KRA, } A & A F \\
B Q Q 5 & \text { IMP } 2 C 53_{\text {MEX }} & C 3532 C \\
\hline B Q Q Q & \text { CALL } 25 A 2 \text { MEX } & \\
B Q 43 & \text { MON } A, E
\end{array}
$$

## $\square$ <br> The Compucolor 8001 CRT

TABLEOFONTENTS
PAGE
PART 1
Specifications and RS232C Interface
Start-Up and Initialization
Summary of Control Codes 5-6
Summary of Escape Codes 7-8
Summary of Graphic Plot Submodes 9
CRT Refresh Memory 10
PART 11
Keyboard 11
Detail of Control Codes 12-18
Details of Escape Codes 19
Details of Graphic Plot Submodes 25-36
Light Pen Operation 37
APPENDIX A
Keyboard Layout A-1
Intecolor®8001 Code Set A-2
Input Flow Diagrams A-3
Input Command Delays A-4
CCI Code Assignments A-5
Jl and J2 Pin Assignment A-6
I/O Connector Layout A-7
APPENDIX B
Plot Mode Functions B-l
Plot Mode Characters and Codes B-2
X Point Plot and Y Point Plot B-3
XY Incremental Point Plot Movements B-4
$X$ and $Y$ Bar Graph Modes B-5
X Incremental Bar Graph, Y Incremental Bar Graph B-6
$\mathrm{X}_{\mathrm{O}} \mathrm{Y}_{\mathrm{O}}$ Vector Plot Mode B-7
APPENDIX C
TMS 5501
APPENDIX D
TMS 8080
APPENDIX E

## PROPRIETARY STATEMENT

This document, submitted in confidence, contains proprietary information which shall not be reproduced or transferred to other documents or disclosed to others or used for manufacturing or any other purpose without prior written permission of Intelligent Systems Corp.

## SPECIFICATIONS

Introduction
The Intecolor 8001 is an eight color intelligent CRT data terminal designed as a replacement for teletypes and black and white CRT data terminals. It is a self-contained, desk top unit which offers, with the use of a modem, two-way data communications over common voice telephone lines or teletype compatible current loops. It can also be used in the stand alone mode as a complete desk top computer if equipped with the proper options.

| Power: | 105-125 volts, 60 HZ , 250 watts Option ll: 205-250 volts, 50-60 HZ |
| :---: | :---: |
| Temperature: | $+10^{\circ} \mathrm{C}$ to $+40^{\circ} \mathrm{C}$ operating <br> $-30^{\circ} \mathrm{C}$ to $+70^{\circ} \mathrm{C}$ storage |
| Humidity: | 0 to 95\% non-condensing |
| Package Size | 17 1/2" high |
| Desk Mount | 19 3/8" wide |
| Version: | 22 l/2" long |
| Keyboard | 3 l/4" high |
| Dimensions: | 14 l/l6" wide x 5 l/2" deep |
| Weight: | 85 pounds |
| Screen | 19" diagonal measure |
| Size: | 186 sq. inch screen area $4 \times 3$ aspect ratio |
| Display | 120 sq. inches |
| Area: | (12.0" wide x l0.0" high) |
| Character | 80 characters per line, 25 lines per page |
| Format: | Option 16: 80 characters per line, 48 lines per page |
| Character | 64 ASCII Characters, $5 \times 7$ dot matrix |
| Style: | within a 6x8 dot pattern |
|  | Option 03: 32 Graphic Characters, 6x8 dot matrix |
|  | Option 17: 64 Graphic Characters, $6 \times 8$ dot matrix |

## Standard Interface

Standard I/O Ports

The standard Intecolor 8001 has two input ports.
One port, Jl, is an asynchronous serial RS 232C I/O, or if Option 07 is installed, a serial 20 ma current loop I/O. The other port, J2, accepts parallel input data from the keyboard and provides an 8 bit parallel output. The Intecolor 8001 is furnished with a crystal clock and provides a keyboard selectable baud rate of normal $110,150,300,1200,2400,4800$, and 9600 baud, or a high speed option of $880,1200,2400,9600,19,200$, 38,400, and 76,800 baud.

The serial input port is furnished without parity checking so that when in the Plot Mode, or CCI Mode, eight data bits can be received.

The signals for the standard RS 232C I/O ports are shown on page 3 and on J1 and J2 in Appendix A7.

Pin 2 of the Keyboard J2 connector signals the Data communications equipment that the terminal has received a byte and is processing the last byte received. The Unit's input port has a one byte buffer. So for maximum speed, the communications equipment can send the next byte as soon as it has detected the high to low transition on pin 2 . The wave form is shown below:


Direction
From ISC to DCE＊
From DCE＊to ISC
From ISC to DCE＊
From DCE＊to ISC
NA
From ISC to DCE＊
氐 甾
喁
U
®
唯
8
我
*DCE - Data Communication Equipment
Pin \#

| Pin \＃ | Signal Line <br> 1 |
| :--- | :--- |
| 2 | Protective Ground |
| 3 | Received Data |
| 4 | Request to Send |
| 5 | Clear to Send |

[^0]r
＊DCE－Data Communication Equipment

## START-UP AND INITIALIZATION

## Introduction

BEFORE ATTEMTPING TO OPERATE YOUR INTECOLOR ${ }^{\circledR} 8001$, IT IS
SUGGESTED THAT THIS SECTION BE READ AND UNDERSTOOD. The power switch (SWl) is located in the lower rear panel portion of the CRT case. Also located on this panel are the various input and output port connections. These are shown in Appendix A8. Connection diagrams are shown in Appendix A7.

## Power

Plug the line cord into a l20VAC-60HZ outlet (230VAC-50-60 HZ with Option ll). When the power switch is pushed up the terminal is in the operating state. After the switch is turned on, a 60 second warm up peripd is required before operating the terminal. The unit will come up in the initialized state, $S_{o}$.

Initialized State $-S_{o}$
The unit will always come up in the initialized state- $S_{O}$ when power is turned on after being off for at least 30 seconds.

In State $S_{o}$ the following conditions are true:
A. Visible foreground color = white
B. Visible background color = black
C. Reverse field flag = "O"
D. Visible A7 bit = "O" (unless otherwise noted)
E. Plot Bit = "O"
F. Page Mode Operation (unless otherwise noted)
G. Terminal Mode $=$ Local (unless otherwise noted)
H. Baud Rate $=9600$ with one stop bit (unless otherwise noted)
I. Write left to right with visible cursor
J. Blind foreground color $=$ red
K. Blind background color $=$ black
L. Blind A7 Bit $=$ "O"
M. Blind Plot Bit $=$ "O"
N. Blind Cursor at home or top left corner of screen.

After the above conditions have been set, the cursor is moved to the home position which is the top left hand corner of the screen, and the position of the first character of the first line. The screen will clear by an Erase Page command which effectively makes all 2000 ( 3840 with 80 character x 48 line option) characters; spaces ( 20 HEX) which are white, nonblinking ( 97 HEX ). The unit is now ready to accept commands from the keyboard or the serial input if connected.

Convergence and Purity
The units convergence and purity may need adjusting when initially received. Allow at least a 30 minute warm before setting the final convergence. See Appendix $C$ for convergence alignment.



|  | BIT CODE | LETTER | FUNCTION |
| :---: | :---: | :---: | :---: |
|  | 0 | @ | Visible cursor mode |
|  | 1 | A | Blind cursor mode |
| * | 2 | B | Plot via color pad |
|  | 3 | C | Transmit cursor $\mathrm{X}, \mathrm{Y}$ position |
|  | 4 | D | Not used |
| * | 5 | E | Re-entry to BASIC 8001 |
|  | 6 | F | Sets full duplex mode |
|  | 7 | G | Not used |
|  | 8 | H | Sets half duplex mode |
|  | 9 | I | Not used |
|  | 10 | J | Set write vertical mode |
|  | 11 | K | Sets roll up and write left to right mode |
|  | 12 | L | Sets local mode |
|  | 13 | M | Not used |
|  | 14 | N | Not used |
| * | 15 | 0 | Re-entry to the CPU operating system |
| * | 16 | P | Initializes and transfers control to the CPU operating system |
| * | 17 | Q | * Character insert mode |
|  | 18 | R | Baud rate selection mode <br> A7 on $=1$ stop bit, A7 off $=2$ stop bit |
| * | 19 | S | Transfer control to the 8080 assembler |
| * | 20 | T | Transfer control to the text editor |
| * | 21 | U | Insert one line |
| * | 22 | V | Delete one line |
| * | 23 | W | Initializes and transfers control to BASIC 8001 |


| LETTER | FUNCTION |
| :---: | :---: |
| X | Sets page mode and write left to right mode |
| Y | Test mode - fill page with next character |
| Z | Set write down on 45 degree mode |
| $\square$ | Not used |
| \} | Sets write up on 45 degree mode |
| ] | Set unit up for Block receive mode |
| $\wedge$ | Causes a jump to address 9FA $\mathrm{CH}^{\text {H }}$ |
|  | Transfer control to the CRT mode |

* Must include certain option to be operational


## SUMMARY OF GRAPHIC PLOT SUBMODES

FOR INTECOLOR 8001

PLOT NORMAL KEYSUBMODE

Plot Mode Escape
Charactor Plot

X Point Plot
Y Point Plot BOARD CODE

Control
OPTIONAL FUNCTION
KEYBOARD CODE

F 15
Control $>\quad$ F 14
Control $=\quad$ F 13
Control < F 12
X-Y Incremental Point Plot Control ; ll
$\mathrm{X}_{0}$ of X Bar Graph Control :
Y of X Bar Graph
Control 9 F 9

X max of X Bar Graph
Incremental X Bar Graph
Control 8
F 8

F 7
$Y_{0}$ of $Y$ Bar Graph
Control 6
F 6

X of $Y$ Bar Graph
Control 5
F 5

Y max of $Y$ Bar Graph
Control 4
F 4

Incremental Y Bar Graph
$X_{0}$ Vector Plot
$Y_{0}$ Vector Plot

Incremental Vector Plot
Control 3
F 3

F 2

F 1
$F \varnothing$

SUMMARY OF INCREMENTAL DIRECTION CODES FOR INTECOLOR 8001

|  |  | $\Delta \mathrm{Xl}$ |  | $\Delta \mathrm{Yl}$ |  | $\Delta \mathrm{X} 2$ |  | $\Delta Y 2$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :--- | :---: |
|  | If BIT $=1$ | $\mathrm{~A}_{7}$ | $\mathrm{~A}_{6}$ | $\mathrm{~A}_{5}$ | $\mathrm{~A}_{4}$ | $\mathrm{~A}_{3}$ | $\mathrm{~A}_{2}$ | $\mathrm{~A}_{1}$ |  |
| $\mathrm{~A}_{0}$ |  |  |  |  |  |  |  |  |  |
| Direction | + | - | + | - | + | - | + | - |  |
| Value | 80 | 40 | 20 | 10 | 8 | 4 | 2 | 1 |  |

The 2000 [3840] * characters for display are stored in a 4096 [8192] word RAM memory beginning at 32,768 ( 8000 HEX ) and ending at 36,767 ( 8 F 9 F HEX) $[40,447$ (9DFF HEX)]. The first word is the zero character stored as the $A_{7}$ bit and then the 7 bit ASCII code ( $A_{6}$ to $A_{0}$ ). The second word is the compositestatus for this character. It is composed of Plot Character Bit $\left(A_{7}\right)$, Foreground Blink ( $A_{6}$ ), Background color code ( $A_{5}, A_{4}, A_{3}$ ), and Foreground color code ( $A_{2}, A_{1}, A_{0}$ ).

Therefore, each screen character requires two 8 bit words in memory, (the screen character and the character's compositestatus). The RAM memory location 8 FAO HEX [ $9 \mathrm{FA} \varnothing$ ] to 8 FFF HEX [ 9 FFF HEX ] are used for scratch pad storages. Memory location $8 \mathrm{FBO}[9 \mathrm{FBO} \mathrm{HEX}]$ and $8 \mathrm{FBl}[9 \mathrm{FBl}]$ are the locations of the Cursor character position and line number respectively. With the Roll Mode (Option 15) memory location 8 FB2 [9FB2] provides the number of lines that the home position has been shifted or rolled.

* [Indicates value for 48 Line System]

REFRESH MEMORY WORD FOR ONE CHARACTER

EVEN $O D D=(E V E N+1)$

ASC II Code


Secrets of the printing characters:


SMALL LETTERS
repeats $\Phi \Phi$ to IF, SMALL LETTERS.

BIG LETTE RS... top half is printed on odd lines, bottom half on even lines.
Same ASCII character prints both halves, differing only by line location. repeats 80 to $9 F, B I G L E T T E R S$.

Can get limited extra characters by plotting some letter tops with different letter bottoms...cannot top tops with tops.

$$
\sim
$$

$$
4
$$



7
7



## Keyboard

The Intecolor 8001 has a detachable keyboard which presents the standard ASCII four level code. (See Appendix A-l for keyboard layouts). The keyboard keys are optically encoded by means of phototransistors, a light source and shutters attached to the keys. There are no switches to wear out and the unit is RFI free. The Keyboard does not provide two key rollover.

## CPU Reset

The CPU Reset key provides a reset signal to the 8080 CPU . Its primary function is to allow the operator to regain control of the terminal if the software the customer has installed gets hung in an endless loop. If the reset is operated properly the bell will issue a short beep upon the release of the key. If automatically forces the terminal to the $S_{o}$ state. That is, just as if the power had been turned off and then back on. If additional RAM memory is installed this memory area is not cleared, but the scratch RAM area within the CRT Refresh RAM card is cleared.

## Control Key

The control key must be held down while the proper alpha numeric key is depressed if a control function is desired. The control functions are either color coded or have its desired results engraved on top of the key. Those keys which have a name enclosed within a ( ) parentheses indicate that they are also standarized escape codes. The escape codes only require that the ESC key be depressed then the ( ) parentheses key desired.

Shift Key
The shift key must be held down while the proper alpha numeric key is depressed if a shifted function is desired. Note that both the control and shift key must be held down to generate certain codes from the keyboard using the alpha numeric keys. See Appendix A-2 for the keyboard code set.

```
DETAIL OF CONTROL CODES
```

All of the display commands can be entered either through the serial input port or the keyboard. The keyboard input port has the highest priority of all inputs or outputs. The eight bit Intecolor 8001 code set as shown in Appendix A-2 must be used for the serial input port. The display control commands are a subset of the 32 ASCII control code set, and a flow diagram of these commands is shown in Appendix A-3.

With some display commands, such as the Graphic Plot Mode, delays may be experienced at the higher baud rates. A chart for these delays is shown in Appendix A4.

The Intecolor 8001 display commands has been expanded by an additional 32 commands via the ESC, character sequence as shown in Appendix 5. The terminal employs two input pointer flags, one for the keyboard and one for the RS232C input. Each flag may point to a different Mode of operation and thus the terminal can act differently from the keyboard as compared to the RS232 input. (See blind cursor operation Code 1 on page 19 .)

Code $\varnothing$ Null - (Control @)
Has no effect upon the display
Code 1 Protect $\stackrel{\text { (Control A) }}{ }$
Not presently implemented so it has no effect upon the unit.

Code 2 Graphic Plot Mode - (Control B) - (Option 02)
The general Graphic Plot Mode is entered by a
binary code 2 or a Control Code B. (See Appendix B). It should be noted that the XY Plot Mode is also entered at the same time. If a plot mode other than XY Point Plot is desired, the next word that follows should then be a binary code from 240 to 255 . These codes represent the various plot submodes as shown in the summary of Graphic Plot Submodes.

An additional feature is available to allow a graphic plot to be erased by simply setting the Flag bit on before entering the plot mode. This causes an XOR function to exist when plotting. Therefore, if you plot the same point, bar or vector twice, the second time erases the original.

Once in the general Plot Mode, any of the plot submodes may be entered by sending the corresponding code to the terminal. When this code is received, a flag internal to the terminal, known as PLOFL, is set placing the terminal in the appropriate plot submode. It should be noted that in many of the plot submodes, PLOFL is automatically set to a different value upon completion of the operation of that submode causing the terminal to enter a new submode. This is done to make coding and operation of the terminal in the various plot functions easier for the operator. The various submodes and their interactions are explained in detail in Appendix B.

## Code 3 <br> Cursor $\mathrm{X}-\mathrm{Y}$ (Control C)

The visible cursor may be positioned any where on the screen simply by sending a 3-word sequence beginning with 03. The next two words that follow determine that $X$ character position (0-79) and $Y$ line position $(0-24)$ for 25 line unit or [0-47] for 48 line unit. Both $X$ and $Y$ values must be in binary form with the range indicated. The cursor home position (i.e., the top left hand corner) is position 0 , 0 while the bottom right hand corner is $(79,24)$ or $[79,47]$.

If the cursor is positioned at $X=80$ binary ( 50 HEX ) then the cursor will
disappear. But if a character is typed it will be positioned at the beginning of the line specified by $Y+1$, the cursor then reappears in character position l. Any cursor command will automatically force the cursor to reappear at the proper position in relation to character position O, line Y + l.

If the cursor $X$ values is 81 binary (5l HEX) or larger then the CRT ignores this as the visible cursor $X$ values and sends the unit into the blind cursor addressing mode. Once in the blind cursor $X-Y$ addressing mode three (3) additional words must be sent. They are blind cursor $X$ value, blind cursor $Y$ value, and the blind status word. The blind $X$ value must be in the range of $0-79$ and the blind $Y$ value must be in the range of $0-24$ or $[0-47]$. The blind status word must be in the same format as required in the CCI mode (control F). See the next page,

The blind A7 bit will be set on by sending from 128 binary to 255 binary instead of 81 binary when going from the visible cursor $X, Y$ mode to the blind cursor $X, Y$ mode. The Blind $A 7$ bit will be set off anytime a binary number between 81 and 127 is used to get into the blind $X, Y$ mode.

It should be noted that the $X$ and $Y$ cursor values received are masked to 0-127 and 0-31 [0-63] respectively. Then, if the value is still out of range, the $X$ value has 80 subtracted and the $Y$ values has 25 [48] subtracted.

When exiting from the blind cursor $X-Y$ mode the terminal is left in the blind cursor mode for what ever input device caused the mode to be entered. That is if after CPU reset is operated the keyboard causes the blind cursor $X Y$ to be addressed then the keyboard will be left in the blind cursor mode while the RS232 serial is still in the visible cursor mode.

## Code 4 EOT - (Control D)

Has no effect upon the display

Code 5 ENQ - (Control E)
Has no effect upon the display

Code 6 CCI - (Control F)
When this code is received the system accepts the next eight bit word from the serial input as the new compositestatus for the characters which follow. See CRT Refresh Memory Section.

The first three bits represents the Foreground Color with $\operatorname{Red}_{F}=A_{0}, \quad G_{r e e n}^{F}=A_{1}$, and Blue $F_{2}$. The next three bits represent the Background Color (optional) with $\operatorname{Red}_{B}=A_{3}$, Green $_{B}=A_{4}$, and Blue ${ }_{B}=A_{5}$. The next bit, $A_{6}$ is the Blink bit for the Foreground Color and the last bit, $A_{7}$ is Plot Character bit which causes the display to interpret the ASCII word as a 2 x 4 plot array.

Code 7 Bell - (Control G)
When this code is received a tone will sound for about 150 MS .

* Stays on until BASIC program reads an END, or INPUT..so a loop

Code 8 Home - (Control H)
When this code is received the cursor moves to 0,0 or the top left hand corner of the screen.

Code 9 Tab - (Control I)
When this code is received the cursor moves horizontally to the next tab position. The tab positions are fixed and are at every eight positions from zero.

Code 10 Line Feed - (Control J)
When this code is received the cursor moves down one line.
This is the only code used for cursor down.

Code ll Erase Line - (Control K)
When this code is received a carriage return is initiated and the characters from the beginning to the end of the line are replaced with spaces and have the same color and status as the present visible CCI status. The cursor is always positioned at the beginning of the line.

Code 12 Erase Page - (Control L)
When this code is received the complete screen is replaced with spaces that have the same color and composite status as the present visible CCI status. The cursor always returns to the Home position. The blind cursor is also positioned at home.

Code 13 Carriage Return - (Control M)
When this code is received, the cursor returns to the beginning of the line that it presently is on.

Code $14 \quad \underline{A}_{7} \underline{\text { On }} \quad$ - (Control N$)$
Upon receiving this code, the characters which are to be displayed have $A_{7}$ forced to a "l". This bit is used to allow 2 X character sizes for 48 line units. This effectively doubles the number of displayable character types from 128 to 256.

Code 15 Blink - $\underline{A}_{7}-\underline{O F F}-($ Control O)
When this code is received the characters which follow have $A_{7}$ set to "0" (i.e., opposite to $\underline{A}_{7} \underline{O n}$ as above) and also have the Blink bit, $A_{6}$ of the composite status for the character set to "O" (i.e., the opposite of Blink-On per Code 3l.)

Code 16 to 23 or Color Keys - There are eight color keys

|  |  |  | $\mathrm{A}_{2}$ | $\mathrm{~A}_{1}$ | $\mathrm{~A}_{0}$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  | 0 | 0 | 0 |
| Black | (Control P) | Code 16 | 0 | 0 | 1 |
| Red | (Control 2) | Code 17 | 0 | 1 | 0 |
| Green | (Control R) | Code 18 | 0 | 0 | 1 |
| Yellow | (Control S) | Code 19 | 1 |  |  |
| Blue | (Control T) | Code 20 | 1 | 0 | 0 |
| Magenta | (Control U) | Code 21 | 1 | 0 | 1 |
| Cyan | (Control V) | Code 22 | 1 | 1 | 0 |
| White | (Control W) | Code 23 | 1 | 1 | 1 |

When one of these eight codes is received then one of two things happens, depending upon the Flag bit. If the Flag is off then the key that is depressed will change the compositestatus to that Foreground Color code.

If the Flag is on, then the key that is depressed will change the compositestatus to the Background Color code. If Background Color option is used, then it will display that color. If Background Color option is not supplied, then no effect will be noticed.

Note that when the plot via color pad is selected, one of the eight color select keys will select one of the eight plot blocks. The plot option 2 is installed See Escape B section for details.

Code 24 Transmit - (Control X)

Whenever control $X$ is received the terminal starts transmission from the visible cursor present position to the end of the screen, or until it detects a FF, $\varnothing \varnothing$ Hex sequence in the Refresh memory.

The transmission sequence is terminated by a carriage return, either $\varnothing \mathrm{D}$ Hex or 8 D Hex at the customer option. It should be noted that there may be many $\varnothing \mathrm{D}$ Hex or 8 D Hex imbedded in the data transmission since these are legal words in the refresh memory.

The transmission sends each 8 bit word in memory in sequence. That sequence is the ASCII character, then the status of that character, followed by the next ASCII character and then its status until the $F F, \varnothing \varnothing$ sequence is detected.

The best way to have this data sent back to the terminal is via the ESC] or block receive mode.

Code 25 Cursor Right - (Control Y)
Moves the cursor right one character without destroying
any information.

Code 26 Cursor Left - (Control Z)
Moves the cursor left one character without destroying
any information.

Code 27 Escape - (Control $[$ )
The Escape command effectively expands the control code set by 32 additional code capabilities. This requires at least a two code sequence (ESC, letter) which then performs a given function. At present only 26 of the 32 additional command capabilities have been enabled. These commands are given in the following table. (For Detail see the Escape Code Section).

SEE ESCAPE CODE TABLE - Page 17

| OPTIONS |  |  |  |
| :---: | :---: | :---: | :---: |
|  | CODE | LETTER | FUNCTION |
| * | 0 | @ | Visible Cursor Operation |
|  | 1 | A | Blind Cursor Operation |
|  | 2 | B | Plot Via Color Pad |
|  | 3 | C | Transmit Cursor X,Y Position |
|  | 4 | D | Not Used |
| * | 5 | E | Re-Entry Control to BASIC 8001 system |
|  | 6 | F | Sets Unit to Full Duplex |
|  | 7 | G | Not Used |
|  | 8 | H | Sets Unit to Half Duplex |
|  | 9 | I | Not Used |
|  | 10 | J | Sets Unit to Write Vertical |
| * | 11 | K | Sets Unit to Roll up Mode \& write Left to Right |
|  | 12 | L | Sets Unit to Local Mode |
|  | 13 | M | Not Used |
|  | 14 | N | Not Used |
| * | 15 | 0 | Re-Entry control to the CPU Operating System |
| * | 16 | P | Initializing \& Transfers Control to the CPU Operating System |
| * | 17 | 2 | Allows Operation in Character Insert Mode |
|  | 18 | R | Allows Selection of 1 of 7 Baud Rates |
| * | 19 | S | Transfers Control to the 8080 Assembler |
| * | 20 | T | Transfers Control to the Text Editor |
| * | 21 | U | Inserts one line (80 blanks) |
| * | 22 | V | Deletes one line (80 blanks) |
| * | 23 | W | Transfers Control to BASIC 8001 Software |
|  | 24 | X | Sets Unit to Write Left to Right \& Page Mode |
|  | 25 | Y | Test Mode-Fills Screen with Next Character |
|  | 26 | Z | Sets Unit to Write Down on 45 Degrees |
|  | 27 | L (ESC) | Not Used |
|  | 28 | 入 | Sets Unit to Write Up on 45 Degrees |
|  | 29 | ] | Sets Unit to Block Receive Mode |
|  | 30 | $N$ | Causes a Jump to Ram Address 9FAøH |
|  | 31 | - | Transfers Control to the CRT Mode |

The letters are presented for easy reference; i.e., (full duplex mode requires ESC, $F$ sequence). It should be noted that the Escape control codes can be any 8 bit value so long as the 5 least significant bits are correct for the operation desired. The terminal simply masks off the undesired higher order bits. The Keyboard and RS232C input port has separate and independent Flags which determine some of the CRT modes. Therefore, the Keyboard may be in the character input mode while the RS232 input may be in the Plot mode or vise versa. The input port and the Keyboard can operate completely independently of each other. See Details of Escape Codes section for more information.

Code 28 Cursor Up - (Control <br>)
Moves the cursor up one line without destroying any information. This is effectively the opposite of a Line Feed operation.

Code 29 Flag Off - (Control〕)
When this code is received the Reverse Field flag is set to "0". Effects the special character codes (96 to l27) and the color codes (16 to 23).

Code 30 Flag On - (Control $\wedge$ )
When this code is received the Reverse Field flag is set to one. Effects the special character codes (96 to 127) ; the color codes (16 to 23); and the plot modes.

Code 31 Blink On - (Control_ )
When this code is received the Blink bit $A_{6}$ of the composite status is set to a "l".

This bit is turned off when the Blink-Protect-Off key is operated (see Code 15 ).

Code 32 to 95 - Numbers and Letters
These provide the standard printing ASCII Upper Case characters, punctuations and numbers. See Appendix A-2 for code set of the the Intecolor 8001 .

Code 96 to 127 - Special Characters
These codes provide either 32 special characters (such as lower case ASCII characters) or 64 special characters. The 64 special characters are actually two groups of 32 special characters. A group is selected depending upon the condition of the Flag bit. If the flag bit is off then the codes are not changed when they are placed in the CRT refresh RAM. If the flag bit is on then these codes have 96 subtracted from them before they are replaced in the CRT refresh RAM. Therefore they are mapped into 0 to 31 within the CRT refresh memory.
@ or Code 0 - Visible Cursor Mode
This is the terminal's normal mode of operation and it is also the startup state. A received character is placed at the visible cursor location. The cursor then advances to the right one position awaiting the next character. All normal cursor operations are applicable to placing the cursor at a different location.

## A or Code 1 - Blind Cursor Mode

This optional mode provides for a dual cursor operation. That is, normally the host computer will operate in the blind cursor mode and the keyboard in the visible cursor mode. The two modes will not interact with each other. There is also a blind status which may be different than the visible status. The only blind cursor movements allowed are a subset of the cursor $X-Y$ positioning. See Code 3 or control C. This mode allows operation without delay for ASCII TEXT at rates up to 38.4 K Baud.

## B or Code 2 - Character Plot Via Color Pad

When the plot option is installed then this plot mode will be available. It will normally be used via the color pad, but can be used without it. It provides a mix between the Plot Mode and the normal ASCII Character Mode. Instead of responding as described in Character Plot, this mode uses only eight codes to intensify each of the eight blocks within a character. These intensifying codes are the normal color select codes (Control P through W).

This option normally uses the color select pad on the keyboard. The pad is arranged as shown below.

| Black | Blue |
| :--- | :--- |
| Red | Magenta |
| Green | Cyan |
| Yellow | White |

Color Selection
Pad


One
Character Plot Aray

One Plot Block Selected by Green

Figure 2.5.6.1

From the above it is easy to see the one to one correspondence between the 2 x 4 color select pad and the 2 x 4 character plot blocks. Thus, this mode is designed especially for use by the keyboard to simplify the drawing of graphs or the correcting of graphs. Once in this mode a block at the top right hand corner of the cursor present position can be intensified by pushing the top right hand corner key in the color select pad, (in this case the blue key or Control $T$ or Code 20). Once that plot block has been intensified, any other plot block at that same character location can also be intensified since the cursor does not automatically advance. If the blue key was to be pushed the second time, then the already intensified plot block will be extinguished. This effectively allows any one plot block to be erased. After all desired plot blocks have been either intensified or extinguished, the cursor may be conventionally moved without escaping from this special text and character plot mode. In fact, all of the control codes are effective while in this mode except the color select control codes, and any of the ASCII Text characters can be entered and displayed. Any code that requires a two key or more sequence (such as cursor $X-Y, C C I$, and ESC) will terminate the mode. It should be noted that the ASCII Character when entered and displayed advances the cursor as previously done in the visible mode, but the plot blocks (generated by the color pad) do not advance the cursor. Therefore, when a character position has been used to display plot blocks a cursor command must be given to advance the cursor to the next character position.

## C or Code 3 - Transmit Cursor $X, Y$

When this code is selected the terminal sends out the following 7 word sequence:

03, X, Y, 06, Status, ASCII Character, CR.

The $X$ and $Y$ values represents the cursor position on the screen. The status is the status of the ASCII character at that cursor location. The $C R$ may be either a $\varnothing \mathrm{D}$ or an 8 D HEX at the customer request.

This sequence of transmission is the same that the light pen would provide if the unit is so equipped.

E or Code 5 - Re-Entry to BASIC 8001
Return to BASIC 8001 without destroying the BASIC 8001 source program which is in Ram memory.

## F or Code 6 - Full Duplex Mode

When this mode is selected then the Keyboard characters are only sent to the RS232C serial port. They are not processed by the terminal. Therefore, once the unit is put in the full duplex mode via the keyboard, then the only "normal" way the mode can be changed to local or half duplex is via the RS232C serial port. There are two other ways that have been provided to regain local control. One way is to operate the CPU Reset key on the Keyboard, which will initialize the terminal as if power has been just turned on. The other way is to
operate the break key on the keyboard. When this is done a break of 150 MS will be transmitted on the RS232C serial port, and the terminal will be forced into the half duplex mode.

## H or Code 8 - Half Duplex Mode

When this mode is selected then the keyboard characters are not only processed by the terminal but are also sent to the RS232C serial port.

J or Code 10 - Write Vertical Mode
This effects the visible cursor mode only and causes the terminal to enter characters vertically one below the other. All other cursor movements are possible via the cursor mode. After a character is entered the cursor is moved down one character awaiting the next character. Upon reaching the last line the next character will be on the top line, i.e. wrap around occurs.

## K or Code ll - Roll Mode (option 15) Write left to right

When this mode is selected the terminal will cause a page roll up when the last line has been filled. All 48 line units roll two lines at a time while 25 line units roll only one line. Note the plot mode and blind cursor mode only work in non-roll mode. This mode also sets the visible cursor to write left to right.

L or Code 12 - Local Mode
When this mode is selected then the keyboard characters are displayed on the terminal, but they are not sent to the RS232C serial port. In this mode the RS232C serial input port can receive data or change this mode. The terminal can be made to transmit out of the RS232C port, while in the local mode by typing Control $X$ or ESC C.

## O or Code 15 - Re-Entry to CPU Operating System Mode

Causes the same result as Code 16 below but does not reinitialize the $I / O$ Byte or the second RS232C channel Baud rate.

## P or Code 16 - Initialize CPU Operating System Mode

When this optional mode is selected the terminal enters into the CPU Operating System. It then obeys all the commands that are allowed in the CPU Operating System. See the CPU Operating System Manual.

Q or Code 17 - Character Insert Mode
Once in this mode the CRT acts exactly like the normal visible cursor system for all control commands except for those requiring a 2 or more character sequence (such as Cursor XY, CCI, and ESC). When any character is typed or received via the RS232C input, it is inserted within the line at the cursor present position and every character
after the cursor to the end of the line is shifted right one character position. The last character on the line is lost forever. The cursor is also advanced one position. The above is trueexcept for control codes, and "Delete" or (shift '__') keys (code l27).

When the "delete" key is depressed or code 127 is received via the RS232 input port then the character at the cursor present position is deleted and all characters to the end of the line are shifted left one character position. The last character on the line becomes a space. The cursor does not advance.

When the "ESC" key is depressed then the character insert-delete mode is terminated after the second character is selected. The terminal then normally returns to the visible character mode.

## R or Code 18 - Baud Rate Selection Mode

When this mode is entered the unit then accepts the next character as one of seven baud rates. It does this by looking at only the first three bits. Therefore, any 8 bit character that has the desired 3 lower order bits will do. Normally the keyboard numbers 1 to 7 are used. The baud rates and the corresponding numbers are indicated in the table below:

| Number | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Normal Baud <br> Rate | 110 | 150 | 300 | 1200 | 2400 | 4800 | 9600 |
| High Speed <br> Baud Rate | 880 | 1200 | 2400 | 9600 | 19,200 | 38,400 | 76,800 |

The unit is initialized with power up at normally 9600 baud, with one stop bit. This initialized baud rate can be specified by the customer at any of the fourteen above rates when ordered. It should be noted that only in certain modes (blind cursor mode) can the 38,400 Baud be used with delays. In no case can 76,800 Baud be used without delays. The unit may be ordered with either normal baud rates or with the High Speed Baud rates. The two different rate systems cannot be mixed.

The number of stop bits will be determined when the baud rate is set by the condition of the A7 flag. If A7 was on before the rate is selected, 1 stop bit is selected; if A7 was off before the rate is selected 2 stop bits are selected.

S or Code 19 - 8080 Assembler Mode
When this optional mode is selected the terminal enters into the 8080 Assembler Mode. It then obeys all the commands that are allowed in the 8080 Assembler. At present this option is not available.

T or Code 20 - Text Editor Mode
When this optional mode is selected the terminal enters into the Text Editor Mode. It then obeys all the commands that are allowed in the Text Editor. At present this option is not available.

U or Code 21 - Insert Line Mode
When this mode is selected the cursor moves to the beginning of the line it is presently on and this line and all lines to the end of the page is shifted down by one line. Then a new line of 80 spaces ( or blanks) are inserted with the cursor remaining at the beginning of that new line.

Normally the cursor will be at the beginning of the line to be inserted when this mode is used. After a line has been inserted the terminal returns to the normal visible character mode.
$\underline{V}$ or Code 22 - Delete Line Mode
When this mode is selected the cursor moves to the beginning of the line it is presently on and this line is deleted. All lines to the end of the page are shifted up by one line. Then a new line of 80 spaces ( or blanks) are inserted at the bottom of the page. The cursor will remain at the beginning of the line that had been deleted. After a line has been deleted the terminal returns to the normal visible character mode.

W or Code 23 - BASIC 8001 Language Mode
When this optional mode is selected the terminal enters into the BASIC 8001 Languaqe mode. It then obeys all the commands that are allowed in Basic 800l. See the "BASIC 8001 Manual".

X or Code 24 - Page Mode Write Left to Right
When this mode is selected the terminal will not roll up when the last line has been filled, but will begin at home again. The terminal is also placed in the write left to right mode. This is the normal power up mode. This mode affects all modes that use the visible cursor. The blind cursor and plot modes will only operate in the page mode.

Y or Code 25 - TEST Mode
When this mode is selected the next character that follows causes the complete screen to be filled with that character. Note use ESC, $\underline{Y}$, - for a convergence test pattern.

When this mode is selected the terminal will place the character at the present visible cursor and will then cause a cursor right followed by a line feed to occur. Therefore, the next character entered will be to the right and down one position from the previous character. When the bottom of the page is reached the next character will appear on the top of the screen, i.e., wrap around occurs.

$$
\text { L or ESC or Code } 27 \text { - No Effect Code }
$$

Performs a return to visible character mode.
\or Code 28 - Write Up 45 Mode
When this mode is selected the terminal will place the character at the present visible cursor and will then cause a cursor right followed by a cursor up to occur. Therefore, the next character entered will be to the right and up one position from the previous characters. When the top of the page is reached the next character will appear on the bottom of the screen, i.e., wrap around occurs.

## J or Code 29 - Block Receive Mode

Causes the unit to enter into the block receive mode. Uses the blind cursor to position the data. Looks for a (FF), (OO) HEX sequence to terminate back to the visible cursor mode. Note this is same format as when control (x) or page transmit is requested. Note page transmit starts at visible cursor and ends at end of page or when an ( FF ), ( OO ) HEX sequence is found.
^ or Code 30 - Jump to RAM 9FAøH
When this code is received the CRT O.S. branches to location 9FAØH. Therefore, the user must patch into RAM address 9FAøH a jump to his program.
— or Code 31 - Transfers Control to the CRT Operating System
When this code is received, the unit is forced to the CRT O.S. mode. If Option 34, the CPU O.S.r is also installed, then a message will be printed saying:

YOU ARE NOW IN THE 8001 CRT MODE

## DETAIL OF GRAPHIC PLOT SUBMODES

Code 2 Graphic Plot Mode - (Control B) - (Option 02)
The general Graphic Plot Mode is entered by a binary code 2 or a Control Code B. (See Appendix B). It should be noted that the XY Plot Mode is also entered at the same time. If a plot mode other than XY Point Plot is desired, the next word that follows should then be a binary code from 240 to 255 . These codes represent the various plot submodes as shown in the summary of Graphic Plot Submodes.

An additional feature is available to allow a graphic plot to be erased by simply setting the Flag bit on before entering the plot mode. This causes an XOR function to exist when plotting. Therefore, if you plot the same point, bar or vector twice, the second time erases the original.

Once in the general Plot Mode, any of the plot submodes may be entered by sending the corresponding code to the terminal. When this code is received, a flag internal to the terminal, known as PLOFL, is set placing the terminal in the appropriate plot submode. It should be noted that in many of the plot submodes, PLOFL is automatically set to a different value upon completion of the operation of that submode causing the terminal to enter a new submode. This is done to make coding and operation of the terminal in the various plot functions easier for the operator. The various submodes and their interactions are explained in detail in Appendix B.

In addition to being able to enter the plot submodes from the general Plot Mode, any plot submode may be entered from any other plot submode with the exception of the Character Plot Mode.

Colors may be defined on a character by character basis only and the color of an individual plot block as well as all other intensified plot blocks within a character will be the most recent color defined when a new block is intensified in that character. To change a color, it is required that the Plot Mode or plot submode be terminated, the color changed, and the Plot Mode be re-entered.

The character grid is made up of 80 characters wide by 25 [48] lines high. The 0 reference point for all plotting is always the lower left corner. Each character is further broken up into 2 blocks wide by 4 blocks high which then causes the plot grid to be 160 blocks wide by 100 [192]blocks high. All plot submodes operate on this size grid and have the same reference point. Positive direction is considered up and to the right and negative direction is considered down and to the left.

All plot submodes and the general Plot Mode are terminated or exited by the binary code, 255. Whenever this code is received, the modes are terminated and must be re-entered as described above.

Appendix B-2 gives a convenient summary of the codes required to enter the Plot Mode and the various plot submodes as well as the status of PLOFL before and after each operation and the ranges of each operation.
Plot Mode - Escape - (255 binary)

This code is used to exit from the Plot Mode or any of the plot submodes. The control "?"or Fl5 is used to escape from the Plot Mode from the Keyboard.

## Character Plot - (254 binary)

The Character Plot is entered by a 254 after the general Plot Mode, "2" or Control Code B, is entered. From the Keyboard use Control " $>$ " or Fl4. It may also be entered directly from any of the other plot submodes. After entering the Character Plot, the next word will be treated as a plot character except for code 255 binary or (FF) hexadecimal (i.e. all eight bits are "l's"). See Appendix B-

The general Plot Mode and the Character Plot terminate upon receipt of a 255 code. The above procedure must be repeated after a 255 code terminates the Plot Mode and the plot submodes.

Other plot submodes may not be entered from the Character Plot. To enter other plot submodes, the Character Plot must be terminated, the general Plot Mode entered and the plot submode entered with its associated code.

The procedures for entering and exiting the Character Plot are shown below.

| Function | Code |
| :---: | :---: |
| Plot Mode | 2 |
| Character Plot | 254 |
| Plot Character 1 | 0 to 254 |
| - |  |
| - |  |
| - |  |
| Plot Character n | 0 to 254 |
| Plot Escape | 255 |

The Character Plot causes the 6 wide by 8 high dot matrix to be divided into 8 blocks organized 2 blocks wide by 4 blocks high. Each block consists of a sub-dot matrix of 3 dots wide by 2 dots high. Each block may be individually intensified by defining the bit (one of eight bits) associated with the block in the plot character. Bits may by "ORed" together for a combination of blocks in a plot character, creating a form of graphics for plotting data or drawing diagrams. Large characters may also be created by utilizing the blocks of several character positions to create a large 5x7 dot matrix.

$$
\underline{X} \text { Point Plot }- \text { (binary 253) }
$$

The X Point Plot is automatically entered upon receipt of the general Plot Mode code, binary code 2 , or Control Code B. It also may be entered directly from any of the other plot submodes except Character Plot. From the Keyboard use Control " = " or Fl3. After entering the $X$ Point Plot, the next word defines the $X$ value of the block that is desired to be plotted See Appendix BThe $X$ value in this mode may range from binary 0 to 159 and all other values will cause 160 to be subtracted and the resultant value of $X$ to be computed.

The $X$ Point Plot may be terminated by code 255 which causes the general Plot Mode to be terminated also. Any of the other plo submodes may be entered directly from the $X$ Point Plot by simply entering the appropriate plot submode codes which range from binary 240 to 254.

It should be noted that this mode does not cause a block to be intensified, but only causes the $X$ value to be defined. Once the $X$ value is sent, the terminal is automatically placed in the $Y$ Point Plot mode. Thus, the next code sent will be the $Y$ value, which may range from binary 0 to 99 [0-191]. Upon receipt of the $Y$ value, a plot block will be intensified on the CRT screen at the $X$ value and $Y$ value intersection. The terminal is then automatically placed in the $X$ Point Plot mode and the next word sent will be interpreted as an $X$ value.

Therefore, once in the $X$ Point Plot mode, new blocks may be defined by simply sending $X$ values and $Y$ values consecutively, without the necessity of re-entering the $X$ or the $Y$ Point Plot modes.

The procedures for entering and exiting the $X$ Point Plot mode are shown below:


* Automatically X Point Plot mode also

NOTE: SEND Code 253 between X,Y data sets if necessary for timing considerations. See Appendix A-4 for delays.

The X Point Plot in conjunction with the Y Point Plot allows any block on a 160 wide by 100 ( 192 for 48 Line) high block matrix to be positioned to and intensified. If the new block is within a character position that is a previously intensified ASCII character, then the ASCII character is replaced completely by the new block and its associated color.

## Y Point Plot - (binary 252)

The $Y$ Point Plot is entered by a binary 252 code after the general Plot Mode is entered. See Appendix BFrom the Keyboard use Control " $<$ " or Fl2. It may also be entered directly from any of the other plot submodes except Character Plot (binary 254). It is more commonly entered automatically from the $X$ Point Plot mode. After entering the Y Point Plot, the next word defines the $Y$ value of the block that is desired to be plotted and causes the block to be intensified in accordance with the Section on (binary 253). The Y value in this mode may range from binary 0 to 99 (0-191) and all larger values will cause 100 (192) to be subtracted from the new value of $y$ to be calculated.

Upon receipt of the $Y$ value, the $X$ Point Plot is automatically entered by the terminal. The $X$ value of the next block to be plotted may then be sent as explained in the Section on (binary 253).

The Y Point Plot is terminated by Code 255 which causes the general Plot Mode to be terminated also. Any of the other plot submodes may be entered directly from the $Y$ Point Plot by simply entering the appropriate plot submode codes which range from binary 240 to 254.

Therefore, once in the $Y$ Point Plot mode, new points may be defined by simply sending $X$ values and $Y$ values consecutively without the necessity of re-entering the $X$ or the $Y$ Point Plot modes. The procedures for entering and exiting the $Y$ Point Plot mode are shown below:

Function
Plot Mode
Plot Submode
$Y_{1}$ Value*
$\mathrm{X}_{2}^{1}$ Value
$Y_{2}$ Value
-
-
-
$x_{n}$ Value $\quad 0$ to 159
$Y_{n}$ Value 0 to 99
Plot Excape 255
or
Plot Submode

## Code

2
252
0 to 99
0 to 159
0 to 99
or
240 to 254

* Plots point using whatever previous X Value left in memory.

NOTE: Send Code 253 between $\mathrm{X}, \mathrm{Y}$ data sets if necessary for timing considerations. See Appendix A-4 for Delays.

XY Incremental Point Plot - (binary 251.)
The XY Incremental Point Plot is entered by code 251 after the general Plot Mode is entered. From the Keyboard use Control ";" or Fll. It may also be entered directly from any of the other plot submodes, except Character Plot. After entering the XY Incremental Point Plot mode, the next word defines the next two increments as shown in Figure below. This word may have a range from binary 0 to 239 since binary 240 to 255 is used for the plot submode codes.

| $\mathrm{b}_{7}$ | $\mathrm{b}_{6}$ | $\mathrm{b}_{5}$ | $\mathrm{b}_{4}$ | $\mathrm{b}_{3}$ | $\mathrm{b}_{2}$ | $\mathrm{b}_{1}$ | $\mathrm{b}_{0}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\Delta{ }^{1}$ |  | $\Delta Y_{1}$ |  | $\Delta \mathrm{X}_{2}$ |  | $4 Y_{2}$ |  |
|  | Plock |  |  |  |  |  |  |

$$
\begin{array}{cll}
\begin{array}{ll}
b \\
n+1 \\
0 &
\end{array} & \begin{array}{l}
\text { b } \\
n
\end{array} & \\
1 & 0 & \text { No Change } \\
0 & 0 & \text { Positive Increment } \\
1 & 1 & \text { Negative Increment } \\
\mathrm{n}=0,2,4,6 & & \text { No Change } \\
& &
\end{array}
$$

If $\mathrm{b}_{\mathrm{o}}$ through $\mathrm{b}_{3}$ are "0"s, then the plot block will not print but will increment one increment according to the coding of $\mathrm{b}_{4}$ through $\mathrm{b}_{7}$. This allows the user to easily "skip" a plot increment by plotting with an invisible block.

It should be noted that the XY Incremental Plot mode does not automatically transfer the terminal to any other plot submode upon receipt of an incremental change word,but remains in the $X Y$ Incremental Plot mode ready to receive another incremental change word. Therefore, a series of incremental movements may be made by sending consecutive incremental change words.

The XY Incremental Plot mode may be terminated by code 255 which causes the general Plot Mode to be terminated also. Any of the other plot submodes may besentered directly from the XY Incremental Point Plot by simply entering the appropriate plot submode codes which range from binary 240 to 254.

The procedures for entering and exiting the XY
Incremental plot mode are shown below:

| Function | Code |
| :---: | :---: |
| Plot Mode | 2 |
| or | or |
| Plot Submode | 240 to 253 |
| XY Incremental |  |
| Point Plot | 251 |
| Incremental Change |  |
| Word 1 | 0 to 239 |
| - |  |
| - |  |
| - |  |
| Incremental Change |  |
| Word n | 0 to 239 |
| Plot Escape | 255 |
| or | or |
| Plot Submode | 240 to 254 |

NOTE: Send code 251 between XY incremental point words if necessary for timing considerations. See Appendix A-4 for Delays.

The X Bar Graph, $X_{o}$ Value is entered by a binary 250 code after the general Plot mode is entered. From the Keyboard use Control ":" or Fl0. It may also be entered from any of the other plot submodes except Character Plot. After entering the X Bar Graph, $X_{O}$ Value Mode, the next word sent defines the $X_{O}$ Value or the left horizontal start block of the horizontal bar graph. The graph grid is referenced to the lower left hand corner of the face of the CRT. The $\mathrm{X}_{\mathrm{O}}$ may range in value from 0 to 159 and all other values have 160 subtracted and the new value calculated for $X_{0}$.

Upon receipt of the $X_{o}$ Value, the value of $X_{o}$ is stored in memory and the terminal is automatically placed in the $X$ Bar Graph, Y Value mode (binary 249). The terminal is now ready to receive the next eight bit word as the $Y$ position of the bar graph. Upon receipt of the $Y$ value, the terminal is then automatically placed in the $X$ Bar Graph, $X$ Max Value mode (binary 248). The terminal is now ready to receive the next eight bit word as the X Max Value. Upon receipt of the $X$ Max Value, the bar is drawn on the CRT and the terminal is placed back into the $X$ Bar Graph, $Y$ Value mode (binary 25l) ready to receive a new $Y$ value to begin the bar graph drawing process over again as outlined above. This process is shown below and in Appendix B.

| Function | Code |
| :---: | :---: |
| Plot Mode | 2 |
| or | or |
| Plot Submode | 240 to 253 |
| X Bar Graph, X O Value | 250 |
| $\mathrm{X}_{\mathrm{O}}$ Value Word 1 | 0 to 159 |
| Y Value Word l | 0 to 99 (0-191) |
| X Max Value Word l | 0 to 159 |
| Y Value Word 2 | 0 to 99 (0-191) |
| X Max Value Word 2 | 0 to 159 |
| - |  |
| - |  |
| Y Value ${ }^{\text {- }}$ |  |
| Y Value Word n | 0 to 99 (191) |
| X Max Word n | 0 to 159 |
| Plot Escape | 255 |
| or | or |
| Plot Submode | 240 to 254 |

NOTE: Use Code 251 between $Y$ value, $X$ max Value data sets for timing considerations. Timing delays depends directly upon the length of the bar being intensified. See Appendix A-4 for delays both minimum and maximum.

As can be seen from the above process, once in the $X$ Bar Graph, $X_{O}$ mode, it is necessary to send only two words, $Y$ and $X$ Max, to completely define other bar graphs with the same $X_{O}$ in the horizontal direction. As before, any of the submodes can be entered independently. After the first bar graph sequence, additional bar graphs can be described by a new $Y$ position for the graph and a new $X$ Max

Value for the graph. The bar is drawn after the $X$ Max Value is received using the original value of $X_{0}$.

Any of the other plot submodes may be entered directly from the X Bar Graph, entering the appropriate plot submode codes which range from binary 240 to 254.

This mode allows bar graphs in any color or multiple colors to be drawn with a width as small as one plot block wide or areas under curves may be easily filled in.

X Bar Graph, Y Value - (binary 249)

The X Bar Graph, Y Value is entered by a binary 249 code after the general Plot Mode is entered. From the Keyboard use Control "9" or F9. It is more commonly entered from the X Bar Graph, $x_{0}$ Value automatically, and may also be entered from any of the other plot submodes except Character Plot (binary 254). After entering the X Bar Graph, Y Value mode, the next word sent defines the $Y$ or vertical position of the horizontal bar graph being drawn. The Y value may range from binary 0 to 99 ( 0 to 191) and all other values will have 100 (192) subtracted from it and the new value calculated for the $Y$ value.

Upon receipt of the $Y$ value word, the value of $Y$ is stored in memory and the terminal is automatically placed in the $X$ Bar Graph, X Max Value mode, as explained more completely in the Section on (binary 248).

Any of the other plot submodes may be entered directly from the X Bar Graph, Y Value mode by simply entering the appropriate plot submode codes which range from binary 240 to 254.

X Bar Graph, X Max Value - (binary 248)
The X Bar Graph, X Max Value is entered by a binary 248 code after the general Plot Mode is entered. From the Keyboard use Control "8" or F8. It is more commonly entered from the X Bar Graph, $Y$ Value automatically, and may also be entered from any of the other plot submodes except Character Plot. After entering the X Bar Graph, X Max Value mode, the next word received defines the $x$ Max horizontal point of the horizontal bar graph being drawn. The $X$ Max Value may range from 0 to 159 and all other values will have 160 subtracted from it and the new value calculated for X Max Value.

Upon receipt of the $X$ Max Value word, the bar graph is drawn in the predefined color on the face of the CRT according to the $X_{O}$ and $Y$ value stored in memory from previous operations. The terminal is then automatically placed in the X Bar Graph, Y Value mode, binary 249, for the beginning of a new bar graph as more completely explained in the Section on (binary 248).

Any of the other plot submodes may be entered directly from the X Bar Graph, X Max Value mode by simply entering the appropriate
plot submode codes which range from binary 240 to 254.

```
X Incremental Bar Graph - (binary 247)
```

The X Incremental Bar Graph is entered by a binary 247
code after the general Plot Mode is entered. From the Keyboard use Control "7" or F7. It may also be entered from any of the other plot submodes except Character Plot. After entering the X Incremental Bar Graph mode, the next word sent defines the next two horizontal and vertical increments for two horizontal bar graphs. Thus, one may position a bar graph each side of the present location and add or subtract an increment to the bar graph previously defined. The coding and composition is the same as explained in the Section on (binary 25l). An example is shown in Appendix B-6.
$\underline{\text { Y Bar Graph, } Y_{0} \text { Value }-(b i n a r y ~ 246) ~}$
The Y Bar Graph, $Y_{O}$ Value is entered by a binary 246 code after the general Plot Mode is entered. From the Keyboard use Control "6" or F6. It may also be entered from any of the other plot submodes except Character Plot. After entering the Y Bar Graph, Yo Value mode, the next word sent defines the $Y_{o}$ or the vertical start point of the vertical bar graph being drawn. The range of the $Y_{0}$ word is 0 to 99 (0-191) and all other values have 100 (192) subtracted and will have the new value calculated for $Y_{O}$ Value.

All other operations are identical as explained in the Section on (binary 250), X Bar Graph, X ${ }_{0}$ Value except that Y Bar Graph, X Value and Y Bar Graph, Y Max Value are applicable for drawing vertical bar graphs. An example is shown in Appendix B-5.
$\underline{Y}$ Bar Graph, X Value - (binary 245)
The Y Bar Graph, X Value is entered by a binary 245
code after the general Plot Mode is entered. From the Keyboard use Control "5" or F5. It is more commonly entered from the Y Bar Graph, $Y_{o}$ Value automatically, and may also be entered from any of the other plot submodes except Character Plot. After entering the Y Bar Graph, X Value mode, the next word sent defines the $X$, or horizontal position of the vertical bar graph being drawn. The $X$ Value may range from 0 to 159 and all other values will have 160 subtracted and will have the new value calculated for the $X$ value.

All other operations are identical as explained inthe Section on binary 249, X Bar Graph, Y Value except that Y Bar Graph, YO Value and Y Bar Graph, Max Value are applicable for drawing vertical bar graphs. An example is shown in Appendix B-5.

> Y Bar Graph, Y Max Value - (binary 244)

The Y Bar Graph, Y Max Value is entered by a binary 244 code after the general Plot Mode is entered. From the Keyboard use Control "4" or F4. It is more commonly entered from the $Y$ Bar Graph, $X$

Value automatically, and also may be entered from any of the other plot submodes except Character Plot. After entering the Y Bar Graph, $Y$ Max Value mode, the next word received defines the vertical $Y$ Max point of the vertical bar graph being drawn. The $Y$ Max Value may range from binary 0 to 99 (0-191) and all other values will have 100 (192) subtracted and will have the new value calculated for $Y$ Max Value.

All other operations are identical as explained inthe Section on (binary 248), X Bar Graph, X Value, except that Y Bar Graph, Y Value and Y Bar Graph, X Value are applicable for drawing vertical bar graphs. An example is shown in Appendix B-5.

> Y Incremental Bar Graph - (binary 243)

The Y Incremental Bar Graph is entered by a binary 243 code after the general Plot Mode is entered. From the Keyboard use Control "3" or F3. It may be entered from any of the plot submodes except Character Plot. After entering the $Y$ Incremental Bar Graph mode, the next word sent defines the next two horizontal and vertical increments for two vertical bar graphs.

All other operations are identical as explained in the Section on (binary 247), X Incremental Bar Graph except for the mode being applicable for drawing vertical bar graphs. An example is shown in Appendix B-6.

Vector Mode $X_{0}$ Value - (binary 242)
The Vector Mode is entered by a binary 242 code after the general Plot Mode is entered. From the Keyboard use Control "2" or F2. It may be entered from any of the plot submodes except Character Plot. After entering the Vector Mode, $X_{O}$ Value, the next word defines the $X_{0}$ Value point of the vector being drawn.

The Vector Mode requires the two end points to be defined (i.e. $X_{O} Y_{O}$ and $\left.X_{1} Y_{1}\right)$. The $X_{1}, Y_{1}$ values should previously be defined by way of the $X, Y$ Point Plot Mode.

Upon receipt of the $X_{O}$ Value the terminal is automatically placed in the Vector $Y_{0}$ Value Mode (binary 24l). The terminal is now ready to receive the next eight bit word as the $Y_{0}$ Vector Value. Upon receipt of the $Y_{O}$ Value the terminal then determines the best straight line fit between $X_{O}, Y_{O}$ and $X_{l}, Y_{l}$ using the plot blocks. The terminal will then revert to the Vector Mode $X_{0}$ value (binary 242), ready to receive the new $X_{0}$ Value for another vector. The process is shown below and in Appendix B-7.


NOTE: Send code 242 between $Y_{O}$ vector point and $X_{0}$ vector point words if necessary for timing considerations. See Appendix A-4 for delays.

Vector Mode $Y_{0}$ Value - (binary 241)
The $Y_{0}$ vector is entered by binary 241 code after the general Plot Mode is entered. From the keyboard use Control "l" or Fl. This mode is more commonly entered automatically from the $X_{O}$ Vector mode. After entering the $Y_{O}$ Vector mode, the next word defines the $Y_{0}$ value of the vector being drawn. There is no restriction on $Y_{0}$ with respect to $Y_{1}$ except it must be in the range of 0 to 99 (191). Upon receipt of the $Y_{O}$ value a vector is drawn from $X_{1}, Y_{1}$ to $X_{O}, Y_{O}$, with the new $X_{1} Y_{1}$ now at the old $X_{0} Y_{O}$. If the next vector has a $X_{1} Y_{1}$ value $=$ $X_{O} Y_{O}$ old, then only the new $X_{O} Y_{O}$ need be sent. This would effectively draw a vector from the present $X_{O} Y_{O}$ position to the new $X_{O} Y_{O}$ point. See Appendix B-7.

$$
\underline{X}_{O} \underline{Y}_{0} \text { - Incremental Vector Mode - (binary 240) }
$$

The $X_{O}-Y_{O}$ incremental vector mode is entered by a binary 240 code after the general plot mode is entered. From the keyboard use control " $\varnothing$ " or $F \varnothing$. It may also be entered from any of the other plot submodes except Character Plot. After entering the incremental vector mode, the next word sent defines the increments in $X_{O}, Y_{O}, X_{l}$ and $Y_{1}$ point values for the vector from $X_{1} Y_{1}$ to $X_{O} Y_{0}$. This word may have a range from binary 0 to 239 since binary 240 to 255 are used for the plot submode codes.

Referring to the section on (binary 251), XY Incremental Point Plot it can be seen that there is one two bit element available for each of the 4 points (i.e. $X_{O}, Y_{O}, X_{1}$ and $Y_{1}$ ). The $\Delta X_{1}, \Delta Y_{1}$ refers to the increment in $X_{1}, Y_{1}$ of the vector and the $\Delta X_{2}, \Delta Y_{2}$ refers to the increment in $X_{O}, Y_{0}$ of the vector.

| $\mathrm{b}_{7}$ | $\mathrm{~b}_{6}$ | $\mathrm{~b}_{5}$ | $\mathrm{~b}_{4}$ | $\mathrm{~b}_{3}$ | $\mathrm{~b}_{2}$ | $\mathrm{~b}_{1}$ | $\mathrm{~b}_{0}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $\mathrm{X}_{1}+1$ | $\mathrm{X}_{1}-1$ | $\mathrm{Y}_{1}+1$ | $\mathrm{Y}_{1}-1$ | $\mathrm{X}_{\mathrm{O}}+1$ | $\mathrm{X}_{0}-1$ | $\mathrm{Y}_{\mathrm{O}^{+1}}$ | $\mathrm{Y}_{\mathrm{O}^{-1}}$ |

Therefore, if $\mathrm{b}_{4}$ and $\mathrm{b}_{5}$ are both $l$ or both $\varnothing$ then no increment will take place. If either half of the word is all zero then the corresponding $X, Y$ will be changed but no vector will be drawn. This allows the user to easily "skip" points. The only time a vector will be drawn is when both halfs of the word are non zero.

The incremental vector plot mode does not automatically transfer control to any other mode. It remains in this incremental mode until terminated by a plot submode code. Therefore a series of incremental movements in both $X_{O}, Y_{O}$ and $X_{1} Y_{1}$ may be made by sending consecutive incremental change words.

The procedure for entering and exiting the XY Incremental plot mode are shown below:

```
Function
Plot Mode
    or
Plot Submode
Incremental Vector
            Plot Mode
Incremental change
        in }\mp@subsup{X}{1}{},\mp@subsup{Y}{1}{},\mp@subsup{X}{O}{\prime},\mp@subsup{Y}{O}{
            Word l 0 to 239
                .
                .
                .
                    •
            Word N O to 239
Plot. Escape
                or
Plot Submode 240 to 254
                                    255
                                or
```

NOTE: Send code 240 between incremental vector words if necessary for timing considerations. See Appendix A-4 for input Delay Times.

The Intecolor 8001 Light Pen is designed to move the cursor on the screen of the terminal by simply pointing to the desired location on the screen and touching with the forefinger the touch-sensitive end of the light pen. The touch sensitive end of the light pen acts as a switch or button.

To effect operation of the light pen, the pen is simply pointed to the desired location on the screen. Either the standard lense or the long range lense may be used in the same manner. When the desired location is reached, the forefinger is placed on the touch-sensitive end of the pen and held on the pen until the cursor on the screen resides at the location the pen is pointing to. As long as the finger is kept on the pen the cursor will follow the pen to any location.

When the cursor is at the desired location, lift the forefinger from the tip of the pen and the following 7 word sequence will be transmitted to the Jl RS232 output port.

| 03 | Cursor X-Y (See Code 3) |
| :--- | :--- |
| $X$ | $X$ Cursor Coordinate |
| $Y$ | Y Cursor Coordinate |
| 06 | CCI (See Code 6) |
| Status | Status Character $\quad$ (See Appendix A-6) |
| ASCII or |  |
| Special |  |
| Character |  |
| 8D |  |
|  |  |

Notice that this sequence is not transmitted unless the finger first touches the end of the pen in the touch sensitive area and is effected when the finger is lifted from the end of the pen.

Note that a blue flood is normal operation and occurs every time the touch sensitive end of the pen is touched by the forefinger and will repeat at a 2 cps rate until the finger is lifted.
$3-1$
 AS SHOWN DOTTED.

THIS FUNCTION is ASCII
$C=C H R \$(C): ? C, C \$$

65193 A (UPPer case 14143 PLOT SMALL. CHARACTERS
64 192 @
63191 ?
13 14, CURSOK LEFT
$62190>$
12140 ERASE \& HOME.
$61189=$
60188 く
59187 ;
$58: 86:$
571859
$48 \quad 176$ (1)
$47175 /$
46174.
$45173-$
44172 ,
$43171+$
42170 *
41169)

40168 (
$391^{167}$
38 166 \&
$37165 \%$
$36164 \$$
35163 丮
$34162 "$
33161 !
$160 ?$
II 139 ERASE LINE, CURSOR LEFT.
10 138 CURSOR DOWN I WINE.
9137 TAB (1)
8 136 CURSOR HOME
7135 BEEP, BEEP!
6 134 RED ON RED
5 133 HALT
4 132 HALT
3131 ? DOWN 1/2 PAGE
2 130 $\because$ CRASH


Repeats all things like 128 greater.

KEYBOARD INPUTS: $I=\operatorname{INP}(1)$ is NoT ASCII!
$\left.\begin{array}{l}\phi=64 \\ =65 \\ \downarrow=66 \\ \downarrow=8 \\ q=73\end{array}\right\}$ must sub ic OR ADD $112\left\{\begin{array}{l}\phi=176 \\ \downarrow \\ \downarrow=185 \\ q=1\end{array}\right\}$
$9=73$
$Q=128$
$A=129$
$B=139$
$\left.\begin{array}{l}C=131 \\ D=132\end{array}\right\}$ must sub 64 OR ADD 64 $\left\{\begin{array}{l}a=192 \\ A=193 \\ 1\end{array}\right\}$
$E=133$
$\downarrow=1$
$Z=202$
WHITE $=215$ $\begin{array}{ll}\text { WHAT }= & 215 \\ \text { CYAN } \\ \text { MAGENTA } & 214 \\ 213\end{array}$
$\left.\begin{array}{l}\text { BLUE }=212 \\ \text { YELLOW }=211\end{array}\right\}$ MUST SUB TE OR SUB 64
YELLOW $=211$
$\begin{array}{ll}\text { GREEN } & =210 \\ \text { RED } & =209\end{array}$
BLACK $=298$

$$
\left.\begin{array}{l}
=160 \\
a=161 \\
b=162 \\
c=163 \\
\downarrow=\downarrow \\
z=234
\end{array}\right\} \mathrm{Mu}
$$

MUST SUB 64 OR ADD 64


[^1]㫛昌圆 CHARACTER PLACEMENT GIVES LEFT\＆BOTTOM BOARDERS WITHIN THE $6 \times 8$ MATRIX．


A-3
INPUT CODE FLOW DIAGRAM

$$
\begin{aligned}
& 7 \\
& 7 \\
& 7 \\
& 7
\end{aligned}
$$




7


7
7
7


7

$\underline{\text { STANDARD }^{\text {INTECOLOR }}} \mathrm{R}$ 8001

| $\mathrm{A}_{7}$ | $\mathrm{~A}_{6}$ | $\mathrm{~A}_{5}$ | $\mathrm{~A}_{4}$ | $\mathrm{~A}_{3}$ | $\mathrm{~A}_{2}$ | $\mathrm{~A}_{1}$ | $\mathrm{~A}_{0}$ |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | RED FOREGROUND |
| 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | GREEN FOREGROUND |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | BLUE FOREGROUND |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | FOREGROUND BLINK |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | PLOT CHARACTER |

WITH BACKGROUND COLOR OPTION

| $\mathrm{A}_{7}$ | $\mathrm{~A}_{6}$ | $\mathrm{~A}_{5}$ | $\mathrm{~A}_{4}$ | $\mathrm{~A}_{3}$ | $\mathrm{~A}_{2}$ | $\mathrm{~A}_{1}$ | $\mathrm{~A}_{0}$ |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | RED FOREGROUND |
| 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | GREEN FOREGROUND |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | BLUE FOREGROUND |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | RED BACKGROUND |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | GREEN BACKGROUND |
| 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | BLUE BACKGROUND |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | FOREGROUND BLINK |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | FLOT CHARACTER |

The above codes may be "ORed" for composite functions

$$
A-5
$$

$$
\begin{aligned}
& 7 \\
& 7 \\
& 7 \\
& 7 \\
& 7 \\
& 7
\end{aligned}
$$

| 1. AA | 1. AA | 10. IN OA ${ }^{\text {Key }}$ | 10. IN OC |
| :---: | :---: | :---: | :---: |
| 2. $B A$ | 2. BA | 9. IN lA Data | 9. IN 1C |
| 3. BB | 3. BB | 21. In 2A Bits | 21. IN 2C |
| 4. CP | 4. 470 ohms to +12 V | 8. IN 3A l-4 | 8. IN 3C |
| 5. CB | 5. | 20. IN 4A Control | 20. IN 4C |
| 6. | 6. | 7. IN 5A Shift | 7. IN 5C |
| 7. AB | 7. $A B$ | 19. IN 6A Key Data B5 | 19. IN 6C |
| 8. | 8. | 6. In 7A Key Data B6 | 6. IN 7C |
| 9. RX Response Control | 9. RX Response Control | 12. In 4B Not | 2. OUT OC |
| 10. RX Responce Control | 10. RX Response Control | 23. IN 5B Used | 14. OUT İC |
| ll. | 11. CLR+ | 11. IN 6B | 3. OUT $\overline{2} \mathrm{C}$ |
| 12. TTL TX | 12. TTL TX | 22. In 7B Key Trigger | 15. OUT $\overline{3} \mathrm{C}$ |
| 13. | 13. TX Isolator input | 2. OUT $\bar{O} A \quad \mathrm{RX}$ ACK | 4. OUT $\overline{4} \mathrm{C}$ |
| 14. | 14. | 14. OUT $\overline{1} \mathrm{~A}$ | 16. OUT 5c |
| 15. | 15. | 3. OUT $\overline{2} A$ | 5. OUT ठС |
| 16. | 16. | 15. OUT $\overline{3} \mathrm{~A}$ | 17. OUT 7 C |
| 17. | 17. | 4. OUT $\overline{4} \mathrm{~A}$ | 13. CPU RESET |
| 18. | 18. CLR- | 16.*OUT 5A | 24. 2nd RS232 TX |
| 19. | 19. | 5. OUT $\overline{6 A}$ Bell | 1l. 2nd TTL TX |
| 20. CD | 20. 470 ohms to +12 V | , 17. OUT $\overline{7}$ A -Кеу АСК | 22. 2nd RS232 RX |
| 21. | 21. CLT+ | 13. CPU RESET | 12. +12 V |
| 22. | 22. | 1. SN -Key Inturr. | 23. -12V |
| 23. | 23. | 25. +5 V | 25. +5v |
| $\begin{aligned} & 24 . \\ & 25 . \end{aligned}$ | 23. | 18. GND | 18. GND |
|  | 25. CLT- |  | 1. SN - EXT Inturr. |
|  | An external jumper is required from pin 12 to pin 13. |  |  |
|  | A 2.2 K ohm register is required from pin 3 to pin 4. | $\cdots$ |  |
| STANDARD TTY | OPTIONAL |  |  |
| ElA RS2326 | 20MA Current Loop. |  |  |
|  |  | J2 | J3 |
| SERIAL INPUT/OUTPUT |  | KEYBOARD | OPTIONAL: <br> PARALLEL INPUT/OUTPUT AND 2nd PS232C |
| $A-7$ <br> 1/O Connector |  |  |  |

$$
\square
$$


$B-1$
PLOT MODE FUNCTIONS
 symbols.

## $-1 \quad-1$







$$
B-5
$$



Appendix B6


B-7
$\mathrm{X}_{0} \mathrm{Y}_{0}$ Vector Plot Mode

APPENDIX C

m


๓

# Appendix C <br> TMS 5501 Multifunction Input/Output Controller <br> TABLE OF CONTENTS 

1. INTRODUCTION
1.1 Description ..... 2
1.2 Summary of Operation ..... 3
2. OPERATIONAL AND FUNCTIONAL DESCRIPTION
2.1 Interface Signals ..... 6
2.2 TMS 5501 Commands ..... 8
2.2.1 Read Receiver Buffer ..... 9
2.2.2 Read External Input !ines ..... 9
2.2.3 Read Interrupt Address ..... 9
2.2.4 Read TMS 5501 Status ..... 9
2.2.5 Issue Discrete Commands ..... 10
2.2.6 Load Rate Register ..... 11
2.2.7 Load Transmitter Buffer ..... 12
2.2.8 Load Output Port ..... 12
2.2.9 Load Mask Register ..... 12
2.2.10 Load Timer n ..... 12
3. TMS 5501 ELECTRICAL AND MECHANICAL SPECIFICATIONS
3.1 Absolute Maximum Ratings ..... 12
3.2 Recommended Operating Conditions ..... 12
LIST OF ILLUSTRATIONS
Figure 1 TMS 5501 Block Diagram ..... 2
Figure 2
Figure 3 Data Bus Assignments for TMS 550 । Status ..... 9
Figure 4 Discrete Command Format ..... 10
Figure 5 Data Bus Assignments for Rate Commands ..... 11
Figure 6 Read Cycle Timing ..... 14
Figure 7 Write Cycle Timing ..... 15
Figure 8 Sensor/Interrupt Timing ..... 15

Information contained in this publication is believed to be accurate and reliable. However, responsibility is assumed neither for its use nor for any infringement of patents or rights of others that may result from its use. No license is granted by implication or otherwise under any patent or patent right of Texas Instruments or others.

Copyright © 1975
Texas Instruments Incorporated

## TMS 5501 MULTIFUNCTION INPUT/OUTPUT CONTROLLER

## 1. INTRODUCTION

### 1.1 DESCRIPTION

The TMS 5501 is a multifunction input/output circuit for use with TI's TMS 8080 CPU. It is fabricated with the same N -channel silicon-gate process as the TMS 8080 and has compatible timing, signal levels, and power supply requirements. The TMS 5501 provides a TMS 8080 microprocessor system with an asynchronous communications interface, data I/O buffers, interrupt control logic, and interval timers.


FIGURE 1-TMS 5501 BLOCK DIAGRAM
The I/O section of the TMS 5501 contains an eight-bit parallel input port and a separate eight-bit parallel output port with storage register. Five programmable interval timers provide time intervals from $64 \mu$ sto 16.32 ms .

The interrupt system allows the processor to effectively communicate with the interval timers, external signals, and the communications interface by providing TMS 8080 -compatible interrupt logic with masking capability.

Data transfers between the TMS 5501 and the CPU are carried by the data bus and controlled by the interrupt, chip enable, sync, and address lines. The TMS 8080 uses four of its memory-address lines to select one of 14 commands to which the TMS 5501 will respond. These commands allow the CPU to:
.... read the receiver buffer
.... read the input port
.... read the interrupt address
.... read TMS 5501 status
.... issue discrete commands
.... Ioad baud rate register
.... load the transmitter buffer
.... load the output port
.... load the mask register
... load an interval timer

The commands are generated by executing memory referencing instructions such as MOV (register to memory) with the memory address being the TMS 5501 command. This provides a high degree of flexibility for I/O operations.by letting the systems programmer use a variety of instructions.

### 1.2 SUMMARY OF OPERATION

## Addressing the TMS 5501

A convenient method for addressing the TMS 5501 is to tie the chip enable input to the highest order address line of the CPU's 16 -bit address bus and the four TMS 5501 address inputs to the four lowest order bits of the bus. This, of course, limits the system to 32,768 words of memory but in many applications the full 65,536 word memory addressing capability of the TMS 8080 is not required.

## Communications Functions

The communications section of the TMS 5501 is an asynchronous transmitter and receiver for serial communications and provides the following functions:

Programmable baud rate - A CPU command selects a baud rate of $110,150,300,1200,2400,4800$, or 9600 baud.
Incoming character detection - The receiver detects the start and stop bits of an incoming character and places the character in the receive buffer.

Character transmission - The transmitter generates start and stop bits for a character received from the CPU and shifts it out.

Status and command signals - Via the data bus, the TMS 5501 signals the status of: framing error and overrun error flags; data in the receiver and transmitter buffers; start and data bit detectors; and end-of-transmission (break) signals from external equipment. It also issues break signals to external equipment.

## Data Interface

The TMS 5501 moves data between the CPU and external devices through its internal data bus, input port, and output port. When data is present on the bus that is to be sent to an external device, a Load Output Port (LOP) command from the CPU puts the data on the $\overline{X O}$ pins of the TMS 5501 by latching it in the output port. The data remains in the port until another LOP command is received. When the CPU requires data that is present on the External Input (XI) lines, it issues a command that gates the data onto the internal data bus of the TMS 5501 and consequently onto the CPU's data bus at the correct time during the CPU cycles.

## Interval Timers

To start a countdown by any of the five interval timers, the program selects the particular timer by an address to the TMS 5501 and loads the required interval into the timer via the data bus. Loading the timer activates it and it counts down in increments of 64 microseconds. The 8 -bit counters provide intervals that vary in duration from 64 to 16,320 microseconds. Much longer intervals can be generated by cascading the timers through software. When a timer reaches zero, it generates an interrupt that typically will be used to point to a subroutine that performs a servicing function such as polling a peripheral or scanning a keyboard. Loading an interval value of zero causes an immediate interrupt. A new value loaded while the interval timer is counting overrides the previous value and the interval timer starts counting down the new interval. When an interval timer reaches zero it remains inactive until a new interval is loaded.

## Servicing Interrupts

The TMS 5501 provides a TMS 8080 system with several interrupt control functions by receiving external interrupt signals, generating interrupt signals, masking out undersired interrupts, establishing the priority of interrupts, and generating RST instructions for the TMS 8080. An external interrupt is received on pin 22, SENS. An additional external interrupt can be received on pin 32, XI7, if selected by a discrete command from the TMS 8080 (See Figure 4). The TMS 5501 generates an interrupt when any of the five interval timers count to zero. Interrupts are also generated when the receiver buffer is loaded and when the transmitter buffer is empty.

When an interrupt signal is received by the interrupt register from a particular source, a corresponding bit is set and gated to the mask register. A pattern will have previously been set in the mask register by a load-mask-register command from the TMS 8080. This pattern determines which interrupts will pass through to the priority logic. The priority logic allows an interrupt to generate an RST instruction to the TMS 8080 only if there is no higher priority interrupt that has not been accepted by the TMS 8080. The TMS 5501 prioritizes interrupts in the order shown below:
1st - Interval Timer \#1
2nd - Interval Timer $=2$
3rd - External Sensor
4th - Interval Timer \#3
5th - Receiver Buffer Loaded
6th - Transmitter Buffer Emptied
7th - Interval Timer \#4
8th - Interval Timer \#5 or an External Input (XI 7)

The highest priority interrupt passes through to the interrupt address logic, which generates the RST instruction to be read by the TMS 8080. See Table 3 for relationship of interrupt sources to RST instructions and Figures 6 and 8 for timing relationships.

The TMS 5501 provides two methods of servicing interrupts; an interrupt-driven system or a polled-interrupt system. In an interrupt-driven system, the INT signal of the TMS 5501 is tied to the INT input of the TMS 8080. The sequence of events will be: (1) The TMS 5501 receives (or generates) an interrupt signal and readies the appropriate RST instruction. (2) The TMS 5501 INT output, tied to the TMS 8080 INT input, goes high signaling the TMS 8080 that an interrupt has occured. (3) If the TMS 8080 is enabled to accept interrupts, it sets the INTA (interrupt acknowledge) status bit high at SYNC time of the next machine cycle. (4) If the TMS 5501 has previously received an interrupt-acknowledge-enable command from the CPU (see Bit 3, Paragraph 2.2.5), the RST instruction is transferred to the data bus.

In a polled-interrupt system, INT is not used and the sequence of events will be: (1) The TMS 5501 receives (or generates) an interrupt and readies the RST instruction. (2) The TMS 5501 interrupt-pending status bit (see Bit 5, Paragraph 2.2.4) is set high (the interrupt-pending status bit and the INT output go high simultaneously). (3) At the prescribed time, the TMS 8080 polls the TMS 5501 to see if an interrupt has occurred by issuing a readTMS 5501-status command and reading the interrupt-pending bit. (4) If the bit is high, the TMS 8080 will then issue a read-interrupt-address command, which causes the TMS 5501 to transfer the RST instruction to the data bus as data for the instruction being executed by the TMS 8080.

### 1.3 APPLICATIONS

## Communications Terminals

The functions of the TMS 5501 make it particularly useful in TMS 8080-based communications terminals and generally applicable in systems requiring periodic or random servicing of interrupts, generation of control signals to external devices, buffering of data, and transmission and reception of asynchronous serial data. As an example, a system configuration such as shown in Figure 2 can function as the controller for a terminal that governs employee entrance into a plant or security areas within a plant. Each terminal is identified by a central computer through ID switches. The central system supplies each terminal's RAM with up to 16 employee access categories applicable to that terminal. These categories are compared with an employee's badge character when he inserts his badge into the badge sensor. If a
$\square$
match is not found, a reject light will be activated. If a match is found, the terminal will transmit the employee's badge number and access category to the central system, and a door unlock solenoid will be activated for 4 seconds. The central computer then may take the transmitted information and record it along with time and date of access.

The TMS 4700 is a $1024 \times 8$ ROM that contains the system program, and the TMS 4036 is a $64 \times 8$ RAM that serves as the stack for the TMS 8080 and storage for the access category information. TTL circuits control chip-enable information carried by the address bus. Signals from the CPU gate the address bits from the ROM, the RAM, or the TMS 5501 onto the data bus at the correct time in the CPU cycle. The clock generator consists of four TTL circuits along with a crystal, needed to maintain accurate serial data assembly and disassembly with the central computer.

The TMS 5501 handles the asynchronous serial communication between the TMS 8080 and the central system and gates data from the badge reader onto the data bus. It also gates control and status data from the TMS 8080 to the door lock and badge reader and controls the time that the door lock remains open. The TMS 5501 signals the TMS 8080 when the badge reader or the communication lines need service. The functions that the TMS 5501 is to perform are selected by an address from the TMS 8080 with the highest order address line tied to the TMS 5501 chip enable input and the four lowest order lines tied to the address inputs.

## 2. OPERATIONAL AND FUNCTIONAL DESCRIPTION

This detailed description of the TMS 5501 consists of:
INTERFACE SIGNALS - a definition of each of the circuit's external connections
COMMANDS - the address required to select each of the TMS 5501 commands and a description of the response to the command.

### 2.1 INTERFACE SIGNALS

The TMS 5501 communicates with the TMS 8080 via four address lines: a chip enable line, an eight-bit bidirectional data bus, an interrupt line, and a sync line. It communicates with system components other than the CPU via eight external inputs, eight external outputs, a serial receiver input, a serial transmitter output, and an external sensor input. Table 1 defines the TMS 5501 pin assignments and describes the function of each pin.

| TABLE 1 |  |  |
| :---: | :---: | :---: |
| TMS 5501 PIN ASSIGNMENTS AND FUNCTIONS |  |  |
| SIGNATURE | PIN | DESCRIPTION INPUTS |
| CE | 18 | Chip enable-When CE is low, the TMS 5501 address decoding is inhibited, which prevents execution of any of the TMS 5501 commands. |
| A3 | 17 | Address bus-A3 through A0 are the lines that are addressed by the TMS 8080 to select a particular |
| A2 | 16 | TMS 5501 function. |
| A1 | 15 |  |
| A0 | 14 |  |
| SYNC | 19 | Synchronizing signal-TheSYNC signal is issued by the TMS 8080 and indicates the beginning of a machine cycle and availability of machine status. When the SYNC signal is active (high), the TMS 5501 will monitor the data bus bits DO (interrupt acknowledge) and D1 ( $\overline{\mathrm{WO}}$, data output function). |
| $\overline{\mathrm{RCV}}$ | 5 | Receiver serial data input line- $\overline{\mathrm{RCV}}$ must be held in the inactive (high) state when not receiving data. A transition from high to low will activate the receive circuitry. |

## TABLE 1 (continued) TMS 5501 PIN ASSIGNMENTS AND FUNCTIONS

SIGNATURE PIN

XI $0 \quad 39$
XI 138
XI 237
XI 336
XI 435
XI 5
XI 633
XI7 72
SENS 22
$\overline{\mathrm{XO}} 0 \quad 24$
XO $1 \quad 25$
XO 26
$\overline{\mathrm{XO}} 3 \quad 27$
$\overline{\mathrm{XO}} 4 \quad 28$
$\overline{\mathrm{XO}} 5 \quad 29$
प्र 630
XO 731
XMT 40

D0 13
D1 12
D2 11
D3 10
D4 9
D5 8
D6 7
D7 6
INT 23

| $V_{S S}$ | 4 |
| :---: | :---: |
| $V_{\text {BB }}$ | 1 |
| $V_{C C}$ | 2 |
| $V_{D D}$ | 3 |
| $\phi 1$ | 20 |
| $\phi 2$ | 21 |

register, which, if enabled, generates an interrupt to the TMS 8080.

## OUTPUTS

External outputs-These eight external outputs are driven by the complement of the output
DESCRIPTION
INPUTS
External inputs-These eight external inputs are gated to the data bus when the read-external-inputs function is addressed. External input n is gated to data bus bit n without conversion.

External interrupt sensing - A transition from low to high at SENS sets a bit in the interrupt register; i.e., if output register bit n is loaded with a high (low) from data bus bit n by a loadoutput register command, the external output $n$ will be a low (high). The external outputs change only when a load-output-register function is addressed.

Transmitter serial data output line-This line remains high when the TMS 5501 is not transmitting.

DATA BUS INPUT/OUTPUT
Data bus - Data transfers between the TMS 5501 and the TMS 8080 are made via the 8 -bit bidirectional data bus. D0 is the LSB. D7 is the MSB.

Interrupt-When active (high), the INT output indicates that at least one of the interrupt conditions has occurred and that its corresponding mask-register bit is set.

POWER AND CLOCKS
Ground reference
Supply voltage ( -5 V nominal)
Supply voltage ( 5 V nominal)
Supply voltage ( 12 V nominal)
Phase 1 clock
Phase 2 clock

The TMS 5501 operates as memory device for the TMS 8080. Functions are initiated via the TMS 8080 address bus and the TMS 5501 address inputs. Address decoding to determine the command function being issued is defined in Table 2.

TABLE 2
COMMAND ADDRESS DECODING When Chip Enable Is High

| A3 | A2 | A1 | AO | COMMAND | FUNCTION | PARAGRAPH |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| L | L | L | L | Read receiver buffer | $\mathrm{RBn} \rightarrow \mathrm{Dn}$ | 2.2.1 |
| L | L | L | H | Read external inputs | $\mathrm{XIn} \rightarrow \mathrm{Dn}$ | 2.2.2 |
| L | L | H | L | Read interrupt address | RST $\rightarrow$ Dn | 2.2.3 |
| L | L | H | H | Read TMS 5501 status | (Status) $\rightarrow$ Dn | 2.2.4 |
| L | H | L | L | Issue discrete commands | See Figure 4 | 2.2.5 |
| L | H | L | H | Load rate register | See Figure 4 | 2.2.6 |
| L | H | H | L | Load transmitter buffer | $\mathrm{Dn} \rightarrow \mathrm{TBn}$ | 2.2.7 |
| L | H | H | H | Load output port | $\mathrm{Dn} \rightarrow \overline{\mathrm{XO}} \mathrm{n}$ | 2.2.8 |
| H | L | L | L | Load mask register | $\mathrm{Dn} \rightarrow \mathrm{MRn}$ | 2.2.9 |
| H | L | L | H | Load interval timer 1 | Dn $\rightarrow$ Timer 1 | 2.2.10 |
| H | L | H | L | Load interval timer? | Dn $\rightarrow$ Timer 2 | 2.2.10 |
| H | L | H | H | Load interval timer 3 | Dn $\rightarrow$ Timer 3 | 2.2.10 |
| H | H | L | L | Load interval timer 4 | Dn $\rightarrow$ Timer 4 | 2.2.10 |
| H | H | L | H | Load interval timer 5 | Dn $\rightarrow$ Timer 5 | 2.2.10 |
| H | H | H | L | No function |  |  |
| H | H | H | H | No function |  |  |

RBn Receiver buffer bit $n$
Dn Data bus I/O terminal $n$
XIn External input terminal $n$
RST $11\left(\mid A_{2}\right)\left(\mid A_{1}\right)\left(\mid A_{0}\right) 111$ (see Table 3)
TBn Transmit buffer bit $n$
XOn Output register bit $n$
MRn Mask register bitn

TABLE 3 RST INSTRUCTIONS
DATA BUS BIT

| INTERRUPT CAUSED BY |  |  |  |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{2}$ | 3 | 4 | 5 | 6 | $\mathbf{7}$ |  |
| H | H | H | L | L | L | H | H | Interval Timer 1 |
| H | H | H | H | L | L | H | H | Interval Timer 2 |
| H | H | H | L | H | L | H | H | External Sensor |
| H | H | H | H | H | L | H | H | Interval Timer 3 |
| H | H | H | L | L | H | H | H | Receiver Buffer |
| H | H | H | H | L | H | H | H | Transmitter Buffer |
| H | H | H | L | H | H | H | H | Interval Timer 4 |
| H | H | H | H | H | H | H | H | Interval Timer 5 or X17 |

The following paragraphs define the functions of the TMS 5501 commands.

### 2.2.1 Read receiver buffer

Addressing the read-receiver-buffer function causes the receiver buffer contents to be transferred to the TMS 8080 and clears the receiver-buffer-loaded flag.

### 2.2.2 Read external input lines

Addressing the read-external-inputs function transfers the states of the eight external input lines to the TMS 8080.

### 2.2.3 Read interrupt address

Addressing the read interrupt address function transfers the current highest priority interrupt address onto the data bus as read data. After the read operation is completed, the corresponding bit in the interrupt register is reset.

If the read-interrupt-address function is addressed when there is no interrupt pending, a false interrupt address will be read. TMS 5501 status function should be addressed in order to determine whether or not an interrupt condition is pending.

### 2.2.4 Read TMS 5501 status

Addressing the read-TMS 5501-status function gates the various status conditions of the TMS 5501 onto the data bus. The status conditions, available as indicated in Figure 3, are described in the following paragraphs.

| BIT: | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | START | FULL | INTRPT | XMIT | RCV | SERIAL | OVERRUN | FRAME |
|  | BIT | BIT | PENDING | BUFFER | BUFFER | RCVD | ERROR | ERROR |
|  | DETECT | DETECT |  | EMPTY | LOADED |  |  |  |

FIGURE 3-DATA BUS ASSIGNMENTS FOR TMS 5501 STATUS

## Bit 0, framing error

A high in bit 0 indicates that a framing error was detected on the last character received (either one or both stop bits were in error). The framing error flag is updated at the end of each character. Bit 0 of the TMS 5501 status will remain high until the next valid character is received.

## Bit 1, overrun error

A high in bit 1 indicates that a new character was loaded into the receiver buffer before a previous character was read out. The overrun error flag is cleared each time the read-I/O-status function is addressed or a reset command is issued.

## Bit 2, serial received data

Bit 2 monitors the receiver serial data input line. This line is provided as a status input for use in detecting a break and for test purposes. Bit 2 is normally high when no data is being received.

## Bit 3, receiver buffer loaded

A high in bit 3 indiciates that the receiver buffer is loaded with a new character. The receiver-buffer-loaded flag remains high until the read-receiver-buffer function is addressed (at which time the flag is cleared). The reset function also clears this flag.

## Bit 4, transmitter butfer empty

A high in bit 4 indicates that the transmitter buffer register is empty and ready to accept a character. Note, however, that the serial transmitter register may be in the process of shifting out a character. The reset function sets the transmitter-buffer-empty flag high.

## Bit 5, interrupt pending

A high in bit 5 indicates that one or more of the interrupt conditions has occured and the corresponding interrupt is enabled. This bit is the status of the interrupt signal INT.

## Bit 6, full bit detected

A high in bit 6 indicates that the first data bit of a receive-data character has been detected. This bit remains high until the entire charac:ter has been received or until a reset is issued and is provided for test purposes.

Bit 7, start bit detected
A high in bit 7 indicates that the start bit of an incoming data character has been detected. This bit remains high until the entire character has been received or until a reset is issued and is provided for test purposes.

### 2.2.5 Issue discrete commands

Addressing the discrete command function causes the TMS 5501 to interpret the data bus information according to the following descriptions. See Figure 4 for the discrete command format. Bits 1 through 5 are latched until a different discrete command is received.


## FIGURE 4-DISCRETE COMMAND FORMAT

## Bit 0, reset

A high in bit 0 will cause the following:

1) The receiver buffer and register are cleared to the search mode including the receiver-buffer-loaded flag, the starl-bit-detected flag, the full-bit-detected flag, and the overrun-error flag. The receiver buffer is not cleared and will contain the last character received.
2) The transmitter data output is set high (marking). The transmitter-buffer-empty flag is set high indicating that the transmitter buffer is ready to accept a character from the TMS 8080.
3) The interrupt register is cleared except for the bit corresponding to the transmitter buffer interrupt, which is set high.
4) The interval timers are inhibited.

A low in bit 0 causes no action. The reset function has no affect on the output port, the external inputs, interrupt acknowledge enable, the mask register, the rate register, the transmitter register, or the transmitter buffer.

## Bit 1, break

A low in bit 1 causes the transmitter data output to be reset low (spacing).

If bit 0 and bit 1 are both high, the reset function will override.

## Bit 2, interrupt 7 select

Interrupt 7 may be generated either by a low to high transition of external input 7 or by interval timer 5 .

A high in bit 2 selects the interrupt 7 source to be the transition of external input 7 . A low in bit 2 selects the interrupt 7 source to be interval timer 5 .

## Bit 3, interrupt acknowledge enable

The TMS 5501 decodes data bus (CPU status) bit 0 at SYNC of each machine cycle to determine if an interrupt acknowledge is being issued.

A high in bit 3 enables the TMS 5501 to accept the interrupt acknowledge decode. A low in bit 3 causes the TMS 5501 to ignore the interrupt acknowledge decode.

Bit 4 and bit 5 are used only during testing of the TMS 5501. For correct system operation both bits must be kept low.

Bit 6 and bit 7 are not used and can assume any value.

### 2.2.6 Load rate register

Addressing the load-rate-register function causes the TMS 5501 to load the rate register from the data bus and interpret the data bits (See Figure 5) as follows.


FIGURE 5-DATA BUS ASSIGNMENTS FOR RATE COMMANDS

## Bits 0 through 6, rate select

The rate select bits (bits 0 through 6 ) are mutually exclusive, i.e., only one bit may be high. A high in bits 0 through 6 will select the baud rate for both the transmitter and receiver circuitry as defined below and in Figure 5:

| Bit 0 | 110 baud |
| :--- | ---: |
| Bit 1 | 150 baud |
| Bit 2 | 300 baud |
| Bit 3 | 1200 baud |
| Bit 4 | 2400 baud |
| Bit 5 | 4800 baud |
| Bit 6 | 9600 baud |

If more than one bit is high, the highest rate indicated will result. If bits 0 through 6 are all low, both the receiver and the transmitter circuitry will be inhibited.

## Bit 7, stop bits

Bit 7 determines whether one or two stop bits are to be used by both the transmitter and receiver circuitry. A high in bit 7 selects one stop bit. A low in bit 7 selects two stop bits.

### 2.2.7 Load transmitter buffer

Addressing the load-transmitter-buffer function transfers the state of the data bus into the transmitter buffer.

### 2.2.8 Load output port

Addressing the load-output-port function transfers the state of the data bus into the output port. The data is latched and remains on $\overline{X O O}$ through $\overline{X O 7}$ as the complement of the data bus until new data is loaded.

### 2.2.9 Load mask register

Addressing the load-mask-register function loads the contents of the data bus into the mask register. A high in data bus bit n enables interrupt n . A low inhibits the corresponding interrupt.

### 2.2.10 Load timer n

Addressing the load-timer-n function loads the contents of the data bus into the appropriate interval timer. Time intervals of from $64 \mu \mathrm{~s}$ (data bus $=$ LLLLLLLH) to $16,320 \mu \mathrm{~s}$ (data bus HHHHHHHH ) are counted in $64-\mu \mathrm{s}$, steps. When the count of interval timer n reaches 0 , the bit in the interrupt register that corresponds to timer n is set and an interrupt is generated. Loading a! lows causes an interrupt immediately.

## 3. TMS 5501 ELECTRICAL AND MECHANICAL SPECIFICATIONS

### 3.1 ABSOLUTE MAXIMUM RATINGS OVER OPERATING FREE-AIR TEMPERATURE RANGE (UNLESS OTHERWISE NOTED)*



[^2]
### 3.2 RECOMMENDED OPERATING CONDITIONS

|  | MIN | NOM | MAX | UNIT |
| :---: | :---: | :---: | :---: | :---: |
| Supply voltage, $\mathrm{V}_{\text {BB }}$ | -4.75 | -5 | -5.25 | $\checkmark$ |
| Supply voltage, $\mathrm{V}_{\mathrm{CC}}$ | 4.75 | 5 | 5.25 | V |
| Supply voltage, $\mathrm{V}_{\text {DD }}$ | 11.4 | 12 | 12.6 | V |
| Supply voltage, $\mathrm{V}_{\text {SS }}$ |  | 0 |  | V |
| High-level input voltage, $\mathrm{V}_{\text {IH }}$ (all inputs except clocks) | 3.3 |  | $\mathrm{V}_{\mathrm{CC}}+1$ | V |
| High-level clock input voltage, $\mathrm{V}_{\mathrm{IH}}(\phi)$ | $V_{D D^{-1}}$ |  | $V_{\text {DD }}{ }^{+1}$ | V |
| Low-level input voltage, $\mathrm{V}_{\text {IL }}$ (all inputs except clocks) (see Note 2) | -1 |  | 0.8 | V |
| Low-level clock input voltage, $\mathrm{V}_{\text {IL }}(\boldsymbol{\phi})$ (see Note 2 ) | -1 |  | 0.6 | V |
| Operating free-air temperature, ${ }^{\text {/ } A}$ | 0 |  | 70 | ${ }^{\circ} \mathrm{C}$ |

NOTE 2: The algebraic convention where the most negative limit is designated as minimum is used in this specification for logic voltage levels only.
3.5 SWITCHING CHARACTERISTICS OVER FULL RANGE OF RECOMMENDED OPERATING CONDITIONS (SEE FIGURES 6 AND 7)

| PARAMETER |  | TESTCONDITIONS | MIN MAX | UNIT |
| :---: | :---: | :---: | :---: | :---: |
| tpZ | Data bus output enable time | $\begin{aligned} & C_{L}=100 \mathrm{pF}, \\ & R_{L}=1.3 \mathrm{ks} 2 \end{aligned}$ | 200 | ns |
| tPXZ | Data bus output disable time to high-impedance state |  | 180 | ns |
| tPD | External data output propagation delay tıme from si |  | 200 | ns |


$C_{L}$ inclucles probe and jig capacitance
LOAD CIRCUIT


NOTE: For, 1 or 2 inputs, high and low timing points are $90 \%$ and $10 \%$ of $V_{1 H(c)}$. Allother timing points are the $50 \%$ level.
FIGURE 6-READ CYCLE TIMING


NOTE: For $\phi 1$ and $\phi 2$ inputs, high and low timing points are $90 \%$ and $10 \%$ of $V_{1} H(\phi)$. All other timing points are the $50 \%$ level.

FIGURE 7-WRITE CYCLE TIMING


NOTES: 1. The RST instruction occurs during the output data valid time of ine read cycle.
2. All timing points are $50 \%$ of $V_{\text {IH }}$.

FIGURE 8-SENSOR/INTERRUPT TIMING

## Appendix D

## TMS 8080 Microprocessor

## TABLE OF CONTENTS

1. ARCHITECTURE
1.1 Introduction ..... 2
1.2 The Stack ..... 2
1.3 Registers ..... 2
1.4 The Arithmetic Unit ..... 3
1.5 Status and Control ..... 3
1.6 I/O Operations ..... 3
1.7 Instruction Timing ..... 3
2. TMS 8080 INSTRUCTION SET
2.1 Instruction Formats ..... 6
2.2 Instruction Set Description ..... 7
2.2.1 Instruction Symbols ..... 7
2.2.2 Accumulator Group Instructions ..... 8
2.2.3 Input/Output Instructions ..... 9
2.2.4 Machine Instructions ..... 9
2.2.5 Program Counter and Stack Control Instructions ..... 10
2.2.6 Register Group Instructions ..... 11
2.3 Instruction Set Opcodes Alphabetically Listed ..... 12
3. TMS 8080 ELECTRICAL AND MECHANICAL SPECIFICATIONS
3.1 Absolute Maximum Ratings ..... 17
3.2 Recommended Operating Conditions ..... 17
3.3 Electrical Characteristics ..... 17
3.4 Timing Requirements ..... 18
3.5 Switching Characteristics ..... 18
3.6 Terminal Assignments ..... 20
3.7 Mechanical Data ..... 20
LIST OF ILLUSTRATIONS
Figure 1 TMS 8080 Functional Block Diagram ..... 2
Figure 2 Voltage Waveforms ..... 19

## TMS 8080 MICROPROCESSOR

## 1. ARCHITECTURE

### 1.1 INTRODUCTION

The TMS 8080 is an 8 -bit parallel central processing unit (CPU) fabricated on a single chip using a high-speed N -channel silicon-gate process. (See Figure 1). A complete microcomputer system with a $2 \cdot \mu \mathrm{~s}$ instruction cycle can be formed by interfacing this circuit with any appropriate memory. Separate 8 -bit data and 16 -bit address buses simplify the interface and allow direct addressing of 65,536 bytes of memory. Up to 256 input and 256 output ports are also provided with direct addressing. Control signals are brought directly out of the processor and all signals, excluding clocks, are TTL compatible.
1.2 THE STACK

The TMS 8080 incorporates a stack architecture in which a portion of external memory is used as a pushdown stack for storing data from working registers and internal machine status. A 16 -bit stack pointer (SP) is provided to facilitate stack location in the memory and to allow almost unlimited interrupt handling capability. The CALL and RST (restart) instructions use the SP to store the program counter (PC) into the stack. The RET (return) instruction uses the SP to acquire the previous PC value. Additional instructions allow data from registers and flags to be saved in the stack.

### 1.3 REGISTERS

The TMS 8080 has three categories of registers: general registers, program control registers, and internal registers. The general registers and program control registers are listed in Table 1. The internal registers are not accessible by the programmer. They include the instruction register, which holds the present instruction, and several temporary storage registers to hold internal data or latch input output addresses and data.


### 1.4 THE ARITHMETIC UNIT

Arithmetic operations are performed in an 8 -bit parallel arithmetic unit that has both binary and decimal capabilities. Four testable internal flag bits are provided to facilitate program control, and a fifth flag is used for decimal corrections. Table 2 defines these flags and their operation. Decimal corrections are performed with the DAA instruction. The DAA corrects the result of binary arithmetic operation on BCD data as shown in Table 3.

### 1.5 STATUS AND CONTROL

Two types of status are provided by the TMS8080. Certain status is indicated by dedicated control lines. Additional status is transmitted on the data bus during the beginning of each instruction cycle (machine cycle). Table 4 indicates the pin functions of the TMS8080. Table 5 defines the status information that is presented during the beginnirig of each machine cycle (SYNC time) on the data bus.

### 1.6 I/O OPERATIONS

Input/output operations (I/O) are performed using the IN and OUT instructions. The second byte of these instructions indicates the device address ( 256 device addresses). When an $I N$ instruction is executed, the input device address appears in duplicate on $A 7$ through $A 0$ and $A 15$ through $A 8$, along with $\overline{W O}$ and INP status on the data bus. The addressed input device then puts its input data on the data bus for entry into the accumulator. When an OUT instruction is executed, the same operation occurs except that the data bus has OUT status and then has output data.

Direct memory access channels (DMA) can be OR-tied directly with the data and address buses through the use of the HOLD and HLDA (hold acknowledge) controls. When a HOLD request is accepted by the CPU, HLDA goes high, the address and data lines are forced to a high-impedance or "floating" condition, and the CPU stops until the HOLD request is removed.

Interfacing with different speed memories is easily accomplished by use of the WAIT and READY pins. During each machine cycle, the CPU polls the READY input and enters a wait condition until the READY line becomes true. When the WAIT output pin is high, it indicates that the CPU has entered the wait state.

Designing interrupt driven systems is simplified through the use of vectored interrupts. At the end of each instruction, the CPU polls the INT input to determine if an interrupt request is being made. This action does not occur if the CPU is in the HOLD state or if interrupts are disabled. The INTE output indicates if the interrupt logic is enabled (INTE is high). When a request is honored, the INTA status bit becomes high, and an RST instruction may be inserted to force the CPU to jump to one of eight possible locations. Enabling or disabling interrupts is controlled by special instructions (EI or DI). The interrupt input is automatically disabled when an interrupt request is accepted or when a RESET signal is received.

### 1.7 INSTRUCTION TIMING

The execution time of the instructions varies depending on the operation required and the number of memory references needed. A machine cycle is defined to be a memory referencing operation and is either 3,4 , or 5 state times long. A state time (designated S ) is a full cycle of clocks $\phi 1$ and $\phi 2$. (NOTE: The exception to this rule is the DAD instruction, which consists of 1 memory reference in 10 state times). The first machine cycle (designated M1) is either 4 or 5 state times long and is the "instruction fetch" cycle with the program counter appearing on the address bus. The CPU then continues with as many $M$ cycles as necessary to complete the execution of the instruction (up to a maximum of 5). Thus the instruction execution time varies from 4 state times (several including ADDr) to 18 (XTHL). The WAIT or HOLD conditions may affect the execution time since they can be used to control the machine (for example to "single step") and the HALT instruction forces the CPU to stop until an interrupt is received. As the instruction execution is completed (or in the HAIT state) the INT pin is polled for an interrupt. In the event of an interrupt, the PC will not be incremented during the next M 1 and an RST instruction can be inserted.

TABLE 1
TMS 8080 REGISTERS

| NAME | DESIGNATOR | LENGTH | PURPOSE |
| :---: | :---: | :---: | :---: |
| Accumulator | A | 8 | Used for arithmetic, logical, and I/O operations |
| B Register | B | 8 | General or most significant 8 bits of double register BC |
| C Register | C | 8 | General or least significant 8 bits of double register BC |
| D Register | D | 8 | General or most significant 8 bits of double register DE |
| E Register | E | 8 | General or least significant 8 bits of double register DE |
| H Register | H | 8 | General or most significant 8 bits of double register HL |
| L Register | L | 8 | General or least significant 8 bits of double register HL |
| Program Counter | PC | 16 | Contains address of next byte to be fetched |
| Stack Pointer | SP | 16 | Contains address of the last byte of data saved in the memory stack |
| Flag Register | F | 5 | Five flags ( $C, ~ Z, ~ S, ~ P, ~ C 1) ~$ |

NOTE: Registers $B$ and $C$ may be used together as a single 16 bit register, likewise, $D$ and $E$, and $H$ and $L$.
TABLE 2
FLAG DESCRIPTIONS

| SYMBOL | TESTABLE | DESCRIPTION |
| :---: | :---: | :---: |
| C | YES | C is the carry/borrow out of the MSB (most significant bit) of the ALU (Arithment Logic Unit). A TRUE condition ( $C=1$ ) indicates overflow for addition or underflow for subtraction. |
| $z$ | YES | A TRUE condition ( $Z=1$ ) indicates that the output of the ALU is equal to zero. |
| S | YES | A T.RUE condition ( $S=1$ ) indicates that the MSB of the ALU output is equal to a one (1). |
| P | YES | A TRUE condition ( $\mathrm{P}=1$ ) indicates that the output of the $A L U$ has even parity (the number of bits equal to one is even). |
| C1 | NO | C1 is the carry out of the fourth bit of the ALU (TRUE condition). C1 is used only for BCD correction with the DAA instruction. |

TABLE 3
FUNCTION OF THE DAA INSTRUCTION
Assume the accumulator ( $A$ ) contains two BCD digits, $X$ and $V$

|  | 4 |  |
| :--- | :--- | :--- | |  | 3 |
| :---: | :---: |


| ACCUMULATOR |  |  |  | ACCUMULATOR |  |  |  |
| :---: | :--- | :---: | :---: | :---: | :---: | :---: | :---: |
| BEFORE DAA |  |  |  |  |  |  |  |$|$

NOTE: The corrections shown in Table 3 are sufficient for addition. For subtraction, the programmer must account for the borrow condition that can occur and give erroneous results. The most straight forward method is to set $A=9916$ and carry $=1$. Then ads the minuend to $A$ after subtracting the subtrahend from $A$.

TABLE 4
TMS 8080 PIN DEFINITIONS

| SIGNATURE | PIN | $1 / 0$ | DESCRIPTION |
| :---: | :---: | :---: | :---: |
| A15 (MSB) | 36 | OUT | A15 through A0 comprise the address bus. True memory or I/O device addresses appear on |
| A14 | 39 | OUT | this 3 -state bus during the first state time of each instruction cycle. |
| A13 | 38 | OUT |  |
| A12 | 37 | OUT |  |
| A11 | 40 | OUT |  |
| A 10 | 1 | OUT |  |
| A9 | 35 | OUT |  |
| A8 | 34 | OUT |  |
| A7 | 33 | OUT |  |
| A6 | 32 | OUT |  |
| A5 | 31 | OUT |  |
| A4 | 30 | OUT |  |
| A3 | 29 | OUT |  |
| A2 | 27 | OUT |  |
| A1 | 26 | OUT |  |
| AO (LSB) | 25 | OUT |  |
| D7 (MSB) | 6 | IN/OUT | D7 through D0 comprise the bidirectional 3-state data bus. Memory, status, or 1/O data is |
| D6 | 5 | IN/OUT | transferred on this bus. |
| D5 | 4 | IN/OUT |  |
| D4 | 3 | IN/OUT |  |
| D3 | 7 | IN/OUT |  |
| D2 | 8 | IN/OUT |  |
| D1 | 9 | IN/OUT |  |
| D0 (LSB) | 10 | IN/OUT |  |
| $\mathrm{V}_{\text {SS }}$ | 2 |  | Ground reference |
| $V_{\text {BB }}$ | 11 |  | Supply voltage ( -5 V nominal) |
| $\mathrm{V}_{\mathrm{CC}}$ | 20 |  | Supply voltage ( 5 V nominal) |
| $V_{\text {DD }}$ | 28 |  | Supply voltage ( 12 V nominal) |
| $\phi 1$ | 22 | IN | Phase 1 clock. |
| ¢2 | 15 | IN | Phase 2 clock. See page 19 for $\phi 1$ and $\phi 2$ timing. |
| RESET | 12 | IN | Reset. When active (high) for a minimum of 3 clock cycles, the RESET input causes the TMS 8080 to be reset. PC is cleared, interrupts are disabled, and after RESET, instruction execution starts at memory location 0 . To prevent a lockup condition, a HALT instruction must not be used in location 0 . |
| HOLD | 13 | IN | Hold signal. When active (high) HOLD causes the TMS 8080 to enter a hold state and float (put the 3 -state address and data bus in a high-impedance state). The chip acknowledges entering the hold state with the HLDA signal and will not accept interrupts until it leaves the hold state. |
| INT | 14 | IN | Interrupt request. When active (high) INT indicates to the TMS 8080 that an interrupt is being requested. The TMS 8080 polls INT during a HALT or at the end of an instruction. The request will be accepted except when INTE is low or the CPU is in the HOLD condition. |
| INTE | 16 | OUT | Interrupts enabled. INTE indicates that an interrupt will be accepted by the TMS 8080 unless it is in the hold state. INTE is set to a high logic level by the EI (Enable Interrupt) instruction and reset to a low logic level by the DI (Disable Interrupt) instruction. INTE is also reset when an interrupt is accepted and by a high on RESET. |
| DBIN | 17 | OUT | Data bus in. DBIN indicates whether the data bus is in an input or an output mode. (high = input, low $=$ output). |

TABLE 4 (CONTINUED)

| SIGNATURE | PIN | I/O | DESCRIPTION |
| :--- | :---: | :---: | :---: |
| $\overline{\text { WR }}$ | 18 | OUT | Write. When active (low) WR indicates a write operation on the data bus to memory or to an <br> I/O port. |
| SYNC | 19 | OUT | Synchronizing control line. When active (high) SYNC indicates the beginning of each <br> machine cycle of the TMS $8080 . ~ S t a t u s ~ i n f o r m a t i o n ~ i s ~ a l s o ~ p r e s e n t ~ o n ~ t h e ~ d a t a ~ b u s ~ d u r i n g ~$ |
| SYNC for external latches. |  |  |  |

TABLE 5
TMS 8080 STATUS

| SIGNATURE | DATA BUS BIT | DESCRIPTION |
| :---: | :---: | :---: |
| INTA | DO | Interrupt acknowledge. |
| $\overline{W O}$ | D1 | Indicates that current machine cycle will be a read (input) (high = read) or a write (output) (low = write) operation. |
| STACK | D2 | Indicates that address is stack address from the SP. |
| HLTA | D3 | HALT instruction acknowledge. |
| OUT | D4 | Indicates that the address bus has an output device address and the data bus has output data. |
| M1 | D5 | Indicates instruction acquisition for fifst byte. |
| INP | D6 | Indicates address bus has address of input device. |
| MEMR | D7 | Indicates that data bus will be used for memory read data. |

2. TMS 8080 INSTRUCTION SET
2.1 INSTRUCTION FORMATS

TMS 8080 instructions are either one, two, or three bytes long and are stored as binary integers in successive memory locations in the format shown below.

One-Byte Instructions D7 D6 D5 D4 D3 D2 D1 D0

OP CODE
Two-Byte Instructions
D7 D6 D5 D4 D3 D2 D1 D0
D7 D8 D5 D4 D3 D2 D1 D0
OP CODE

OPERAND
Three-Byte Instructions
D7 D6 D5 D4 D3 D2 D1 D0

D7 D6 D5 D4 D3 D2 D1 D0
D7 D6 D5 D4 D3 D2 D1 D0
OP CODE

LOW ADDRESS OR OPERAND 1
HIGH ADDRESS OR OPERAND 2

### 2.2 INSTRUCTION SET DESCRIPTION

Operations resulting from the execution of TMS 8080 instructions are described in this section. The flags that are affected by each instruction are given after the description.

### 2.2.1 INSTRUCTION SYMBOLS



| MNEMONIC | OPERANDS | BYTES | M CYCLES/ STATES |
| :---: | :---: | :---: | :---: |
| ACl | $\mathrm{b}_{2}$ | 2 | $2 / 7$ |
| ADC | M | 1 | 2/7 |
| ADC | ${ }^{\text {a }}$ | 1 | 1/4 |
| ADD | M | 1 | 2/7 |
| ADD | ${ }^{\text {a }}$ | 1 | 1/4 |
| ADI | $\mathrm{b}_{2}$ | 2 | $2 / 7$ |
| ANA | M | 1 | $2 / 7$ |
| ANA | ${ }^{\text {ra }}$ | 1 | 1/4 |
| ANI | $\mathrm{b}_{2}$ | 2 | $2 / 7$ |
| CMA |  | 1 | 1/4 |
| CMC |  | 1 | 1/4 |
| CMP | M | 1 | 2/7 |


| CMP | $r_{\text {a }}$ | 1 | $1 / 4$ |
| :---: | :---: | :---: | :---: |
| CPI | $b_{2}$ | 2 | $2 / 7$ |
| DAA |  | 1 | $1 / 4$ |


| DAD | rb | $\uparrow$ | $1 / 10$ |
| :--- | :---: | :---: | :---: |
| LDA | $b_{3} b_{2}$ | 3 | $4 / 13$ |
| LDAX | $r_{c}$ | 1 | $2 / 7$ |
| ORA | $M$ | 1 | $2 / 7$ |
|  |  |  |  |
| ORA | $r_{a}$ | 1 | $1 / 4$ |
| ORI | $b_{2}$ | 2 | $2 / 7$ |
| RAL |  | 1 | $1 / 4$ |
| RAR |  | 1 | $1 / 4$ |
| RLC |  | 1 | $1 / 4$ |
|  |  | 1 | $1 / 4$ |


| MNEMONIC | OPERANDS | BYTES | M CYCLES STATES |
| :---: | :---: | :---: | :---: |
| SBB | M | 1 | 2/7 |
| SBB | $\mathrm{r}_{\mathrm{a}}$ | 1 | 1/4 |
| SBI | $\mathrm{b}_{2}$ | 2 | 2/7 |
| STA | $b_{3} b_{2}$ | 3 | 4/13 |
| STAX | $\mathrm{r}_{\mathrm{c}}$ | 1 | 2/7 |
| STC |  | 1 | 1/4 |
| SUB | M | 1 | 2/7 |
| SUB | $\mathrm{r}_{\mathrm{a}}$ | 1 | 1/4 |
| SUI | $\mathrm{b}_{2}$ | 2 | 2/7 |
| XRA | M | 1 | 2/7 |
| XRA | ra | 1 | 1/4 |
| XRI | $\mathrm{b}_{2}$ | 2 | 2/7 |

### 2.2.3 INPUT/OUTPUT INSTRUCTIONS

| $\frac{\text { MNEMONIC }}{\text { IN }}$ | $\frac{\text { OPERANDS }}{\mathrm{b}_{2}}$ | $\frac{\text { BYTES }}{2}$ |  |
| :---: | :---: | :---: | :---: |
| MCYCLES/ <br> STATES |  |  |  |
| OUT | $\mathrm{b}_{2}$ | 2 | $3 / 10$ |

### 2.2.4 MACHINE INSTRUCTIONS

$\frac{\text { MNEMONIC }}{\text { HLT }}$ OPERANDS $\frac{\text { BYTES }}{1} \frac{\text { STATES }}{2 / 7}$

## DESCRIPTION

(A)-(input data from data bus), byte 2 is sent on bits A7-A0 and A15-A8 as the input device address. INP status is given on the data bus.
(Output data) $\leftarrow(A)$, byte 2 is sent on bits $A 7-A 0$ and $A 15-A 8$ as the output device address. OUT status is given on the data bus.

## DESCRIPTION

Halt, all machine operations stop. All registers are maintained. Only an interrupt can return the TMS 8080 to the run mode, Note that a HLT should not be placed in location zero, otherwise after the reset pin is active, the TMS 8080 will enter a nonrecoverable state (until power is removed), i.e., in halt with interrupts disabled. This condition also occurs if a HLT is executed while interrupts are disabled. HLTA status is given on the data bus.
(PC) - (PC) +1 , no operation.

$\frac{\text { MNEMONIC }}{\text { CALL }} \frac{\text { OPERANDS }}{\mathrm{b}_{3} \mathrm{~b}_{2}} \quad \frac{\text { BYTES }}{3} \quad$| M CYCLES/ |
| :---: |
| STATES |

Conditional call instructions for true flags:

| (f) |  | $5 / 17$ (Pass) |  |
| :--- | :--- | :--- | :--- |
| CC (carry) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 | $3 / 11$ (Fail) |
| CPE (parity) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 |  |
| CM (sign) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 |  |
| CZ (zero) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 |  |

Conditional call instructions for false flags:

| (f) |  | $5 / 17$ (Pass) |  |
| :---: | :---: | :---: | :---: |
| CNC (carry) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 | $3 / 11$ (Fail) |
| CPO (parity) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 |  |
| CP (sign) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 |  |
| CNZ (zero) | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 |  |
| Di |  | 1 | $1 / 4$ |
| EI |  | 1 | $1 / 4$ |
|  |  | 3 | $3 / 10$ |

Conditional jump instructions for true flags:

| (f) |  |  |  |
| :--- | :--- | :--- | :--- |
| JC | (carry) | $b_{3} b_{2}$ | 3 |
| JPE | (parity) | $b_{3} b_{2}$ | 3 |
| JM | (sign) | $b_{3} b_{2}$ | 3 |
| JZ | (zero) | $b_{3} b_{2}$ | 3 |

Conditional jump instructions for false flaas:

| (f) |  |  | 3/10 |
| :---: | :---: | :---: | :---: |
| JNC (carry) | $b_{3} b_{2}$ | 3 |  |
| JPO (parity) | $b_{3} b_{2}$ | 3 |  |
| JM (sign) | $b_{3} b_{2}$ | 3 |  |
| JNZ (zero) | $b_{3} b_{2}$ | 3 |  |
| PCHL |  | 1 | 1/5 |
| POP | PSW | 1 | 3/10 |
| POP | ${ }^{\text {r }}$ d | 1 | 3/10 |
| PUSH | PSW | 1 | 3/11 |
| PUSH | ${ }^{\text {r }}$ | 1 | 3/11 |
| RET |  | 1 | 3/10 |

## DESCRIPTION

[(SP)-1] [(SP)-2]-(PC), $\left.(S P)+-(S P)-2, \quad(P C)-<b_{3}\right\rangle<b_{2}>$, transfer $P C$ to the stack address given by $S P$, decrement $S P$ twice, and jump unconditionally to address given in bytes 2 and 3.

If $\left.(f)=1,[(S P)-1][(S P)-2]-(P C),(S P) \cdot-(S P)-2,(P S)-<b_{3}\right\rangle$ $<b_{2}>$, otherwise $(P C)-(P C)+3$. If the flag specified, $f$, is 1 , then execute a call. Otherwise, execute the next instruction.

If $\left.(f)=0,[(S P)-1][(S P)-2]-(P C),(S P) \leftarrow(S P)-2,(P C)-<b_{3}\right\rangle$ $\left.<\mathrm{b}_{2}\right\rangle$, otherwise $(\mathrm{PC}) \leftarrow(\mathrm{PC})+3$.

Disable interrupts. INTE is driven false to indicate that no interrupts will be accepted.
Enable interrupts. INTE is driven true to indicate that an interrupt will be accepted. Execution of this instruction is delayed to allow the next instruction to be executed before the INT input is polled.
( PC ) $\leftarrow \mathrm{b}_{3}><\mathrm{b}_{2}>$, jump unconditionally to address given in bytes 2 and 3 .

If $\left.(f)=1,(P C) \leftarrow<b_{3}\right\rangle<b_{2}>$, otherwise $(P C) \leftarrow(P C)+3$. If the flag specified, $f$, is 1, execute a JMP. Otherwise, execute the next instruction.

If $\left.(f)=0,(P C) \leftarrow<b_{3}\right\rangle\left\langle b_{2}\right\rangle$, othewise $(P C)+-(P C)+3$.
(PC) $-(H L)$
$(F) \leftarrow[(S P)],(A) \leftarrow[(S P)+1], \quad(S P) \leftarrow-(S P)+2$, restore the last stack values addressed by $S P$ into $A$ and $F$. Increment $S P$ twice. $\left(r_{d L}\right) \leftarrow[(S P)],\left(r_{d H}\right) \leftarrow[(S P)+1],(S P) \leftarrow(S P)+2$.
$[(S P)-1] \leftarrow(A),[(S P)-2] \leftarrow-(F),(S P) \leftarrow-(S P)-2$, save the contents of $A$ and $F$ into the stack addressed by $S P$. Decrement $S P$ twice. $[(S P)-1] \leftarrow\left(r_{d L}\right),[(S P)-2] \leftarrow\left(r_{d H}\right),(S P) \leftarrow(S P)-2$.
$(P C) \leftarrow[(S P)][(S P)+1]$, $(S P) \leftarrow-(S P)+2$, return to program at memory address given by last values in the stack. The SP is incremented by two.

$\underline{\text { MNEMONIC OPERANDS BYTES MCYCLES/ }}$| MTATES |
| :---: |

Conditional return instructions for true flags

| (f) |  |  |
| :--- | :--- | :--- |
| RC | (carry) | C |
| RPE | (parity) | P |
| RM | (sign) | S |
| RZ | (zero) | Z |

Conditional return instructions for talse flags:

| (f) |  |  | $3 / 11$ (Pass) |
| :---: | :---: | :---: | :---: |
| RNC (carry) | C | 1 | $1 / 5$ (Fail) |
| RPO (parity) | P | 1 |  |
| RP (sign) | S | 1 |  |
| RNZ (zero) | Z | 1 |  |
| RST |  | 1 | $3 / 11$ |

SPHL
1
$1 / 5$

### 2.2.6 REGISTER GROUP INSTRUCTIONS

| MNEMONIC | OPERANDS | BYTES | M CYCLES/ STATES |
| :---: | :---: | :---: | :---: |
| DCR | M | 1 | 3/10 |
| DCR | $\mathrm{r}_{\mathrm{a}}$ | 1 | 1/5 |
| DCX | $\mathrm{r}_{\mathrm{b}}$ | 1 | 1/5 |
| INR | M | 1 | 3/10 |
| INR | $\mathrm{r}_{\mathrm{a}}$ | 1 | 1/5 |
| INX | $r^{\text {b }}$ | 1 | 1/5 |
| LHLD | $\mathrm{b}_{3} \mathrm{~b}_{2}$ | 3 | 5/16 |
| LXI | $r_{b} b_{3}{ }^{2}$ | 3 | 3/10 |
| MVI | $\mathrm{M}, \mathrm{b}_{2}$ | 2 | 3/10 |
| MVI | $\mathrm{ra}_{\mathrm{a}} \mathrm{b}_{2}$ | 2 | 2/7 |
| MOV | $\mathrm{Mra}_{\mathrm{a}}$ | 1 | 2/7 |
| MOV | $\mathrm{ram}^{\text {M }}$ | 1 | 2/7 |
| MOV | $r^{\prime} 1^{r} a 2$ | 1 | 1/5 |
| SHLD | $b_{3} b_{2}$ | 3 | 5/16 |
| XCHG |  | 1 | 1/4 |
| XTHL |  | 1 | 5/18 |

## DESCRIPTION

If $(f)=1,(P C) \cdots[(S P)][(S P+1],(S P)<(S P)+2$. If the flag specified, $f$, is 1 , execute a RET. Otherwise, execute the next instruction.

If $(f)=0,(P C)-[(S P)][(S P)+1],(S P)+-(S P)+2$.
[(SP)-1] [(SP)-2] <-(PC) (SP)--(SP)--2, (PC) $\leftarrow 00000_{8}$ where $R$ is a 3 bit field in RST (RST $=3 R 78$ ). Transfer PC to the stack address given by SP, decrement SP twice, and jump to the address specified by $R$.
$(S P) \leftarrow(H L)$

## DESCRIPTION

(M) $-(M)-1$, decrement the contents of memory location specified by $H$ and $L .\{Z, S, P, C 1\}$
$\left(r_{a}\right) \leftarrow\left(r_{a}\right)-1$, decrement the contents of register $r_{a} .\{Z, S, P, C 1\}$ $\left(r_{b}\right) \leftarrow\left(r_{b}\right)-1$, decrement double registers $B C, D E, H L$, or $S P$. $(M)<-(M)+1$, increment the contents of memory location specified by $H$ and L. $\{Z, S, P, C 1\}$
$\left(r_{a}\right) \leftarrow\left(r_{a}\right)+1$, increment the contents of register $r_{a} \cdot\{Z, S, P, C 1\}$ $\left(r_{b}\right) \leftarrow\left(r_{b}\right)+1$, increment double registers $B C, D E, H L$, or $S P$. $(L) \leftarrow\left[<b_{3}><b_{2}>\right]$; $\left.\left.(H)-\left[<b_{3}\right\rangle<b_{2}\right\rangle+1\right]$, load registers $H$ and $L$ with contents of the two memory locations specified by bytes 3 and 2 .
$\left(r_{b H}\right)-<b_{3}>;\left(r_{b L}\right) \leftarrow<b_{2}>$, load double registers $B C, D E, H L$, or SP immediate with bytes 3, 2, respectively.
$(\mathrm{M}) \leftarrow<\mathrm{b}_{2}>$, store immediate byte 2 in the address specified by HL
$\left(r_{\mathrm{a}}\right) \leftarrow<\mathrm{b}_{2}>$, load register $\mathrm{r}_{\mathrm{a}}$ immediate with byte 2 of the instruc. tion.
(M) $-\left(r_{a}\right)$, store register $r_{a}$ in the memory location addressed by $H$ and $L$.
$\left(r_{a}\right)-(M)$, load register $r_{a}$ with contents of memory addressed by HL.
$\left(r_{a 1}\right) \leftarrow\left(r_{a 2}\right)$, load register $r_{a 1}$ with contents of $r_{a 2}, r_{a 2}$ contents remain unchanged.
[<b3 $\left.><\mathrm{b}_{2}\right\rangle$ ] $\leftarrow(\mathrm{L})$ : $\left.\left.\left.\left|<\mathrm{b}_{3}\right\rangle<\mathrm{b}_{2}\right\rangle+1\right)\right]:-(H)$, store the contents of $H$ and $L$ into two successive memory locations specified by bytes 3 and 2.
$(H)-(D) ;(L)-(E)$, exchange double registers $H L$ and $D E$
$(L) \rightarrow[(S P)],(H) \rightarrow[(S P)+1],(S P)=(S P)$, exchange the top of the stack with register HL.

| MNEMONIC | BYTES | DESCRIPTION | POSITIVE-LOGIC |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | REGISTER | HEX OPCODE |  | CLOCK |
|  |  |  | AFFECTED | D7-D4 | $\underbrace{\text { D3-D0 }}$ | CYCLES* |
| ACl | 2 | Add immediate to A with carry ${ }^{\dagger}$ |  | C |  | 7 |
| ADC M | 1 | Add memory to A with carry ${ }^{\dagger}$ |  | 8 | E | 7 |
| ADC r | 1 | Add register to $A$ with carry ${ }^{\dagger}$ | B | 8 | 8 | 4 |
|  |  |  | C | 8 | 9 |  |
|  |  |  | D | 8 | A |  |
|  |  |  | E | 8 | B |  |
|  |  |  | H | 8 | C |  |
|  |  |  | L | 8 | D |  |
|  |  |  | A | 8 | F |  |
| ADD M | 1 | Add memory to $\mathrm{A}^{\dagger}$ |  | 8 | 6 | 7 |
| ADD r | 1 | Add register to $A^{\dagger}$ | B | 8 | 0 | 4 |
|  |  |  | C | 8 | 1 |  |
|  |  |  | D | 8 | 2 |  |
|  |  |  | E | 8 | 3 |  |
|  |  |  | H | 8 | 4 |  |
|  |  |  | L | 8 | 5 |  |
|  |  |  | A | 8 | 7 |  |
| ADI | 2 | Add immediate to $\mathrm{A}^{\dagger}$ |  | C | 6 | 7 |
| ANA M | 1 | AND memory with $A^{\dagger}$ |  | A | 6 | 7 |
| ANAr | 1 | AND register with $\mathrm{A}^{\dagger}$ | B | A | 0 | 4 |
|  |  |  | C | A | 1 |  |
|  |  |  | D | A | 2 |  |
|  |  |  | E | A | 3 |  |
|  |  |  | H | A | 4 |  |
|  |  |  | L | A | 5 |  |
|  |  |  | A | A | 7 |  |
| ANI | 2 | AND immediate with $\mathrm{A}^{\dagger}$ |  | E | 6 | 7 |
| CA.LL | 3 | Call unconditional |  | C | D | 17 |
| CC | 3 | Call on carry |  | D | C | 11/17 |
| CM | 3 | Call on minus |  | F | C | 11/17 |
| CMA | 1 | Complement A |  | 2 | F | 4 |
| CMC | 1 | Complement carry |  | 3 | F | 4 |
| CMP M | 1 | Compare memory with $\mathrm{A}^{\dagger}$ |  | B | E | 7 |
| CMP r | 1 | Compare register with $A$ |  |  |  |  |
|  |  |  | B | B | 8 | 4 |
|  |  |  | C | B | 9 |  |
|  |  |  | D | B | A |  |
|  |  |  | E | B | B |  |
|  |  |  | H | B | C |  |
|  |  |  | L | B | D |  |
|  |  |  | A | B | F |  |
| CNC | 3 | Call on no carry |  | D | 4 | 11/17 |
| CNZ | 3 | Call on no zero |  | C | 4 | 11/17 |
| CP | 3 | Call on positive |  | F | 4 | 11/17 |
| CPE | 3 | Call on parity even |  | E | C | 11/17 |
| CPI | 2 | Compare immediate with $\mathrm{A}^{\dagger}$ |  | F | E | 7 |
| CPO | 3 | Call on parity odd |  | E | 4 | 11/17 |
| CZ | 3 | Call on zero |  | C | C | 11/17 |
| DAA | 1 | Decimal adjust $A^{\dagger}$ |  | 2 | 7 | 4 |

[^3]| MNEMONIC | BYTES | DESCRIPTION | POSITIVE-LOGIC |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | REGISTER AFFECTED | HEX OPCODE |  | CLOCK cyCLES |
|  |  |  |  | D7-D4 | D3-D0 |  |
| DAD B | 1 | Add B\&C to H\&L* |  | 0 | 9 | 10 |
| DAD C | 1 | Add D\&E to H\&L* |  | 1 | 9 | 10 |
| DAD H | 1 | Add H\&L to H\&L+ |  | 2 | 9 | 10 |
| DAD SP | 1 | Add stack pointer to H\&L \$ |  | 3 | 9 | 10 |
| DCR M | 1 | Decrement Memory |  | 3 | 5 | 10 |
| DCR r | 1 | Decrement Register $\stackrel{\text { S }}{ }$ | B | 0 | 5 | 5 |
|  |  |  | C | 0 | D |  |
|  |  |  | D | 1 | 5 |  |
|  |  |  | E | 1 | D |  |
|  |  |  | H | 2 | 5 |  |
|  |  |  | L | 2 | D |  |
|  |  |  | A | 3 | D |  |
| DCX B | 1 | Decrement B\&C |  | 0 | B | 5 |
| DCX D | 1 | Decrement D\&E |  | 1 | B | 5 |
| DCX H | 1 | Decrement H\&L |  | 2 | B | 5 |
| DCX SP | 1 | Decrement stack pointer |  | 3 | B | 5 |
| DI | 1 | Disable interrupts |  | F | 3 | 4 |
| EI | 1 | Enable interrupts |  | F | B | 4 |
| HLT | 1 | Halt |  | 7 | 6 | 7 |
| IN | 2 | Input |  | D | B | 10 |
| INR M | 1 | Increment memory |  | 3 | 4 | 10 |
| INR r | 1 | Increment register $\$$ | B | 0 | 4 | 5 |
|  |  |  | C | 0 | C |  |
|  |  |  | D | 1 | 4 |  |
|  |  |  | E | 1 | C |  |
|  |  |  | H | 2 | 4 |  |
|  |  |  | L | 2 | C |  |
|  |  |  | A | 3 | C |  |
| INX B | 1 | Increment B\&C register |  | 0 | 3 | 5 |
| INX D | 1 | Increment D\&E register |  | 1 | 3 | 5 |
| INX H | 1 | Increment H\&L register |  | 2 | 3 | 5 |
| INX SP | 1 | Increment stack pointer |  | 3 | 3 | 5 |
| JC | 3 | Jump on carry |  | D | A | 10 |
| JM | 3 | Jump on minus |  | F | A | 10 |
| JMP | 3 | Jump unconditional |  | C | 3 | 10 |
| JNC | 3 | Jump on no carry |  | D | 2 | 10 |
| JNZ | 3 | Jump on no zero |  | C | 2 | 10 |
| JP | 3 | Jump on positive |  | F | 2 | 10 |
| JPE | 3 | Jump on parity even |  | E | A | 10 |
| JPO | 3 | Jump on parity odd |  | E | 2 | 10 |
| JZ | 3 | Jump on zero |  | C | A | 10 |
| LDA | 1 | Load A direct |  | 3 | A | 13 |
| LDAX B | 1 | Load A indirect |  | 0 | A | 7 |
| LDAX D | 1 | Load A indirect |  | 1 | A | 7 |
| LHLD | 3 | Load H\&L direct |  | 2 | A | 16 |
| LXIB | 3 | Load immediate register pair B\&C |  | 0 | 1 | 10 |
| LXID | 3 | Load immediate register pair D\&E |  | 1 | 1 | 10 |
| LXIH | 3 | Load immediate register |  | 2 | 1 | 10 |
| LXISP | 3 | Load immediate stack pointer |  | 3 | 1 | 10 |

[^4]| MNEMONIC | BYTES | DESCRIPTION | POSITIVE-LOGIC |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | AFFECTED | $\underbrace{D 7-D 4}$ | $\underbrace{D 3-D 0}$ | CYCLES |
| MOV M, r | 1 | Moveregister to memoly | B | 7 | 0 | 7 |
|  |  |  | C | 7 | 1 |  |
|  |  |  | D | 7 | 2 |  |
|  |  |  | E | 7 | 3 |  |
|  |  |  | H | 7 | 4 |  |
|  |  |  | L | 7 | 5 |  |
|  |  |  | A | 7 | 7 |  |
| MOV r,M | 1 | Move memory to register | B | 4 | 6 | 7 |
|  |  |  | C | 4 | E |  |
|  |  |  | D | 5 | 6 |  |
|  |  |  | E | 5 | E |  |
|  |  |  | H | 6 | 6 |  |
|  |  |  | L | 6 | E |  |
|  |  |  | A | 7 | E |  |
| MOV $r_{1}, r_{2}$ | 1 | Move register to register | B, B | 4 | 0 | 5 |
|  |  |  | B, $C$ | 4 | 1 |  |
|  |  |  | B, ${ }^{\text {d }}$ | 4 | 2 |  |
|  |  |  | B, E | 4 | 3 |  |
|  |  |  | B, H | 4 | 4 |  |
|  |  |  | B, L | 4 | 5 |  |
|  |  |  | B, A | 4 | 7 |  |
|  |  |  | C, B | 4 | 8 |  |
|  |  |  | C, C | 4 | 9 |  |
|  |  |  | C, D | 4 | A |  |
|  |  |  | C, E | 4 | B |  |
|  |  |  | C, H | 4 | C |  |
|  |  |  | ..-C, L | 4 | D |  |
|  |  |  | C, A | 4 | F |  |
|  |  |  | D, B | 5 | 0 |  |
|  |  |  | D, C | 5 | 1 |  |
|  |  |  | D, D | 5 | 2 |  |
|  |  |  | D, E | 5 | 3 |  |
|  |  |  | D, H | 5 | 4 |  |
|  |  |  | H,L | 5 | 5 |  |
|  |  |  | D, A | 5 | 7 |  |
|  |  |  | E,B | 5 | 8 |  |
|  |  |  | E,C | 5 | 9 |  |
|  |  |  | E, D | 5 | A |  |
|  |  |  | E, E | 5 | B |  |
|  |  |  | E, H | 5 | C |  |
|  |  |  | E, L | 5 | D |  |
|  |  |  | E, A | 5 | F |  |
|  |  |  | H, B | 6 | 0 |  |
|  |  |  | H,C | 6 | 1 |  |
|  |  |  | H, D | 6 | 2 |  |
|  |  |  | H,E | 6 | 3 |  |
|  |  |  | H, H | 6 | 4 |  |
|  |  |  | H,L | 6 | 5 |  |
|  |  |  | H, A | 6 | 7 |  |
|  |  |  | L, B | 6 | 8 |  |


| MNEMONIC | BYTES | DESCRIPTION | POSITIVE-LOGIC |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | REGISTER | HEX O | CODE | CLOCK |
|  |  |  | AFFECTED | $\underbrace{\text { D7-D4 }}$ | $\underbrace{\text { D3-D0 }}$ | CYCLES* |
| MOV $\mathrm{r}_{1}, \mathrm{r}_{2}$ | 1 | Move register to register (continued) | L, C | 6 | 9 |  |
|  |  |  | L,D | 6 | A |  |
|  |  |  | L, E | 6 | B |  |
|  |  |  | L, H | 6 | C |  |
|  |  |  | L, L | 6 | D |  |
|  |  |  | L, A | 6 | F |  |
|  |  |  | A, B | 7 | 8 |  |
|  |  |  | A, C | 7 | 9 |  |
|  |  |  | A, D | 7 | A |  |
|  |  |  | A, E | 7 | B |  |
|  |  |  | A, H | 7 | C |  |
|  |  |  | A, L | 7 | D |  |
|  |  |  | A, A | 7 | F |  |
| MVI M | 2 | Move immediate memory |  | 3 | 6 | 10 |
| MVIr | 2 | Move immediate register | B | 0 | 6 | 7 |
|  |  |  | C | 0 | E |  |
|  |  |  | D | 1 | 6 |  |
|  |  |  | E | 1 | E |  |
|  |  |  | H | 2 | 6 |  |
|  |  |  | L | 2 | E |  |
|  |  |  | A | 3 | E |  |
| NOP | 1 | No operation | 4 | 0 | 0 | 4 |
| ORA M | 1 | OR memory with $\mathrm{A}^{\dagger}$ |  | B | 6 | 7 |
| ORA r | 1 | OR register with $\mathrm{A}^{\dagger}$ | B | B | 0 | 4 |
|  |  |  | C | B | 1 |  |
|  |  |  | D | B | 2 |  |
|  |  |  | E | B | 3 |  |
|  |  |  | H | B | 4 |  |
|  |  |  | L | B | 5 |  |
|  |  |  | A | B | 7 |  |
| ORI | 2 | OR immediate with $\mathrm{A}^{\dagger}$ |  | F | 6 | 7 |
| OUT | 2 | Output |  | D | 3 | 10 |
| PCHL | 1 | H\&L to program counter |  | E | 9 | 5 |
| POP B | 1 | Pop register pair B\&C off stack |  | C | 1 | 10 |
| POP D | 1 | Pop register pair D\&E off stack |  | D | 1 | 10 |
| POP H | 1 | Pop register pair H\&L off stack |  | E | 1 | 10 |
| POP PSW | 1 | Pop $A$ and flags off stack ${ }^{\dagger}$ |  | F | 1 | 10 |
| PUSH B | 1 | Push register pair B\&C |  | C | 5 | 11 |
| PUSH D | 1 | Push register pair D\&C |  | D | 5 | 11 |
| PUSH H | 2 | Push register pair H\&L on stack |  | E | 5 | 11 |
| PUSH PSW | 1 | Push A and Flags on stack |  | F | 5 | 11 |
| RAL | 1 | Rotate A left through carry |  | 1 | 7 | 4 |
| RAR | 1 | Rotate A right through carry |  | 1 | F | 4 |
| RC | 1 | Return on carry |  | D | 8 | 5/11 |
| RET | 1 | Return |  | C | 9 | 10 |
| RLC | 1 | Rotate A left ${ }^{\text {\% }}$ |  | 0 | 7 | 4 |
| RM | 1 | Return on minus |  | F | 8 | 5/11 |
| RNC | 1 | Return on no carry |  | D | 0 | 5/11 |
| RNZ | 1 | Return on no zero |  | C | 0 | 5/11 |
| RP | 1 | Return on positive |  | F | 0 | 5/11 |

[^5]| MNEMONIC | BYTES | DESCRIPTION | POSITIVE-LOGIC |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | AFFECTED | $\underbrace{D 7-D 4}$ | $\underbrace{\text { D3-D0 }}$ | CYCLES* |
| RPE | 1 | Return on parity even |  | E | 8 | 5/11 |
| RPO | 1 | Return on parity odd |  | E | 0 | 5/11 |
| RRC | 1 | Rotate A right ${ }^{\text {\% }}$ |  | 0 | F | 4 |
| RST | 1 | Restart |  |  |  | 11 |
|  |  |  | PC $\leftarrow-000016$ | C | 7 |  |
|  |  |  | PC $\leftarrow 000816$ | C | F |  |
|  |  |  | PC $<-0010_{16}$ | D | 7 |  |
|  |  |  | PC $\leftarrow 0001816$ | D | F |  |
|  |  |  | PC $<-002016$ | E | 7 |  |
|  |  |  | PC $\leftarrow-002816$ | E | F |  |
|  |  |  | PC +003016 | F | 7 |  |
|  |  | . | PC $<003816$ | F | F |  |
| RZ | 1 | Return on Zero |  | C | 8 | 5/11 |
| SBB M | 1 | Subtract memory from A with borrow ${ }^{\dagger}$ |  | 9 | E | 7 |
| SBB r | 1 | Subtract register from $A$ with borrow ${ }^{\dagger}$ | B | 9 | 8 | 4 |
|  |  |  | C | 9 | 9 |  |
|  |  |  | D | 9 | A |  |
|  |  |  | E | 9 | B |  |
|  |  |  | H | 9 | C |  |
|  |  |  | L | 9 | D |  |
|  |  |  | A | 9 | F |  |
| SBI | 2 | Subtract immediate from A with borrow ${ }^{\dagger}$ |  | D | E | 7 |
| SHLD | 3 | Store H\&L direct |  | 2 | 2 | 16 |
| SPHL | 1 | H\&L to stack pointer |  | F | 9 | 5 |
| STA | 3 | Store A direct |  | 3 | 2 | 13 |
| STAX B | 1 | Store A indirect |  | 0 | 2 | 7 |
| STAX D | 1 | Store A indirect |  | 1 | 2 | 7 |
| STC | 1 | Set carry $\ddagger$ |  | 3 | 7 | 4 |
| SUB M | 1 | Subtract memory from $A^{\dagger}$ |  | 9 | 6 | 7 |
| SUB r | 1 | Subtract register from $\mathrm{A}^{\dagger}$ | B | 9 | 0 | 4 |
|  |  |  | C | 9 | 1 |  |
|  |  |  | D | 9 | 2 |  |
|  |  |  | E | 9 | 3 |  |
|  |  |  | H | 9 | 4 |  |
|  |  |  | L | 9 | 5 |  |
|  |  |  | A | 9 | 7 |  |
| SUI | 2 | Subtract immediate from $\mathrm{A}^{\dagger}$ |  | D | 6 | 7 |
| XCHG | 1 | Exchange D\&E, H\&L registers |  | E | B | 4 |
| XRA M | 1 | Exclusive OR memory with $\mathrm{A}^{\dagger}$ |  | A | E | 7 |
| XRA r | 1 | Exclusive OR register with $\mathrm{A}^{\dagger}$ | B | A | 8 | 4 |
|  |  |  | C | A | 9 |  |
|  |  |  | D | A | A |  |
|  |  |  | E | A | B |  |
|  |  |  | H | A | C |  |
|  |  |  | L | A | D |  |
|  |  |  | A | A | F |  |
| XRI | 2 | Exclusive OR immediate with $\mathrm{A}^{\dagger}$ |  | E | E | 7 |
| XTHL | 1 | Exchange top of stack $\mathrm{H} \& \mathrm{~L}$ |  | E | 3 | 18 |

[^6]
## 3. TMS 8080 ELECTRICAL AND MECHANICAL SPECIFICATIONS

### 3.1 ABSOLUTE MAXIMUM RATINGS OVER OPERATING FREE-AIR TEMPERATURE RANGE (UNLESS OTHERWISE NOTED)*


"Stresses beyond those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions beyond those indicated in the "Recommended Operating Conditions" section of this specification is not implied. Exposure to absolute-maximum-rated conditions for extended periods may affect device reliability.
NOTE 1: Under absolute maximum ratings voltage values are with respect to the normally most negative supply voltage, $\mathrm{V}_{\mathrm{BB}}$ (substrate) Throughout the remainder of this data sheet, voltage values are with respect to $\mathrm{V}_{\text {SS }}$ unless otherwise noted.

### 3.2 RECOMMENDED OPERATING CONDITIONS

|  | MIN | NOM | MAX | UNIT |
| :---: | :---: | :---: | :---: | :---: |
| Supply voltage, $\mathrm{V}_{\text {BB }}$ | -4.75 | -5 | -5.25 | $\checkmark$ |
| Supply voltage, $\mathrm{V}_{\mathrm{CC}}$ | 4.75 | 5 | 5.25 | V |
| Supply voltage, $\mathrm{V}_{\text {DD }}$ | 11.4 | 12 | 12.6 | $\checkmark$ |
| Supply voltage, V ${ }_{\text {SS }}$ |  | 0 |  | V |
| High-level input voltage, $\mathrm{V}_{1 \mathrm{H}}$ (all inputs except clocks) (see Note 2) | 3.3 |  | $V_{C C}+1$ | V |
| High-level clock input voltage, $\mathrm{V}_{\mathrm{I}} \mathrm{H}(\boldsymbol{\phi})$ | $\mathrm{V}_{\mathrm{DD}}{ }^{-1}$ |  | $\vee_{\text {DD }}{ }^{+1}$ | V |
| Low-level input voltage, $\mathrm{V}_{\text {IL }}$ (all inputs except clocks) (see Note 3) | -1 |  | 0.8 | V |
| Low-level clock input voltage, $\mathrm{V}_{1} \mathrm{~L}(\mathrm{\phi})$ (see Note 3) | -1 |  | 0.6 | V |
| Operating free-air temperature, TA | 0 |  | 70 | C |

### 3.3 ELECTRICAL CHARACTERISTICS OVER FULL RANGE OF RECOMMENDED OPERATING CONDITIONS (UNLESS OTHERWISE NOTED)

|  | PARAMETER | TEST CONDITIONS | MIN TYP ${ }^{+}$ | MAX | UNIT |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 11 | Input current lany input except clocks and data bus) | $\mathrm{V}_{1}=0 \mathrm{~V}$ to $\mathrm{V}_{\mathrm{CC}}$ |  | $\pm 10$ | $\mu \mathrm{A}$ |
| $11(\%)$ | Clock input current | $\mathrm{V}_{1(\phi)}=0 \mathrm{~V}$ to $\mathrm{V}_{\text {DD }}$ |  | $\pm 10$ | $\mu \mathrm{A}$ |
| I/(DB) | Input current, data bus | $\mathrm{V}_{1(\mathrm{DB})}=0 \mathrm{~V}$ to $\mathrm{V}_{\text {CC }}$ |  | -100 | $\mu \mathrm{A}$ |
| II(hold) | Address or data bus input current during hold | $\mathrm{V}_{\text {I }}(\mathrm{ad})$ or $\mathrm{V}_{1}(\mathrm{DB})=\mathrm{V}_{\text {CC }}$ |  | 10 | $\mu \mathrm{A}$ |
|  |  | $\mathrm{V}_{\text {I }}(\mathrm{ad})$ or $\mathrm{V}_{\text {I( }}(\mathrm{DB})=0 \mathrm{~V}$ |  | -100 |  |
| $\mathrm{V}_{\mathrm{OH}}$ | High-level output voltage | $\mathrm{I}^{\text {OH }}=100 \mu \mathrm{~A}$ | 3.7 |  | V |
| V OL | Low-level output voltage | $\begin{aligned} & \mathrm{I} \mathrm{OL}(\mathrm{DB})=1.7 \mathrm{~mA}, \\ & \mathrm{I} \mathrm{OL}=0.75 \mathrm{~mA} \text { (any output except } \mathrm{DB} \text { ) } \end{aligned}$ |  | 0.45 | V |
| ${ }^{\text {I BB }}$ (av) | Average supply current from $\mathrm{V}_{\mathrm{BB}}$ | Operating at $\mathrm{t}_{\mathrm{C}}(\phi)=480 \mathrm{~ns}$, $T_{A}=25^{\circ} \mathrm{C}$ | -0.01 | -1 | mA |
| ICC(av) | Average supply current from $\mathrm{V}_{\mathrm{CC}}$ |  | 60 | 75 |  |
| ${ }^{\text {I D }}$ (av) | Average supply current from $\mathrm{V}_{\text {DD }}$ |  | 40 | 67 |  |
| $\mathrm{C}_{\mathrm{i}}$ | Capacitance, any input except clock | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}=\mathrm{V}_{\mathrm{DD}}=\mathrm{V}_{\mathrm{SS}}=0 \mathrm{~V}, \\ & \mathrm{~V}_{\mathrm{BB}}=-4.75 \text { to }-5.25 \mathrm{~V}, \mathrm{f}=1 \mathrm{MHz}, \end{aligned}$ <br> All other pins at 0 V | 10 | 20 | pF |
| $\mathrm{C}_{\mathrm{i}}(\boldsymbol{\phi})$ | Clock input capacitance |  | 5 | 10 |  |
| $\mathrm{C}_{0}$ | Output capacitance |  | 10 | 20 |  |

${ }^{\dagger}$ All typical values are at $T_{A}=25^{\circ} \mathrm{C}$ and nominal voltages.
NOTES: 2. Active pull-up resistors of nominally $2 k \leq 2$ will be switched onto the data bus when DBIN is high and the data input voltage is more positive than $\mathrm{V}_{1 \mathrm{H}} \mathrm{min}$
3. The algebraic convention where the most negative limit is designated as minimum is used in this specification for logic voltage levels only.
3.4 TIMING REQUIREMENTS OVER FULL RANGE OF RECOMMENDED OPERATING CONDITIONS (SEE FIGURE 2)

|  | . | MIN MAX | UNIT |
| :---: | :---: | :---: | :---: |
| $\mathrm{t}_{\mathrm{c}}(\mathrm{o})$ | Clock cycle time (see Note 5) | 4802000 | ns |
| $\mathrm{t}_{\mathrm{r}}(\rho)$ | Clock rise time | 550 | ns |
| $\mathrm{t}_{\mathrm{f}}\left(\mathrm{c}^{\circ}\right)$ | Clock fall time | 550 | ns |
| ${ }_{\text {tw }}$ (c) 1 ) | Pulse width, clock 1 high | 60 | ns |
| ${ }^{\text {tw }}$ (C) ${ }^{\text {c }}$ | Fulse width, clock 2 high | 220 | ns |
|  | Delay time, clock 1 low to clock 2 | 0 | ns |
|  | Delay time, clock 2 to clock 1 | 70 | ns |
| $\mathrm{t}_{\mathrm{d}}(\mathrm{o} 1 \mathrm{H} \cdot \mathrm{r} 2)$ | Delay time, clock 1 high to clock 2 (time between leading edges) | 130 | ns |
| $\mathrm{t}_{\text {su }}$ (da-¢, 1) | Data setup time with respect to clock 1 | 50 | ns |
| ${ }^{\text {tsu}}$ su(da-r, 2 ) | Data setup time with respect to clock 2 | 150 | ns |
| $\mathrm{t}_{\text {su }}$ (hold) | Hold input setup time | 140 | ns |
| $\mathrm{t}_{\text {sul }}$ (int) | Interrupt input setup time | 180 | ns |
| $\mathrm{t}_{\text {su }}$ (rdy ${ }^{\text {d }}$ | Ready input setup time | 120 | ns |
| $\mathrm{t}_{\mathrm{h}}$ (da) | Data hold time (see Note 6) | tPD(DBI) | ns |
| $t_{\text {h }}$ (hold) | Hold input hold time | 0 | ns |
| th(int) | Interrupt input hold time | 0 | ns |
| $t^{\text {f }}$ (rdy $)$ | Ready input hold time | 0 | ns |


6. The data input should be enabled using the DBIN status signal. No bus conflict can then occur and the data hold time requirement is thus assured
3.5 SWITCHING CHARACTERISTICS OVER FULL RANGE OF RECOMMENDED OPERATING CONDITIONS (SEE FIGURE 2)

|  | PARAMETER | TEST CONDITIONS | MIN MAX | UNIT |
| :---: | :---: | :---: | :---: | :---: |
| tPD(ad) | Propagation delay time, clock 2 to address outputs | $\begin{aligned} & C_{L}=100 \mathrm{pF}, \\ & R_{L}=1.3 \mathrm{ks} \end{aligned}$ | 200 | ns |
| tPD(da) | Propagation delay time, clock 2 to data bus |  | 220 | ns |
| tPD(cont) | Propagation delay time, clocks to control outputs |  | 120 | ns |
| tPD(DBI) | Propagation delay time, clock 2 to DBIN output |  | $25 \quad 140$ | ns |
| tPD(int) | Propagation delay time, clock 2 to INTE output |  | 200 | ns |
| tDI | Time for data bus to enter input mode |  | tPD(DBI) | ns |
| tPXZ | Disable time to high-impedance state during hold (address outputs and data bus) |  | 120 | ns |

The time that the address outputs and output data will remain stable after $\overline{W R}$ goes high, $\mathrm{t}_{\mathrm{W}}$ and $\mathrm{t}_{\mathrm{WD}} \geqslant \mathrm{t}_{\mathrm{d}}\left(\varphi_{1} \mathrm{H}-\mathrm{Q}_{2}\right)$.
The time between address outputs becoming stable and $\bar{W} \bar{R}$ going low, $\mathrm{t}_{\mathrm{A}} \mathrm{W} \leqslant 2 \mathrm{t}_{\mathrm{c}}(\phi)-\mathrm{t}_{\mathrm{d}}(\phi 1 \mathrm{H} \cdot \phi 2)-\mathrm{t}_{\mathrm{r}}(\phi)-120 \mathrm{~ns}$.
The time between output data becoming stable and $\overline{W R}$ going low, $\left.t_{D W} \geqslant t_{c}(\phi)-t_{d}(\phi) 1 H \cdot \phi 2\right)-t_{1}(\phi)-150 \mathrm{~ns}$
The following are relevant when interfacing to devices requiring $\mathrm{V}_{1} \mathrm{H}$ min of 3.3 V :
a) Maximum output rise time ( $\mathrm{t} T \mathrm{TH}$ ) from 0.8 V to 3.3 V is 140 ns with $C_{L}$ as specified for the propagation delay times above.
b) Maximum propagation delay times when measured to $V_{\text {ref }}(\mathrm{H})=3 \mathrm{~V}$ (instead of 2 V ) will be 60 ns more than as specified above with $C_{L}$ as specified.

$C_{L}$ includes probe and jig capacitance.
LOAD CIRCUIT
voltage waveforms (see notes a and b)

$\begin{array}{ll}\text { NOTES: a. This timing diagram shows timing relationships only, it does not represent any specific machine cycle. } \\ & \left.\text { b. Time measurements are made at the following reference voltages: Clock, } V_{r e f(H)}=9.5 \mathrm{~V}, \mathrm{~V}_{\text {ref }}(\mathrm{L})=1 \mathrm{~V} . \mathrm{Other} \text { inputs, } \mathrm{V} \text { ref(H) }=2 \mathrm{~V}, \mathrm{~V} \text { ref( }\right)=0.8 \mathrm{~V} .\end{array}$

| TMS 8080 |  |  |  |
| :---: | :---: | :---: | :---: |
| A10 | 1 | 40 | A11 |
| $\mathrm{V}_{\text {SS }}$ | 2 | 39 | A14 |
| D4 | 3 | 38 | A13 |
| D5 | 4 | 37 | A12 |
| D6 | 5 | 36 | A15 |
| D7 | 6 | 35 | A9 |
| D3 | 7 | 34 | A8 |
| D2 | 8 | 33 | A7 |
| D1 | 9 | 32 | A6 |
| DO | 10 | 31 | A5 |
| $\mathrm{V}_{\mathrm{BB}}$ - | 11 | 30 | A4 |
| RESET | 12 | 29 | A3 |
| HOLD | 13 | 28 | $V_{\text {DD }}$ |
| INT | 14 | 27 | A2 |
| 02 | 15 | 26 | A1 |
| INTE | 16 | 25 | AO |
| DBIN | 17 | 24 | WAIT |
| $\overline{W R}$ | 18 | 23 | READY |
| SYNC | 19 | 22 | \$1 |
| $\mathrm{V}_{\mathrm{CC}}$ | 20 | 21 | HLDA |

40-PIN CERAMIC PACKAGE

APPENDIX E

APPENDIX E
How to Align the Intecolor 8001 .

C O N T E N T S

### 1.0 SAFETY PRECAUTIONS

1.0.1 HIGH VOLTAGE
1.0.2 X-RADIATION PRECAUTIONS
2.0 INSTALLATION AND SERVICE ADJUSTMENTS
2.0.1 SERVICING PRECAUTIONS
2.0.2 AC LINE TAP SELECTOR
2.0.3 VERTICAL DEFLECTION
2.0.4 HORIZONTAL DEFLECTION
2.0.5 HIGH WOLTAGE ADJUSTMENT
2.0.6 FOCUS ADJUSTMENT
2.0.7 PURITY ADJUSTMENT
2.0.8 COLOR TEMPERATURE ADJUSTMENTS
2.0.9 TOP, BOTTOM, AND SIDE PINCUSHION ADJUSTMENT
2.0.10 CONVERGENCE ADJUSTMENT PRELIMINARIES
2.0.11 CONVERGENCE STATIC ADJUSTMENTS
2.0.12 FINAL CONVERGENCE
(C) 1976

### 1.0 SAFETY PRECAUTIONS

WARNING: The following precautions should be observed:

1. Do not install, remove, or handle the picture tube in any manner unless shatter-proof goggles are worn. People not so equipped should be kept away while picture tubes are handled. Keep picture tube away from the body while handling.
2. Part of the High Voltage is connected to the AC line directly. This circuitry, found on the Analog Module (100047), is isolated from the remainder of the circuitry by optical isolator, U3, and driver transformer, T101. Should service of the High Voltage be required it is recommended that an isolation transformer be inserted in the power line between the Intecolor ${ }^{\circledR} 8001$ and the AC supply before any service is performed. When the Chassis must be operated directly from the AC supply, the power plug should always be inserted in the correct polarity to connect the High Voltage common (emitter of Q5) to the ground side of the AC line. Check with a VOM or oscilloscope to see if a potential exists between this point and a known earth ground. A zero reading should be obtained. If any voltage reading is obtained, reverse the power plug and recheck for zero meter reading.
3. When service is required, observe the original lead dress. Extra precaution should be given to assure correct lead dress in the high voltage circuitry and video area. Where a short circuit has occurred, replace those components that indicate evidence of overheating. Always use the manufacturer's recommended replacement component.

NOTE: THE NOMINAL HIGH VOLTAGE FOR THE INTECOLOR ${ }^{\circledR} 8001$ 17" or 19" TERMINAL IS 25 KV . THE HIGH VOLTAGE MUST NOT, UNDER ANY CIRCUMSTANCES, EXCEED 27.5KV.

Each time a terminal's High Voltage requires servicing, measurements should be made at normal viewing settings of the Brightness Control. This will afford assurance that;

1. The High Voltage is within limits specified.
2. The High Voltage regulation circuit is functioning properly.
3. X-Radiation is at a minimum.

If the High Voltage measures abnormally high or the High Voltage Regulation Circuit is not functioning properly, the Terminal should be restored to normal operation through service or adjustments. (See 2.0.5 for High Voltage Adjustment procedure.)

IT IS IMPORTANT TO USE AN ACCURATE AND RELIABLE HIGH VOLTAGE METER.
1.0.2 X -RADIATION PRECAUTIONS

The primary source of X Radiation in this Terminal is the picture tube.

The tube utilized for the above mentioned function in the terminal is specifically constructed to limit X Radiation emissions.

For continued X -Radiation protection, the replacement tube must be the same type as the original, including suffix letter, or an ISC approved type.

### 2.0.1 SERVICING PRECAUTIONS

Purity, Color, Temperature, and Convergence adjustments for the Intecolor ${ }^{\circledR} 8001$ are essentially the same as for conventional shadow mask color tubes. Certain precautions should be taken, however, in servicing the Intecolor ${ }^{\circledR} 8001$ terminal.

Some precautions to observe while servicing the solid state chassis are listed below:

1. Always connect the ground lead of a test instrument to the chassis before connecting the positive lead; conversely, always remove the ground lead of a test instrument last.
2. Do not check for high voltage by drawing an arc. Use a high voltage meter or a high voltage probe with a VOM.
3. Do not bridge electrolytic capacitors since resultant surges may damage solid state devices.
4. Some transistors are equipped with heat sinks. Do not operate the transistor with the heat sink removed.
5. A11 soldering irons used where transistors and integrated chips are concerned should be 35 watt ( 6 volts) irons and grounded in such a way that no voltage will be applied to the solid state device during the soldering operation. This precaution is to prevent possible damage to the device due to excessive heat or voltage applied under no bias conditions.
6. When servicing the video circuitry it is recommended that an oscilloscope of at least 100 MHZ bandwidth, such as the Tektronix 454A, be used.

The AC Line Tap Selector is located inside the chassis on the right hand side as viewed from the rear (See Figure 2.0.2.1). In areas having a 115VAC 1ine supply, this tap should be left in the 115 VAC position. Other taps are shown depending on the line voltage.
2.0.3 VERTICAL DEFLECTION

At 115 volts line voltage adjust the VERTICAL HEIGHT CONTROL, R3, (See Figure 2.0.3.1) and the VERTICAL POSITION CONTROL, R4, so that the picutre is centered and there is a $12^{\prime \prime}$ wide by $10^{\prime \prime}$ high display. A suitable display is found by filling up the screen with a single character or erasing the screen with a background color.
2.0.4 HORIZONTAL DEFLECTION

Adjust the HORIZONTAL WIDTH CONTROL, R6, (Analog Module, 100047) (See Figure 2.0.3.1) so that the picture has a $12^{\prime \prime}$ wide by $10^{\prime \prime}$ high display. HORIZONTAL CENTERING is accomplished by adjusting R3 on the rear edge of the Display Generator Card, 100117. Adjusting the Pot R3 causes one character movements to the right or left of the screen.


## AC LINE TRANSFORMER TAP SELECTION

 FIGURE 2.0.2.1
2.0.5 HIGH VOLTAGE ADJUSTMENT

Preset High Voltage Adjustment Control R8 (Analog Module 100047) to $1 / 2$ clockwise, and Brightness Control R17, to maximum counterclockwise (minimum brightness).

Remove the High Voltage Anode Cap from the tube and connect a Pomona 非2900A or equivalent to the High Voltage Cap. CAUTION: BE SURE HV PROBE GROUND IS GROUNDED. İNSURE THAT ANODE CAP IS ISOLATED FROM ALL PERSONS AND EQUIPMENT. Adjust High Voltage Control, R8 for 25 KV .
2.0.6 FOCUS ADJUSTMENT

Create a full page of white dots on the CRT screen by utilizing the following procedure:

1. Select Foreground Color - WHITE
2. Select Background Color - BLACK
3. Press keyboard "." (period) and allow to repeat until screen is full of white dots.

Adjust the FOCUS pot (found on the right side (viewing from rear) of the Analog Card mounting bracket. Remove the external case with 6 screws) for optimum focus over the entire screen. (See Figure 2.0.6.1)

### 2.0.7 PURITY ADJUSTMENT

The Intecolor ${ }^{\circledR} 8001$ should always be facing either north or south during p.rity adjustment. This assures that any effect of the earth's normal magnetic field upon beam landing will be negligible when the terminal is placed in its normal viewing location.

The instrument should be at room temperature ( $60^{\circ} \mathrm{F}$ or above) for at least 30 minutes before set-up adjustments are made. Allow a minimum of ten minutes operation at high beam current (brightness full without bloom) before attempting purity or convergence adjustments.


Should any parts of the chassis become magnetized, it will be necessary to manually degauss the affected areas. Move a manual (GC 9317 or equivalent) degaussing coil slowly around those areas and the face of the CR Tube and slowly withdraw to a distance of six feet before disconnecting the coil from the AC power source.

Before performing the purity adjustments, the center of the raster must be converged and the dynamic convergence set roughly as explained in Section 2.0.12. Check that the focus control is properly set (See Section 2.0.6). The focus adjustment should be made with the brightness control set at maximum beam current without bloom.

1. Purity adjustments are most accurate while observing one screen only, preferably red. Erase the screen with the background color 'RED".
2. Loosen the yoke wing nuts and move the yoke to the rear as far as possible. (See Figure 2.0.7.1)
3. Rotate the purity magnets and adjustment tabs so that a clean red area is produced at the center of the screen. Push the yoke forward until a uniform red raster is obtained. Tighten the yoke wing nuts.
4. Erase the screen with the background color "WHITE". Check for a uniform white screen (see COLOR TEMPERATURE ADJUSTMENTS, Section 2.0.8, for procedure). If uniformity has not been obtained, reconverge the center of the screen and repeat the purity adjustments.
5. It should be noted that purity adjustments also affect the focus and DC Horizontal and Vertical screen positions and these parameters may have to be readjusted as outlined under Sections 2.0.3, 2.0 .4 , and 2.0.6.


YOKE; BLUE LATERAL, AND PURITY
LOCATIONS AND ADJUSTMENTS
FIGURE 2.0.7.1

1. Place a screen full of WHITE characters or ERASE the screen in WHITE. Turn the screen grid drive controls R14 (RED), R15 (GREEN), R16 (BLUE) (Analog Module 100047) to minimum drive (Fully CCW) then turn the BRIGHTNESS Control, R17 to maximum brightness (Fully CW).
2. Turn the RED contro1, R14, clockwise until the red vertical retrace raster line at the top of the screen
is just visible. Turn the GREEN Control, R15, clock-(f) wise until the green vertical retrace raster line at the top of the screen is just visible. Repeat the same for the BLUE Control, R16. ( $R-15)$
3. Adjust the BRIGHTNESS Contro1, R17, until there is no visible vertical retrace raster line and the brightness is at a comfortable viewing level with a minimum of color saturation.
4. Adjust each screen grid drive control, RED (R14), GREEN (R15), and BLUE (R16), until a white screen is obtained, or a $9300^{\circ} \mathrm{K}$ color temperature (WHITE).
2.0.9 TOP, BOTTOM, AND SIDE PINCUSHION ADJUSTMENT

Place a suitable test pattern on the screen such as all "+" (plus) symbols or all "." (periods). (See Section 2.0.6 for pattern set-up). Any color or WHITE may be used.

The top and bottom (Vertical) pin cushion adjustment is made, if necessary, by adjusting R5 on the Analog Module (100047) for straight horizontal lines at the top and bottom of the raster as shown in Figure 2.0.3.1 and Figure 2.0.9.1.

The side (Horizontal) pin cushion adjustment is made by adjusting R7 on the Analog Module (100047) for straight vertical lines on the left and right side of the raster.
$\mathrm{V}=$ Vertical
Pincushion= R7
H= Horizontal Pincushion= R5


## PINCUSHION ADJUSTMENT

FIGURE 2.0.9.1


HORIZONTAL AND VERTICAL RAMP ADJUSTMENTS
ANALOG MODULE (100047)
FIGURE 2.0.10.1.1

### 2.0.10 CONVERGENCE ADJUSTMENT PRELIMINARIES

The CONVERGENCE ADJUSTMENT PRELIMINARIES are necessary only if convergence cannot be obtained as outlined under FINAL CONVERGENCE ADJUSTMENTS (Section 2.0.12), or if these areas have required service or parts replacements, or the adjustment pots have been tampered with. An oscilloscope, such as the Tektronix 454, or equivalent will be necessary for these adjustments.

### 2.0.18.1 PRELIMINARY HORIZONTAL RAMP ADJUSTMENT

The Horizontal Ramp U11/6 amplitude is adjusted by R1 on the Analog Module (100047). The ramp is adjusted so that the positive peak is equal in height to the negative peak (symmetrical about ground or VH $=\mathrm{V}-$ ). (See Figure 2.0.10.1.1).

### 2.0.10.2 PRELIMINARY VERTICAL RAMP ADJUSTMENT

The VERTICAL RAMP U10/6 amplitude is adjusted by R2 on the Analog Module (100047) in the same manner as the HORIZONTAL RAMP ADJUSTMENT (See Figure 2.0.10.1.1).

### 2.0.10.3 PRELIMINARY HORIZONTAL PARABOLA ADJUSTMENT (U’7/3) RIGHT \& LEFT CENTER, TUBE AREAS 4 \& 5 (See Figure 2.0.12.2). <br> Adjust R9 on the Analog Module (100047) until the bottom of the Parabola is at Ground leve1. See Figure 2.0.10.3.1.

```
2.0.10.4 PRELIMINARY VERTICAL PARABOLA ADJUSTMENT (U8/3) TOP \& BOTTOM CENTER, TUBE AREAS \(2 \& 3\) (See Figure 2.0.12.2).
```

Adjust Rl0 on the Analog Module (100047) until the bottom of the Parabola is at ground level. See Figure 2.0.10.3.1.
2.0.10.5 HORIZONTAL AND VERTICAL RAMP ADJUSTMENTS.

Monitor the HORIZONTAL PARABOLA at U7/3 on the Analog Module (100047). Superimpose a small amount of the video signals (with a screen full of WHITE characters) by adding a small amount of the " B " trace (connect a scope probe to the collector of Q26 or Q27 or Q28) on the oscilloscope (CHOP, INVERT B, ADD) to the" A " trace (connected to U7/3). The above may also be accomplished
by simply connecting the " $A$ " channel Scope ground to a ground in the vicinity of $\mathrm{Q} 26, \mathrm{Q} 27$, or Q 28 . The video will be apparent on the parabola, as shown in Figure 2.0.10.5.1.

Adjust R1 until the superimposed video is as shown in Figure 2.0.10.5.1.

Monitor the VERTICAL PARABOLA at $\mathrm{U} 8 / 3$ and adjust R2 of the Analog Module (100047) until the end points of the parabola are equal in height.

The above procedure is shown in Figure 2.0.10.5.2.
2.0.10.6 VACANT
2.0.10.7 CORNER PARABOLA ADJUSTMENTS

TUBE AREAS 6, 7, 8, and 9 (See Figure 2.0.12.2)
The CORNER PARABOLA ADJUSTMENTS are made by R11, R12 and R13 on the Analog Module 100047 and monitoring the waveform as shown at U14/3 as in Figure 2.0.10.7.1. OFFSET is adjusted to zero by R13 by adjusting the waveform baseline to ground as shown in Figure 2.0.10.7.1, Waveform A.

BASELINE SLANT is adjusted by R12 on Analog Module (100047) as shown in B of Figure 2.0.10.7.1. Adjust for ${ }^{\mathrm{v}} \mathrm{SC}$ equal to " 0 " volts.

VERTICAL SYMMETRY is adjusted as shown in C of Figure 2.0.10.7.1 using R11 on Analog Module (100047). Alignment is made by adjusting R 11 until $+\mathrm{V}_{\mathrm{HC}}=-\mathrm{V}_{\mathrm{HC}}$.
2.0.10.8 HORIZONTAL, VERTICAL and CORNER PARABOLA TOUCH-UP

Touch up of the HORIZONTAL, VERTICAL, and CORNER PARABOLAS can best be accomplished by monitoring the waveforms on the J1 on the Convergence Module (100014).
A. Adjust the HORIZONTAL PARABOLA offset, $\mathrm{V}_{\mathrm{HP}}$ with R9 on the Analog Module (100047) by monitoring the waveform at J1/1 on the Convergence Module (100014) as shown in Figure 2.0.10.8.1, A.
B. Adjust the VERTICAL PARABOLA offset, $\mathrm{V}_{\mathrm{VP}}$ with R10 on the Analog Module (100047) by monitoring the waveform at J1/5 on the Convergence Module (100014) as shown in Figure 2.0.10.8.1, B
C. Adjust the CORNER PARABOLA offset, $\mathrm{V}_{\mathrm{CP}}$ with R13 on the Analog Module (100047) by monitoring the waveform at J1/7 on the Convergence Module (100014) as shown in Figure 2.0.10.8.1, C.


HORIZONTAL AND .VERTICAL PARABOLA ADJUSTMENTS
FIGURE 2.0.10.3.1

Superimposed Video


HORIZONTAL PARABOLA VIDEO ADJUSTMENT
FIGURE 2.0.10.5.1

A. CORNER PARABOLA OFFSET



1. Adjust R11 Analog Module (100047) to $+\mathrm{V}_{\mathrm{HC}}=-\mathrm{V}_{\mathrm{HC}}$
2. Monitor waveform at U14/3 on Analog Module (100047)

CORNER PARABOLA ADJUSTMENTS
FIGURE 2.0.10.7.1

A.


VERTICAL PARABOLA
J1/5 ON CONVERGENCE MODULE (100014)
B.

Adjust R13 on Analog Module (100047) until $\mathrm{V}_{\mathrm{CP}}=$


## C.

HORIZONTAL, VERTICAL, AND CORNER PARABOLA TOUCH-UP
FIGURE 2.0.10.8.1

Place a dot pattern on the screen in the following manner from the Keyboard.

$$
\begin{aligned}
& \text { Define FOREGROUND COLOR AS "WHITE"' } \\
& \text { BACKGROUND COLOR AS "BLACK" } \\
& \text { Depress "." (period) Key and allow to repeat } \\
& \text { until the screen is full of white dots. }
\end{aligned}
$$

The above will fill up the screen with dots. Now place " + " symbols utilizing the keyboard as shown in Figure 2.0.11.1

Turn all the pots on the Convergence Module (100014) to the straight up position as shown in Figure 2.0.11.3.

Now adjust the static magnets and the Blue Lateral Magnet to align the " + " symbols R,G,B, colors in Screen Sector 1, as shown in Figure 2.0.11.2, so as to appear as "WHITE". This will occur when the RED, GREEN, AND BLUE colors are accurately superimposed on top of each other. With the exception of BLUE lateral which is explained below.

For the above to be accurate the tube must have been externally degaused, the Purity adjusted, the FOCUS R18 adjusted for sharp, and the BRIGHTNESS, R17, Analog Module (100047), set for a low level with the color temperature being set to $9600^{\circ} \mathrm{K}$ as $\operatorname{explained}$ in previous sections. DO NOT ATTEMPT FURTHER CONVERGENCE UNLESS THE above has been previously performed. (See Sections

The beams move at approximately the same angle as the convergence magnets are offset from the vertical plane. Blue, since it is mounted in the vertical plane moves


+ . . . . . . . . . + . . . . . . . . . +
$\qquad$
$\qquad$
$\qquad$

FIGURE 2.0.11.1
the beam up and down vertically; red and green move the respective beams on a line at about a $60^{\circ}$ angle from the vertical. The blue lateral magnet moves all three beams in the horizontal plane, the blue beam in one direction and the red and green beams in the opposite direction in a 5 to 1 ratio. The blue beam has the greatest lateral shift.

The thumb screw adjustment of red, green, and blue center convergence magnets can be rotated in either direction continuously. Flux change is accomplished by rotating the pole position of the magnets, not by moving the magnets farther from or closer to the respective guns.

Adjust the Static Blue so that the Blue in the center of the screen is superimposed on the RED and GREEN.

### 2.0.12 FINAL CONVERGENCE

Touch up the center convergence with the pots R13 (GREEN), R14 (RED) and R15 (BLUE) on the Convergence Module (100014) as shown in Figure 2.0.11.2 and Figure 2.0.11.3.

Once center convergence has been adjusted proceed to the next convergence Screen Sector, 2, as shown in Figure 2.0.11.2. Proceed with the alignment in the order of the sector numbers as shown in Figure 2.0.11.2. After each Sector is aligned, check and touch up the center convergence. Note that the adjustment pots on the Convergence Module (100014) are arranged in the same location as each Screen Sector as viewed on the tube face (and the component side of the board) and the trio of pot groups in each sector are arranged as GREEN, RED, AND BLUE corresponding to the location of the GREEN, RED, and BLUE electron beams as viewed from the tube face.

When completed with the above, touch up each Screen Sector as needed in the SAME ORDER as outlined above. Do not violate the order of the Screen Sector numbers in the adjustment procedure.

Never attempt a convergence procedure without first setting the Convergence Module (100014) pots to the center position as shown in Figure 2.0.12.3 and following the Screen Sector numbers. It is seldom necessary for the static magnets to be adjusted unless shipment vibration causes convergence coil or static magnet movements or unless convergence coil or yoke replacements become necessary.

| 8 | 2 | 6 |
| :---: | :---: | :---: |
| 5 | 1 | 4 |
| 9 | 3 | 7 |

CR TUBE CONVERGENCE SECTORS (SCREEN VIEW)
FIGURE 2.0.11.2

(TOP VIEW)
NOTE: Green and Red Pots are interchanged on al1 17" Tubes.
FIGURE 2.0.12.3

TABLEOF O O O N TENTS


## $-1 \times-1$

| TERMS: |  |
| :---: | :---: |
| TERM | DESCRIPTION |
| Address | A 16 bit number assigned to a memory location corresponding to its sequential position. |
| Bit | The smallest unit of information which can be represented. (A bit may be in one of two states, 0 or l). |
| Byte | A group of 8 contiguous bits occupying a single memory location. |
| Console | Refers to the 8001 CRT Display as the output device, and the 8001 keyboard as the input device. Allows operator interface with the CPU operating system. |
| Instruction | The smallest single operation that the computer can be directed to execute. |
| Object Program | A program which can be loaded directly into the computer's memory and which requires no alteration before execution. An object program is usually on paper tape, and is produced by assembling (or compiling) a source program. Instructions are represented by binary machine code in an object program. |
| Program | A sequence of instructions which, taken as a group, allow the computer to accomplish a desired task. |
| Source Program | A progam which is readable by a programmer but which must be transformed into object program format before it can be loaded into the computer and executed. Instructions in an assembly language source program are represented by their assembly language mnemonic. |
| System Program | A program written to help in the process of creating user programs. |

```
TERMS -- (Continued):
```

TERM

| User Program | A program written by the user to make the <br> computer perform any desired task. |
| :--- | :--- |
| Word | A group of 16 contiguous bits occupying <br> two successive memory locations. (2 bytes). |

ABBREVIATIONS:

Space Bar
nnn represents a number in binary format.
nnn represents a number in decimal format.
nnn represents a number in octal format.
nnn represents a number in octal format.
nnn represents a number in hexadecimal format.

Shaded portions of CPU/operator dialog represent Console output.

CPU O.S. COMMANDS AND MESSAGES
2.0 CPU OPERATING SYSTEM (O.S.)

STARTING ADDRESS - 100 When in 8708 ERASÀBLE PROM

All arguments are in hexadecimal form.
A RAM TEST 2,3,4?
B GO BACK TO CRT O.S.

D DISPLAY IN HEXADECIMAL FORMAT
D low address, high address
Memory from low address to high address is displayed in hexadecimal form.

E END
E address
Endfile mark is created; 60 null characters are written on punch device

F FILL MEMORY
F low address, high address, data
Memory from low address to high address is filled with data.

G GO TO
G Address, bkptl, bkpt2
Program control is transferred to address. Breakpoints are set at bkptl and bkpt2. When break points are executed, all of the CPU registers are automatically displayed.

H HEXADECIMAL ARITHMETIC
H number, number sp
The sum and difference of the two numbers is printed in hexadecimal.

L LOAD HEXADECIMAL TAPE
L Bias address
A hexadecimal format tape is read into memory at tape address plus bias address.

M low address, high address, destination address
A block of memory from low address to high address is moved to location destination address.

N
PUNCH NULL

N

Sixty null characters are punched.

R BAUD RATE FOR SECOND RS-232 CHANNEL
$R$ rate number

The rate number must be between 1 and 7. See the "How to Use the 8001" Manual.

S SUBSTITUTE
S address Sp
Memory at address is displayed, and can be modified by typing in new data. Termination with space opens next sequential address, termination with carriage return ends command.
$X$ EXAMINE REGISTERS OR MEMORY

X reg ident
Register is displayed, and can be modified as in the $S$ command.

W WRITE HEXADECIMAL

W low address, high address
Memory from low address to high address is punched in hexadecimal format.

MESSAGES
. CPU O.S. ready to accept commands
? Error. Reenter command

```
3.0 INTECOLOF: 8001 CONFIGURATION
3.1 I/O SYSTEM
```

The Intecolor ${ }^{\circledR} 3001$ can support a number of input/output devices, from the CRT display and the RS232C I/O to devices supplied by the user. In general, it may be convenient to have two devices which can perform the same function, but to use them for different purposes at various times. For example, if a program is being assembled, you might want the program listing to be written on one device, while any system messages not relevant to the assembly would be written on a separate device.

The I.O system described below permits this type of change. Devices may be assigned functions via the System Monitor $S$ command (see Section 4.2.ll) or via the user's program. That is, it is possible to write programs which read from several different input devices and write to several different output devices of the program's choosing, without requiring any human intervention.

## 3.l.1 LOGICAL AND PHYSICAL DEVICES

Regardless of how many I/O devices a particular Intecolor ${ }^{\circledR} 8001$ has, there are only four operations which can be performed to any of them. For example, a WRITE operation can be performed either to the RS232C channel 1 to a host computer or a high speed tape system. All system programs and user-written programs, therefore, access four LOGICAL DEVICES (i.e., a WRITE device) which are then translated to a FHYSICAL DEVICE (i.e., a high speed tape) by the I/O system.

The four logical devices available to programs are:

| CONSOLE | An interactive, character-oriented device used <br> for both input and output. |
| :--- | :--- |
| READER | A character-oriented, input-only device which <br> transfers data on command and signals the <br> program when where is no more data (an end-of- <br> file condition). |
| WRITE | A character-oriented, output-only device which <br> accepts a character from the program and re- <br> cords it on some external medium. |
| LIST | A character-oriented, output-only device which <br> accepts a character from the program and records <br> it on some external medium in human readable form. |

Each of these four logical devices may be associated with one of four physical devices at any instant, giving a total of 16 physical devices. The mapping from logical to physical devices is specified by an I/O status byte which resides in memory and is accessible to system and user programs via substitute command. The address of the $I / O$ status byte is $9 F 9 \varnothing$ hex. A pointer to the I/O status byte is also contained in memory locations $\varnothing \varnothing 36$ and $\varnothing \varnothing 37$ (low byte of pointer, high byte of pointer). The possible mappings appear as follows:

I/O Status Byte:
Initially

| ${ }^{\text {A }} 7{ }^{\text {A }} 6$ | $\begin{array}{ll}\text { A } & \\ & \\ 5 & 4\end{array}$ | $\begin{array}{lll}\text { A } & \\ \\ & \\ \end{array}$ | $\begin{array}{ccc}\text { A } & & A \\ & 1 & 0\end{array}$ |
| :---: | :---: | :---: | :---: |
| 10 | 00 | 00 | 10 |

$$
\begin{array}{ll}
A_{7} A_{6}=\text { LIST FIELD } & A_{1} A_{0}=\text { CONSOLE FIELD } \\
A_{5} A_{4}=\text { PUNCH FIELD } & A_{3} A_{2}=\text { READER FIELD }
\end{array}
$$

LOGICAL DEVICES
I/O DEV FIELD
PHYSICAL DEVICES

| CONSOLE | 00 | RS 232 Channel 1 |
| :---: | :---: | :---: |
|  | 01 | RS 232 Channel 2 |
|  | 10 | $\begin{aligned} & \text { CR Tube = Console Output } \\ & \text { Keyboard= Console Input } \end{aligned}$ |
|  | 11 | (user console device) |
| READER | 00 | RS232 Channel 1 |
|  | 01 | RS232 Channel 2 |
|  | 10 | Keyboard |
|  | 11 | (user reader device 1) |

LOGICAL DEVICES I/O DEV FIELD PHYSICAL DEVICES

| WRITE | 00 | RS232 Channel l |
| :---: | :---: | :---: |
|  | 01 | RS232 Channel 2 |
|  | 10 | CR Tube |
|  | 11 | (user punch device l) |
| LIST | 00 | RS232 Channel l |
|  | 01 | RS232 Channel 2 |
|  | 10 | CR Tube |
|  | 11 | (user list device l) |

At cold start or system reset, the $I / O$ status byte is set equal to 82 H , causing the CR Tube and keyboard to be selected for console I/O and LIST, while the RS232 Channel l is selected for both READ and WRITE.

### 3.1.2 I/O SUBROUTINES

The way in which a program performs an I/O operation to any of the four logical devices is by calling the appropriate subroutine supplied by the $I / O$ system. The available subroutines and their locations in memory are given in the following table:

| ROUTINE | FUNCTION |  |
| :--- | :--- | :--- |
|  | MEMORY LOCATION |  |
| CI | Console input |  |
| CO | Console Output | 103 H |
| RI | Reader input | 109 H |
| PO | Punch output | 106 H |
| LO | List output | 10 CH |
| SO | Console String Output | 10 FH |
|  |  |  |

The rest of this section gives a description and examples of how to call these subroutines.

This routine returns a character received from the selected console device to the caller in the A register. The A register and the condition bits are affected by this operation.

Example:

$$
\begin{array}{ll}
\text { Assembly } & \text { Language } \\
\\
\text {... } & \\
\text { CALL } & \text { CI } \\
\text { STA } & \text { DATA } \\
\text {... } &
\end{array}
$$

CO - CONSOLE OUTPUT
CO transmits a character, passed from the calling program in the A register, to the device selected for console output. The A register and the condition bits are affected.

Example:
Assembly Language

| MVI | A,'.' |  |
| :--- | :--- | :--- |
| CALL | CO | ;PRINT '.' ON CONSOLE |

RI - READER INPUT
RI returns a character read from the reader device in the A register. If no character was read from the device (i.e., end of file), the CARRY condition bit is set equal to 1 , and the A register is zeroed. If data is ready, the CARRY bit is zeroed. If no character is received from the physical device then striking any key causes an end of file to be simulated and control is returned to the calling program.

Example:
Assembly Language

| CALL | RI |  |
| :--- | :--- | :--- |
| JC | EOF | E END OF FILE |
| STA | DATA |  |

PO transmits a character from the calling program to the device selected as the punch device. PO is identical in format to $C O$.

LO - LIST OUTPUT
LO performs the same function to the selected list device as CO and PO do to their selected devices.

SO - CONSOLE STRING OUTPUT
SO transmits a character string to the device selected for console output. A pointer to the beginning of the string is passed from the calling program in the HL register pair. The string should be terminated by a byte having the value 239 (decimal). SO also provides repeat loops of the form: ..., 237, N, Dl, D2, ..., DM, 238, ... where $N$ is the repeat count for the string of bytes D l through DM.

Example:
Assembly Language

| LXI | H, STR |
| :--- | :--- |
| CALL | SO |

STR: DB 'AB', 237, 3, 'CD', 238, 'EFG', 239
This example will print 'ABCDCDCDEFG' on the console device.

FLOPPY TAPE I/O SUBROUTINES

Three I/O subroutines are provided for the Intecolor Floppy Tape. These routines are:

| ROUTINE | FUNCTION | MEMORY LOCATION |
| :---: | :---: | :---: |
| TWR | Write to Floppy Tape | - 6866017 |
| TRD | Read from Floppy Tape | D133H 68005 H |
| TVF | Compare memory with Floppy Tape | 0036168400 |

The Floppy Tape is a block-transfer device. One record is written per track. The inputs from the calling program to each of the three I/O routines are:

```
HL register pair - pointer to memory buffer
DE register pair - byte count
\begin{tabular}{rl} 
A register \(\quad\) - & Tape drive/track code: \\
& BIT3 - DRIVE: \(\emptyset\) or 1 \\
& BITS2- \(\quad\) - Track: \(\varnothing\) through 7
\end{tabular}
```

After calling any one of the routines, the A register will contain a status code and will have been tested (ORA A):

| $A=\varnothing$ | : | No Errors |
| :---: | :---: | :---: |
| $\mathrm{A}=2$ |  | Keyboard Abort (Pressing any key on the keyboard during the data transfer will abort the operation) |
| $\mathrm{A}=4$ | : | Buffer too large for write. |
| A $=6$ | : | Buffer too small for read. |
| $\mathrm{A}=8$ | : | Read Failure: A complete, correctly formatted record could not be read from the tape. |
| $A=10$ | : | Checksum error. |
| $\mathrm{A}=12$ | : | Verify failure. A mismatch was detected between data in memory and data read from the tape during a memory compare operation (TVF). |

Also, after calling any of the routines, the $H L$ register pair will point one byte past the last byte manipulated in the memory buffer.

### 3.1.3 USER-SUPPLIED DEVICES

This section describes the necessary steps in hooking up a user-supplied I/O device to the I/O system.

The I/O subroutines described in Section 3.3.2 assume that programs (called drivers) exist which perform the actual transfer of data between $I / O$ devices and the CPU. For instance, when the console input routine is called, it checks to see which physical device is assigned to the console, and then branches to the driver appropriate to the device Therefore, when the user supplies his own device, he must:

1) Write a program to perform the data transfer, making sure that the program saves and restores any CPU registers it uses that are not specifically changed by the I/O subroutine.
2) Store a JMP to this driver's address in the appropriate location as defined in the following table:

MEMORY LOCATION

9F91H
9F94H
9F97H
9F9AH 9F9DH

USE
$\begin{array}{lll}\text { USER DEFINED CONSOLE INPUT } \\ \text { USER DEFINED CONSOLE OUTPUT } \\ \text { USER DEFINED READER (1) } \\ \text { USER DEFINED WRITE (1) } \\ \text { USER DEFINED LIST } & (1)\end{array}$

Thus, if the user supplied a custom built listing device, he would write a driver to transfer data to it in an appropriate manner, then store the JMP to the driver's address at location 9F9DH. By assigning LIST=3, his device would receive any listing output generated.

### 4.0 CPU OPERATING SYSTEM

The Intecolor 8001 CPU O.S. enables the operator to easily manipulate the contents of memory, read and produce MAG tapess and execute programs.

The CPU O.S., and all Intecolor ${ }^{\circledR} 8001$ system software in general, use the last 80 memory locations after the refresh area for storage of temporary data. Therefore, if the operator runs a program beginning in these locations, and then uses the CPU O.S. Text Editor, or Assembler, he must re-load these 80 bytes of his program before running it again. Alternatively, programs could be written beginning at any higher location. Then system programs and user programs could be executed in any order, without requiring the re-load operation.

For a 25 line system these locations are $8 F B O H$ to $8 F F F H$. The 48 line system uses locations 9 FBOH to $9 F F F H$.

The CPU O.S. is the operator's interface to the 8080 CPU , and controls loading and execution of user programs, and to some extent the debugging of user programs. Figure 4-1 illustrates memory utilization during various stages of system software use. While the CPU O.S. is running, it uses an area at the top of memory for data storage and scratch work.

```
4.1 CPU OPERATING SYSTEM IMPLEMENTATION AND EXECUTION
```


## 4.l.1 CPU O.S. IMPLEMENTATION

The Intecolor ${ }^{\circledR} 8001$ CPU O.S. program is implemented on two E PROM modules, which are pre-installed into each Intecolor 8001 with Option 34. This allows the CPU to be used with great ease, as it is not necessary to wait for lengthy paper-tape loading operations. All that is required to go on-line with CPU O.S. is to turn the Intecolor 8001 on, hit the ESCAPE key, and then the CPU O.S. key, and begin execution.

## 4.l.2. STARTING SYSTEM MONITOR

To begin operating the CPU O.S., press two keys in sequence, 'ESCAPE', (CPU O.S.) and the Intecolor 8001 will automatically jump to the starting address of the CPU O.S.

### 4.2 CPU O.S. OPERATION AND COMMANDS

The commands consist of a single letter typed into the Intecolor ${ }^{\circledR} 8001$ keyboard followed by a number of arguments, possibly none. The arguments are separated, if there are more than one, by spaces or commas. A command is terminated and executed by typing a carriage return or space, depending upon the command.

A (RAM TEST 2,3,4?) $[4$ is visual $]$

### 4.2.1 B COMMAND (BACK TO CRT O.S.)

### 4.2.2 D COMMAND (DISPLAY DATA)

The format of the $D$ command is:

D low address, high address
Low address is a valid 16 bit memory address.

High address is a valid 16 bit memory address equal to or greater than low address.

Description: Upon execution of this command, memory data from (low address) to (high address) is displayed upon the list device (normally the $C R$ tube). Data are displayed in hexadecimal form. Up to sixteen bytes per line are printed, preceded by the hexadecimal address of the first byte of that line. A carriage return is forced after a byte having a low order digit of $F$ in its memory address is printed.

Example: Enter at the keyboard the command:
.DlOF, 123 (Cr)
and the CR Tube will display:
0l0F AA

0120 EF $12 \quad 3456$
where memory locations 0lOF through 0123 are assumed to contain

$12 \quad 3456$
the $D$ command should be used only to examine memory contents. To punch the memory contents onto a tape, the $W$ command should be used. These commands produce a tape in the proper formats, while the $D$ command causes a simple sequence of characters to be output.

Error conditions:

1. If low address or high address is greater than 16 bits, only the last 4 hex digits of the argument will be used as the address.

Example: The command

$$
. D 30010, A B 0013 \text { (Cr) }
$$

is equivalent to the command

$$
. D 0010,0013(\mathrm{Cr})
$$

2. If low address is greater than high address, only the one byte at low address will be displayed.

Example: The command:
.Dl0,6
is equivalent to the command
.Dlo,lo
3. Non-existent memory is equivalent to a string of bytes all containing FF H .

Example: If memory address 2000 H- 2010 H are invalid, then the command:
.D2000,2010
will cause the teletype to print:

2000 FF FF FF FF FF FF FF FF FF FF FF FF FF FF FF FF 2010 FF
4. If low address or high address contains an invalid character, or if high address is omitted, the CR Tube will immediately display '?(Cr) (lf). and await the next command.

Example: If the user attempts to enter the number OG as an address, the following will be displayed:
.DOG?

### 4.2.3 F COMMAND (FILL MEMORY WITH CONSTANT)

The format of the F command is:

F low address, high address, data
Low address is a valid 16 bit memory address.
High address is a valid 16 bit memory address equal to or greater than low address.

Data is an 8 bit data value.

Description: Execution of this command causes memory locations (low address) through (high address) to be filled with the constant (data).

Example: The command:

$$
. \mathrm{F} 7,14, \mathrm{AA}(\mathrm{Cr})
$$

will set bytes 0007 through 0014 equal to AA $H$.

0007 AA AA AA AA AA AA AA AA AA 0010 AA AA AA AA AA

Error Conditions:
l. If low address of high address is greater than 16 bits (or data is greater than 8 bits), only the last 4 (or 2) hex digits will be used.

Example: The command:

$$
. F 7 A B 0007,0014, \text { FFACAA (Cr) }
$$

is equivalent to the command:

$$
. \mathrm{F} 0007,0014, \mathrm{AA}(\mathrm{Cr})
$$

2. If low address is greater than high address, data will replace only the byte at low address.

Example: If locations 7, 8, and 9 contain AA $H, B B H$, and $C C H$, execution of the command:

$$
. F 7,1,33 \text { (Cr) }
$$

will cause memory to appear as follows:
000733 BB CC
3. If a non-existent memory address is specified, this command has no effect.
4. If low address, high address, or data contain an invalid character, the CR Tube will immediately display '? (Cr) (lf).' and await the next command.

Example: If the user tries to enter $B Q$ as data, the following will be displayed:

$$
. \mathrm{F} 0012,14, \mathrm{BQ} ?
$$

```
4.2.4 G COMMFIND (GO TO)
```

The format of the $G$ command is:

```
    G address, bkptl, bkpt2
```

Address, bkptl, and bkpt2 are valid 16 bit hexadecimal memory addresses.

Description: The G cominand causes program control to be transferred to location address. If either bkptl or bkpt2 is specified, a breakpoint will be set in the program at the corresponding address(es). The specified address must correspond to the first byte of a program instruction. If either breakpoint is encountered during program execution, the CPU O.S. will save and display all program status (CPU registers and condition bits), clear all existing breakpoints, and take control. The user may then examine and/or modify registers or memory, or use any other monitor commands. This feature allows the user to debug portions of a program.

If address is not specified, the program status is restored and the saved value of the program counter is used as the new starting address.

Example: The command:

G24A
will cause program execution to begin at location 24 AH , with no breakpoints being set.

The command:
G,12C
will cause a breakpoint to be set at 12 CH , and program execution to resume at the address indicated by the saved value of the program counter.

The command:

G
will cause program execution to resume at the address indicated by the saved value of the program counter, with all status restored and no breakpoints set.

Error Conditions:

1. If address is greater than 16 bits, only the last 4 hex digits of the argument will be used as the address.

Example: The command:

> .G3C0010 (Cr)
is equivalent to the command
.G0010 (Cr)
2. If address is a non-existent memory address, the system will attempt to transfer control and then return to the CRT O.S. with no response. The CPU O.S. must then be manually restarted.

### 4.2.5. H COMMAND (HEXADECIMAL ARITHMETIC)

The format of the H command is:
.H number, number Sp
Number is a 16 bit hexadecimal number.

Description: The $H$ command is designed to aid the user in performing hexadecimal arithmetic while using the CPU O.S. It causes the sum and difference it arguments to be displayed in two-s complement hexadecimal form. This command is terminated by a space, rather than by a carriage return.

Example:

$$
\begin{array}{lll}
. \mathrm{HlE}, 5 \mathrm{C} & \text { OO7A }
\end{array}
$$

Error Conditions:

1. If either number is greater than 16 bits, only the last 4 hex digits are used.

Example: The command:
. HOOABC, 23Sp
is equivalent to the command:
. $\mathrm{HOABC}, 23 \mathrm{Sp}$
2. If number contains an invalid character, the $C R$ Tube will immediately display '?(Cr) (lf).' and await the next command.

Example: If the user attempts to enter OlP, the following will be displayed:
.HOlP?

### 4.2.6 I COMMAND (RESET CRT TO STATE So)

The format of the $I$ command is:

I causes the same action as the CPU reset key being typed.

### 4.2.7 L COMMAND (LOAD HEXADECIMAL FILE)

The format for the $L$ command is:

L bias address

Bias Address is a 16 bit two's complement hexadecimal number.

Description: This command loads tape written in hexadecimal format (using the $W$ command) into memory. The address at which the tape is loaded is determined by adding the address on the tape to the bias address using two's complement arithmetic. The bias may be negative, but in this case must be in two's complement form. If the tape was produced using an $E$ command with a non-zero entry point address (see section 4.2.ll), control will be transferred to that location in memory. Otherwise, the CPU O.S. will remain in control and request another command.

Example: If a tape was used which began at location 0100 H , the following command:
. LFFBO (Cr)
will cause the tape to be read and loaded into location 50 H . $(1000+\mathrm{FFB}=50)$.

NOTE: If an error occurs while reading the tape (such as a checksum error), the CPU O.S. will immediately stop reading the tape, display '? (Cr) (Lf).' and await the next command. The operation may be retried by backing up the tape to any point before the last colon and issuing another $L$ command, since each data word specifies the address at which it is to be loaded. The CPU O.S. will read up to the first colon it encounters, and then begin loading data.

Note that this means that, if you wish to change data in locations in memory, it is not necessary to regenerate an entirely new tape with the change; instead you may read in the original tape, then read in a patch tape which reloads only the erroneous locations.

Error Conditions:

1. If the bias address is greater than 16 bits, only the last 4 hex digits are used as the bias address.

Example: The command:

## .LOOFFBO (Cr)

is equivalent to the command:
.LFFBO (Cr)
2. If an invalid character is present in the bias address, the CR Tube will immediately display '*(Cr) (lf).' and await the next command.

Example: If the user attempts to enter GOO as a bias address, the following will be displayed:

### 4.2.8 M COMMAND (MOVE MEMORY)

The format of the $M$ command is:
.M low address, high address, destination address

Low address is a valid l6 bit memory address.

High address is a valid 16 bit memory address equal to or greater than low address.

Destination address is a valid 16 bit memory address.

Description: The M command causes the block of memory from low address through high address to be moved to the locations in memory beginning at destination address.

Example: If memory appears as follows:

| LOCATIONS |  | DATA |
| :--- | :--- | :--- |
| $0300-0304$ |  | contain |
| 0200-0204 | Ol020304 |  |
| contain | AlA2A3A4 |  |

Then the command:
M200,204,300
will cause the following:

LOCATIONS
$\begin{array}{lll}\text { 0300-0304 } & \text { contain } & \text { AlA2A3A4 } \\ 0200-0204 & \text { contain } & \text { AlA2A3A4 }\end{array}$

Note: The movement is performed byte by byte: the byte at low address is moved to destination address, then low address +1 is moved to destination address+l, etc. Therefore, the MOVE command may be used to fill memory with a byte or sequence of bytes.

Example: If location 0300 H contains $\mathrm{FF} H$, the command

$$
. \mathrm{M} 300,310,301 \text { (Cr) }
$$

will cause locations 300 through 310 to contain FF H. The FF at 300 is moved to 301, then the byte at 301 (which is now FF), is moved to 302, and so on.

Error Conditions:

1. If any address is greater than 16 bits, only the last 4 hex digits are used as the address.

Example: The command:

$$
. \mathrm{MOO} 02,303,00405(\mathrm{Cr})
$$

is equivalent to the command:

$$
\mathrm{Ni} 302,303,405 \text { (Cr) }
$$

2. If low address is greater than high address, only one byte will be moved from low address to destination address.

Example: The command:

$$
. \mathrm{M} 300,2 \mathrm{FO}, 100(\mathrm{Cr})
$$

is equivalent to the command:

$$
. \mathrm{M} 300,300,100 \text { (Cr) }
$$

3. If low address through high address specifies a nonexistent range of memory, bytes of $\mathrm{FF} H$ will be moved to the memory locations specified by destination address.

Example: If locations 2000 H through 2005 are non-existent, the commend:

$$
\text { . M2000, 2005, } 100 \text { (Cr) }
$$

will cause locations 0100 H through 0105 H to contain FF H .
4. If an invalid character is entered in an address, the CR Tube will display '?(Cr) (lf).' and await the next command.

Example: If the user attempts to enter OBAG as the destination address, the following will be displayed:
Ml00,10F , OBAG*
4.2.9. R COMMAND (BAUD RATE SELECT)

The format of the $R$ command is

R rate value

The rate value must be between 1 and 7. See chart below.

| NUMBER | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| NORMAL BAUD <br> RATE | 110 | 150 | 300 | 1200 | 2400 | 4800 | 9600 |
| HIGH SPEED <br> BAUD RATE | 880 | 1200 | 2400 | 9600 | 19,200 | 38,400 | 76,800 |

```
4.2.10 S COMMAND (SUBSTITUTE MEMORY)
```

The $S$ command is used to display and/or modify the contents of individual memory locations. It is used as follows:

1. Type an S, followed by the hexadecimal address of the first memory location you wish to display. Type space.
2. The data from the selected address is displayed, followed by a dash (-).
3. To modify memory, type in the new data followed by a space or a carriage return. If you do not wish to modify the contents of that location, do not type any data in, but only type a space or carriage return.
4. If a space was typed in step 3, the next memory location will be displayed as in step 2. If a carriage return was typed, operation will be returned to the CPU O.S.

Example: The contents of the first four bytes of memory is 00 AlCE FF . You wish to change it to 00 A 3 CE ll.
. SOOOOSp00Sp Al - A3Sp CE - Sp FF - llCr

User entries are unshaded. Display back is shaded.

Error Conditions:

1. If address is greater than 16 bits, or the data to be substituted is greater than 8 bits, only the last 4 or 2 hex digits respectively are used.

Example: The following sequence is equilvalent to the previous example:
. SOAB0000Sp 00 - Sp Al - BA3Sp CE - Sp FF - OllCr
2. If an invalid character is encountered, the CR Tube will immediately display '?(Cr) (lf).' and await the next command.

### 4.2.11 X COMMAND (EXAMINE AND MODIFY REGISTERS)

The format of the X command is:

X(C.R.) PRINTS ALL REGISTERS, ANNOTATED!
Reg ident is a single character specifying a CPU register as follows:
$\mathrm{A}=\mathrm{A}$ register
$B=B$ register
$C=C$ register
$\mathrm{D}=\mathrm{D}$ register
$\mathrm{E}=\mathrm{E}$ register
F = Flag byte, displayed in the form as it is stored by the instruction PUSH PSW
$\mathrm{H}=\mathrm{H}$ register
$\mathrm{L}=\mathrm{L}$ register
$M=H$ and $L$ registers combined (16 bi"ts)
$P=$ Program counter (16 bits)
$S=$ Stack pointer (l6 bits)

Note: The format of the flag byte $F$ is:

A
S Z O C OPIC

| Sign bit | State of carry bit |
| :--- | :--- |
| Zero bit | Always 1 |
| Always 0 | State of parity bit |
| Auxiliary carry bit | Always 0 |

Description: The X command is used to display and/or modify CPU registers. It operates similar to the $S$ command, as follows:

1. Type an $X$, followed by the register identifier.
2. The data from the selected register is displayed, followed by a dash (-). Four hexadecimal digits are displayed for $M, P$, and $S$; two hex digits for the other register identifiers.
3. To modify the register, type in the new data followed by a space or a carriage return. If you do not wish to modify the register, type only the space or carriage return.
4. If a space was typed in step 3, the next register in alphabetical order is displayed. If carriage return was typed, the X command is terminated. If a space is typed after register $S$ has been displayed, the command is terminated, this being the last register identifier in the list.

Example: The A, B, C, and D registers contain AAH, BBH, CCH, and DDH, respectively. You wish to change the $B$ and $C$ registers to OOH and FFh , respectively.

XASp AA- Sp BB- OOSp CC- FFSp DD-Cr

Note: Values set by the X-command will become the actual contents of the registers after execution of the next GO command.

The values displayed by the X -command are the contents of the registers prior to the execution of the last breakpoint set by the GO command. These displayed values, however, will reflect any changes of register "contents" made by the execution of X -commands since this last breakpoint.

## Error Conditions:

1. If the data to be substituted is greater than 16 bits for registers M, P, S, or 8 bits for the other register identifiers, only the last 4 or 2 hex digits respectively are used.
2. If an invalid register identifier or character is encountered, the CR Tube will immediately display '?(Cr) (Lf).' and await the next command.

### 4.2.12 E COMMAND (END FILE)

The format of the E command is:

> E address

Address is a valid 16 bit memory address.
Description: The E command causes an end-of-file mark and sixty null characters to be written at the end of a hexadecimal output file. The end of file mark is hexadecimal record of length 00. (See Appendix D). If address is 0 or absent, the $L$ command which loads the file will return control to the CPU O.S. If address is non-zero, the $L$ command will transfer control to that memory address immediately after loading the file.

The format of the W command is:

W low address; high address
Low address is a valid 16 bit memory address.

High address is a valid 16 bit memory address equal to or greater than low address.

Description: The $W$ command is used to output memory locations low address through high address to the system punch device in hexadecimal format. A series of $W$ commands may be issued in order to punch various non-contiguous memory locations onto a continuous strip of tape.

Any series of $W$ commands should be terminated with an $E$ command in order to punch a termination character, so that when the tape is read it will be handled properly.

Example: If memory locations l through 3 contain 53F8EC, the command: .W0001,0003 (Cr)
produces:
:0300010053F 8ECC5
(See Appendix D for an explanation of tape format.)

Error Conditions:

1. If low address or high address is greater than 16 bits, only the last 4 hex digits of the argument will be used as the address.

Example: The command:

$$
\text { WAB0010, } 100 \text { (Cr) }
$$

is equivalent to the command:

$$
\text { W0010, } 100 \text { (Cr) }
$$

2. If low address is greater than high address, only the one byte at low address will be written:

Example: The command:
.WlO, O(Cr)
is equivalent to the command:

> .Wl0,10 (Cr)
3. Non-existent memory is equivalent to a string of bytes all containing FF H.
4. An invalid character in either address will cause the CR Tube to display '?(Cr) (lf).' and await the next command.

Example: If the user attempts to enter $3 Z$ as low address, the following will be displayed:
.W3Z?

### 4.2.14 N COMMAND (NULL PUNCH)

The N command consists only of the letter N followed by a carriage return and causes 60 null characters to be written on the punch device.

## APPENDIX A

-- INSTRUCTION SUMMARY --


| RPl | The first register of register pair RP |
| :---: | :---: |
| RP2 | The second register of register pair RP |
| sign | The sign bit |
| SP | The l6-bit stack pointer register |
| SRC | Source register or memory byte |
| zero | The zero bit |
| XY | The value obtained by concatenating the values X and Y |
| [ ] | An optional field enclosed by brackets |
| ( ) | Contents of register or memory byte enclosed by parentheses |
| $\leftarrow$ | Replace value on lefthand side of arrow with value on righthand side of arrow |

CARRY BIT INSTRUCTIONS
Format:

| [LABEL:] | CODE |  |
| :--- | :--- | :--- |
| CODE | DESCRIPTION |  |
| STC | $($ carry $\longleftarrow$ | Set carry |
| CMC | $($ carry $) \longleftarrow$ | Complement carry |

Condition bits affected: Carry

Format:

| [LABEL: ] | INR | REGM |
| :--- | :--- | :--- |
| [LABEL: ] | -or- |  |
| [LABEL: ] | DCR | REGM |
| [LABEL: ] | CMA |  |
| [LAR- | -or- |  |
|  | DAA |  |


| Code | Description |  |
| :---: | :---: | :---: |
| INR | $(\mathrm{REGM}) \longleftarrow(\mathrm{REGM})+1$ | Increment register REGM |
| DCR | $(\mathrm{REGM}) \longleftarrow(\mathrm{REGM})-1$ | Decrement register REGM |
| CMA | $(\mathrm{A}) \longleftarrow(\bar{A})$ | Complement accumulator |
| DAA | If $\left(A_{0}-A_{3}\right)>9$ or (aux. carry $=1$, <br> (A) <br> (A) +6 <br> Then if $\left(A_{4}-A_{7}\right)>9$ or (carry) $=$ $1(A)=(A)+6 * 2^{4}$ | Convert accumulator contents to form two decimal digits |


| Condition bits affected: | INR,DCR | : Zero, sign, parity |
| :--- | :--- | :--- |
|  | CMA | DAA |
|  | : Zone |  |
|  |  |  |

NOP INSTRUCTION

Format:
[LABEL:] NOP

| Code | Description |
| :---: | :---: |
| NOP | ------------- No operation |

[^7]Format:

| [LABEL: $]$ | MOV | DST,SRC |
| :--- | :--- | :--- |
|  | - or- |  |
| $[$ LABEL : ] | CODE | RP |

NOTE: SRC and DST not both $=\mathrm{M}$
NOTE: $R P=B$ or $D$


Condition bits affected: None

REGISTER OR MEMORY TO ACCUMULATOR INSTRUCTIONS

Format:

|  | [LABEL:] | CODE R |  |
| :---: | :---: | :---: | :---: |
| Code | Description |  |  |
| ADD | (A) $\leftarrow$ | (A) $+($ REGM $)$ | Add REGM to accumulator |
| ADC | (A) $\leftarrow$ | $(\mathrm{A})+(\mathrm{REGM})+($ carry $)$ | Add REGM to accumulator with carry |
| SUB | (A) | (A) - (REGM) | Subtract REGM from accumulator |
| SBB | (A) $\leftarrow$ | (A) - (REGM) - (carry) | Subtract REGM from accumulator with borrow |
| ANA | (A) | (A) AND (REGM) | AND accumulator with REGM |
| XRA | (A) $\leftarrow$ | (A) XOR (REGM) | EXCLUSIVE-OR accumulator with REGM |


| Code | Description |
| :---: | :---: |
| ORA | (A) <-_ (A) OR (REGM) OR accumulator with REGM |
| CMP | Condition bits set by (A) -(REGM)Compare REGM with <br> accumulator |

Condition bits affected:
ADD, ADC, SUB, SBB: Carry, sign, zero, parity, aux. carry ANA, XRA, ORA: Sign, zero, parity. Carry is zeroed.
CMP: Carry, sign, zero, parity, aux. carry. Zero set if (A)=(REGM) Carry reset if (A) < (REGM) Carry set if (A) $\geqslant$ (REGM)

ROTATE ACCUMULATOR INSTRUCTIONS

Format:
[LABEL:] CODE


Condition bits affected: Carry

REGISTER PAIR INSTRUCTIONS
Format:

| [LABEL: ] | CODE1 | RP |
| :--- | :--- | :--- |
| [LABEL: ] | -or- |  |
|  | CODE2 |  |

Note: For PUSH and POP, RP=B,D,H or PSW For DAD, INX, and $D C X, R P=B, D, H$, or $S P$

| Codel | Description |  |
| :---: | :---: | :---: |
| PU'SH | $\begin{aligned} & ((S P)-1) \leftarrow \leftarrow(R P 1),((S P)-2) \longleftarrow(R P 2), \\ & (S P) \longleftarrow \leftarrow(S P)-2 \end{aligned}$ | Save RP on the <br> stack <br> RP=A saves accumulator <br> and condition bits. |
| POP | $\begin{aligned} & (\mathrm{RP} 1) \longleftarrow((\mathrm{SP})+1), \quad(\mathrm{RP} 2) \longleftarrow((\mathrm{SP})), \\ & (\mathrm{SP}) \longleftarrow(\mathrm{SP})+2 \end{aligned}$ | Restore RP from the stack |
| DAD | $(H L) \longleftarrow(H L)+(R P)$ | ```RP=A restores accumulator and condition bits. Add RP to the l6-bit number in H and L.``` |
| INX | $(\mathrm{RP}) \leftarrow(\mathrm{RP})+1$ | Increment RP by 1 |
| DCX | $(\mathrm{RP}) \leftarrow(\mathrm{RP})-1$ | Decrement RP by 1 |
| Code2 | Description |  |
| XCHG | $(\mathrm{H}) \longleftrightarrow(\mathrm{D}), \quad(\mathrm{L}) \longleftrightarrow(\mathrm{E})$ | Exchange the 16 bit number in $H$ and $L$ with that in $D$ and $E$. |
| XTHL | $(\mathrm{L}) \longleftrightarrow((\mathrm{SP})), \quad(\mathrm{H}) \longleftrightarrow((\mathrm{SP})+\mathrm{l})$ | Exchange the last values saved in the stack with $H$ and L. |
| SPHL | $(\mathrm{SP}) \longleftarrow \mathrm{CH}^{(H):(L) ~}$ | Load stack pointer from H and L . |

Condition bits affected:

PUSH, INX, DCX, XCHG, XTHL, SPHL: None
POP : If RP=PSW, all condition bits are restored from the stack, otherwise none are affected.
DAD : Carry

IMMEDIATE INSTRUCTIONS

Format:

| [LABEL: ] | LXl | RP, DATA16 |
| :--- | :--- | :--- |
|  | -or- |  |
| [LABEL: ] | MV1 | REGM, DATA |
|  | -or- |  |
| [LABEL: ] | CODE | REGM |

Note: $R P=B, D, H$, or $S P$

| CODE | DESCRIPTION |  |
| :---: | :---: | :---: |
| LXI | $(\mathrm{RP}) \longleftarrow$ DATA 16 | Move 16 bit immediate Data into RP |
| MVI | $($ REGM $) \longleftarrow$ DATA | Move immediate DATA into REGM |
| ADI | $(\mathrm{A}) \leftarrow-\ldots$ (A) + DATA | Add immediate data to accumulator |
| ACI | $(\mathrm{A}) \longleftarrow(\mathrm{A})+\mathrm{DATA}+($ carry $)$ | Add immediate data to accumulator with carry |
| SUI | $(\mathrm{A}) \longleftarrow$ ( L$)-$ DATA | Subtract immediate data from accumulator |
| SBI | $(\mathrm{A}) \longleftarrow$ (A) - DATA - (carry) | Subtract immediate data from accumulator with borrow |
| ANI | (A) $\longleftarrow$ (A) AND DATA | AND accumulator with immediate data |
| XRI | (A) $\longleftarrow$ (A) XOR DATA | EXCLUSIVE-OR accumulator with immediate data |
| ORI | $(\mathrm{A}) \longleftarrow$ (A) OR DATA | OR accumulator with immediate data |
| CPI | Condition bits set by (A)-DATA | Compare immediate data with accumulator |

Condition bits affected:

LXI, MVI: None
ADI, ACI, SUI, SBI: Carry, sign, zero, parity, aux. carry
ANI, XRI, ORI: zero, sign, parity. Carry is zeroed.
CPI: Carry, sign, zero, parity, aux. carry. Zero is set if (A)= DATA Carry reset if (A) < DATA Carry set if (A) $\geqslant$ DATA

DIRECT ADDRESSING INSTRUCTIONS

Format:

|  | LABEL:] CODE |  |
| :---: | :---: | :---: |
| CODE | DESCRIPTION |  |
| STA | $(\mathrm{ADDR}) \longleftarrow$ ( $\left.{ }^{( }\right)$ | Store accumulator at location ADDR |
| LDA | $(\mathrm{A}) \longleftarrow$ (ADDR) | Load accumulator from location ADDR |
| SHLD | $(A D D R) \leftarrow(L),(A D D R+1) \leftarrow(H)$ | Store $L$ and $H$ at ADDR and ADDR+1 |
| LHLD | $(\mathrm{L}) \longleftarrow($ ADDR $),(\mathrm{H}) \longleftarrow($ ADDR +1$)$ | Load $L$ and $H$ from ADDR and ADDR+1 |

Condition bits affected: None

Format:
$\begin{array}{lll}\text { [LABEL: ] } & \text { PCHL } \\ & \text {-or- } & \\ \text { [LABEL: ] } & \text { CODE }\end{array}$

| CODE | DESCRIPTION |  |
| :---: | :---: | :---: |
| PCHL | $(\mathrm{PC}) \longleftarrow(\mathrm{HL})$ | Jump to location specified by register H and L |
| JMP | $(\mathrm{PC}) \longleftarrow \mathrm{ADDR}$ | Jump to location ADDR |
| JC | If $($ carry $)=1, \quad(P C) \longleftarrow \underset{\text { ADDR }}{\text { AD }}=3$ If $($ carry $)=0, \quad(P C) \longleftarrow+3$ | Jump to ADDR if carry set |
| JNC | $\begin{aligned} & \text { If }(\text { carry })=0, \quad(P C) \longleftarrow \quad \text { ADDR } \\ & \text { If }(\text { carry })=1, \quad(P C) \longleftarrow(P C)+3 \end{aligned}$ | Jump to ADDR if carry reset |
| JZ | $\begin{aligned} & \text { If }(\text { zero })=1,(P C) \longleftarrow \text { ADDR } \\ & \text { If }(\text { zero })=0,(P C) \longleftarrow(P C)+3 \end{aligned}$ | Jump to ADDR of zero set |
| JNZ | $\begin{aligned} & \text { If }(\text { zero })=0, \quad(P C) \longleftarrow \text { ADDR } \\ & \text { If }(\text { zero })=1, \quad(P C) \longleftarrow(P C)+3 \end{aligned}$ | Jump to ADDR if zero reset |
| JP | If $($ sign $)=0, \quad(P C) \longleftarrow \begin{aligned} & \text { ADDR } \\ & \text { If }(\text { sign })\end{aligned}=1, \quad(P C) \longleftarrow(P C)+3$ | Jump to ADDR if plus |
| JM | If $($ sign $)=1, \quad(P C) \longleftarrow \begin{aligned} & \text { ADDR } \\ & \text { If }(\text { sign })\end{aligned}=0, \quad(P C) \longleftarrow(P C)+3$ | Jump to ADDR if minus |
| JPE |  | Jump to ADDR if parity even |
| JPO | If (parity) $=0$, <br> $(\mathrm{PC}) \longleftarrow \mathrm{ADDR}$ <br> If $($ parity $)=1,(P C) \longleftarrow(P C)+3$ | Jump to ADDR is parity odd |

Condition bits affected: None

Format:
[LABEL:]
CODE
ADDR

| CODE | DESCRIPTION |
| :---: | :---: |
| CALL | $\begin{aligned} ((S P)-1) \leftarrow(P C H),((S P)-2) \leftarrow(P C L), & (S P) \leftarrow(S P)+2, \quad(P C) \leftarrow \text { ADDR } \\ & \text { call subroutine and push return } \\ & \text { address onto stack } \end{aligned}$ |
| $C \mathrm{C}$ |  |
| CNC | $\begin{array}{ll} \text { If }(\text { carry })=0,((S P)-1 \leftarrow(P C H), & ((S P)-2) \leftarrow(P C L),(S P) \leftarrow(S P)+2, \\ \text { If }(\text { carry })=1,(P C) \leftarrow(P C)+3 & \text { Call subroutine if carry reset } \end{array}$ |
| CZ | $\begin{aligned} & \text { If }(\text { zero })=1,((S P)-1)-(P C H),((S P)-2) \leftarrow(P C L),(S P) \leftarrow-(S P)+2, \\ & \text { If (zero })=0,(P C) \leftarrow(P C)+3 \quad \text { Call subroutine if zero set } \end{aligned}$ |
| CNZ | ```If (zero) = 0, ((SP)-1) \leftarrow(PCH), ((SP)-2) \leftarrow(PCL), (SP) \leftarrow(SP)+2, (PC) \leftarrow- ADDR If (zero) = 1, (PC) \leftarrow(PC)+3 Call subroutine if zero reset``` |
| CP | $\begin{aligned} & \text { If }(\text { sign })=0,((S P)-1) \leftarrow(P C H),((S P)-2) \leftarrow(P C L),(S P) \leftarrow(S P)+2, \\ & \text { If }(\text { sign })=1,(P C) \leftarrow-(P C)+3 \quad \text { Call subroutine if sign plus } \end{aligned}$ |
| CM | $\begin{aligned} & \text { If }(\text { sign })=1,((\mathrm{SP})-1) \leftarrow(\mathrm{PCH}),((\mathrm{SP})-2) \leftarrow-(\mathrm{PCL}),(\mathrm{SP}) \leftarrow(\mathrm{SP})+2, \\ & \text { If }(\text { sign })=0,(\mathrm{PC}) \leftarrow(\mathrm{ADDR})+3 \quad \text { Call subroutine if sign minus } \end{aligned}$ |
| CPE | ```If (parity)=1, ((SP)-1)\leftarrow-(PCH), ((SP)-2)\leftarrow(PCL), (SP) \leftarrow(SP)+2, (PC) <- ADDR If (parity)=0, (PC) \leftarrow(PC)+3 Call subroutine if parity even``` |
| CPO | ```If (parity)=0,((SP)-1) \leftarrow(PCH), ((SP)-2) \leftarrow(PCL), (SP) \leftarrow(SP)+2, (PC) \leftarrow ADDR If (parity)= l, (PC) < (PC)+3 Call subroutine if parity odd``` |

Condition bits affected: None

Format:
[LABEL:] CODE

| CODE | DESCRIPTION |
| :---: | :---: |
| RET | $(\mathrm{PCL}) \leftarrow((\mathrm{SP})),(\mathrm{PCH}) \leftarrow((\mathrm{SP})+1) ;(\mathrm{SP}) \leftarrow(\mathrm{SP})+2$ <br> Return from subroutine |
| RC | $\begin{aligned} & \text { If }(\operatorname{carry})=1,(P C H) \leftarrow((S P)),(P C H) \leftarrow((S P)+1), \quad(S P) \leftarrow(S P)+2 \\ & \text { If }(\operatorname{carry})=0,(P C) \leftarrow(P C)+3 \end{aligned}$ |
| RNC |  |
| RZ | $\begin{aligned} & \text { If }(\text { zero })=1,(P C L) \leftarrow((S P)),(P C H) \leftarrow((S P)+l),(S P) \leftarrow(S P)+2 \\ & \text { If }(\text { zero })=0,(P C) \leftarrow(P C)+3 \quad \text { Return if zero set } \end{aligned}$ |
| RNZ |  |
| RM | ```If (sign) = l, (PCL) \longleftarrow ((SP)), (PCH) \leftarrow((SP)+l), (SP) \leftarrow(SP)+2 If (sign) = 0, (PC) \leftarrow (PC) +3 Return if minus``` |
| RP |  |
| RPE |  |
| RPO | If (parity) $=0,(\mathrm{PCL}) \leftarrow((\mathrm{SP})),(\mathrm{PCH}) \leftarrow((\mathrm{SP})+1),(\mathrm{SP}) \leftarrow(\mathrm{SP})+2$ <br> If (parity) $=1,(\mathrm{PC}) \leftarrow(\mathrm{PC})+3 \quad$ Return if parity $\quad \underset{1}{ }$ |

Condition bits affected: None

RST INSTRUCTION

Format:
[LABEL:]
RST
EXP
Note: 0 EXP 7

| CODE | DESCRIPTION |  |
| :--- | :---: | :---: |
| RST | $((S P)-1) \leftarrow(P C H),((S P)-2)$$\leftarrow(P C L), \quad(S P) \leftarrow(S P)+2$ |  |
|  | $(\mathrm{PC}) \leftarrow-0000000000 E X P 000 B$ | Call subroutine at address <br> specified by EXP |

Condition bits affected: None

## INTERRUPT FLIP FLOP INSTRUCTIONS

Format:
[LABEL:]
CODE

| CODE |  | DESCRIPTION |
| :--- | :--- | :--- |
| EI | $($ INTE $) \longleftarrow 1$ | Enable the interrupt system |
| DI | $($ INTE $) \longleftarrow 0$ | Disable the interrupt system |

Condition bits affected: None

INPUT/OUTPUT INSTRUCTIONS

Format:

| [LABEL:] CODE | EXP |  |
| :--- | :--- | :--- |
| CODE |  | DESCRIPTION |
| IN | $(\mathrm{A}) \longleftarrow \quad$ input device | Read a byte from device EXP into <br> the accumulator |
| output device $\longleftarrow \quad$ (A) | Send the accumulator contents to <br> device EXP |  |

Condition bits affected: None

HLT INSTRUCTION

Format:
[LABEL:] HLT

| CODE | DESCRIPTION |
| :---: | :---: |
| HLT | -------------- |
| Instruction execution halts until <br> an interrupt occurs. |  |

Condition bits affected: None

ORG PSEUDO - INSTRUCTION

Format:
ORG EXP

| Code | De:cription |  |
| :--- | :---: | :--- |
| ORG | LOCATION COUNTER $\longleftarrow \quad$ EXP | Set Assembler location <br> counter to EXP |

EQU PSEUDO - INSTRUCTION
Format:

| NAME EQU |  |  |
| :--- | :--- | :--- |
| Code | Description |  |
| EQU | NAME $\longleftarrow$ EXP | Assign the value EXP <br> to the symbol NAME |

END PSEUDO - INSTRUCTION

Format:

END

| Code | Description |
| :--- | :---: |
| END | End the assembly. |

This appendix summarizes the bit patterns and number of time states associated with every 8080 CPU instruction.

When using this summary, note the following symbology:

1) DDD represents a destination register. SSS represents a source register. Both DDD and SSS are interpreted as follows:

| DDD or SSS | Interpretation |
| :---: | :---: |
| 000 | Register B |
| 001 | Register C |
| 010 | Register D |
| 011 | Register E |
| 100 | Register H |
| 101 | Register L |
| 110 | A memory register |
| 111 | The accumulator |

2) Instruction execution time equals number of time periods multiplied by the duration of a time period.

A time period may vary from 480 nanosecs to 2 microsec.
When two numbers of time periods are shown (eg. 5/ll), it means that the smaller number of time periods will be required if a condition is not met, and the larger number of time periods will be required if the condition is met.

| MNEMONIC | $\mathrm{D}_{7}$ | $\mathrm{D}_{6}$ | $\mathrm{D}_{5}$ | $\mathrm{D}_{4}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{1}$ | $\mathrm{D}_{0}$ | Number of Time Periods |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| CALL | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 17 |
| CC | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 0 | 11/17 |
| CNC | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 11/17 |
| CZ | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 11/17 |
| CNZ | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 11/17 |
| CP | 1 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 11/17 |
| CM | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 11/17 |
| CPE | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 0 | 11/17 |
| CPO | 1 | 1 | 1 | 0 | 0 | 1 | 0 | 0 | 11/17 |
| RET | 1 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 10 |
| RC | 1 | 1 | 0 | 1 | 1 | 0 | 0 | 0 | 5/11 |
| RNC | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 5/11 |
| RZ | 1 | 1 | 0 | 0 | 1 | 0 | 0 | 0 | 5/11 |
| RNZ | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 5/11 |
| RP | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 5/11 |
| RM | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 5/11 |
| RPE | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 5/11 |
| RPO | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 5/11 |
| RST | 1 | 1 | A | A | A | 1 | 1 | 1 | 11 |
| IN | 1 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 10 |
| OUT | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 10 |
| LXI B | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 10 |
| LXI D | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 10 |
| LXI H | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 10 |
| LXI SP | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 1 | 10 |
| PUSH B | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 11 |
| PUSH D | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 11 |
| PUSH H | 1 | 1 | 1 | 0 | 0 | 1 | 0 | 1 | 11 |
| PUSH A | 1 | 1 | 1 | 1 | 0 | 1 | 0 | 1 | 11 |
| POP B | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 10 |
| POP D | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 1 | 10 |
| POP H | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 10 |
| POP A | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 10 |
| STA | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 0 | 13 |
| LDA | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 0 | 13 |
| XCHG | 1 | 1 | 1 | 0 | 1 | 0 | 1 | 1 | 4 |
| XTHL | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 1 | 18 |
| SPHL | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 1 | 5 |
| PCHL | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 5 |
| DAD B | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | $\pm 0$ |
| DAD D | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 10 |
| DAD H | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 10 |
| DAD SP | 0 | 0 | 1 | 1 | 1 | 0 | 0 | 1 | 10 |
| STAX B | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 7 |
| STAX D | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 7 |
| LDAX B | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 7 |
| LDAS D | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 7 |
| INX B | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 5 |
| INX D | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 5 |
| INX H | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 1 | 5 |
| INX SP | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 5 |


| MNEMON IC | $\mathrm{D}_{7}$ | $\mathrm{D}_{6}$ | $\mathrm{D}_{5}$ | $\mathrm{D}_{4}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{1}$ | $\mathrm{D}_{0}$ | Number of Time Periods |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| MOV $r_{1}, r_{2}$, | 0 | 1 | D | D | D | S | S | S | 5 |
| MOV M,r | 0 | 1 | 1 | 1 | 0 | S | S | S | 7 |
| MOV $\mathrm{r}, \mathrm{M}$ | 0 | 1 | D | D | D | 1 | 1 | 0 | 7 |
| HLT | 0 | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 7 |
| MVI r | 0 | 0 | D | D | D | 1 | 1 | 0 | 7 |
| MVI M | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 10 |
| INR | 0 | 0 | D | D | D | 1 | 0 | 0 | 5 |
| DCR | 0 | 0 | D | D | D | 1 | 0 | 1 | 5 |
| INR A | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 5 |
| DCR A | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 1 | 5 |
| INR M | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 10 |
| DCR M | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 10 |
| ADD r | 1 | 0 | 0 | 0 | 0 | S | S | S | 4 |
| ADC r | 1 | 0 | 0 | 0 | 1 | S | S | S | 4 |
| SUB r | 1 | 0 | 0 | 1 | 0 | S | S | S | 4 |
| SBB r | 1 | 0 | 0 | 1 | 1 | S | S | S | 4 |
| NDA $r$ | 1 | 0 | 1 | 0 | 0 | S | S | S | 4 |
| XRA $r$ | 1 | 0 | 1 | 0 | 1 | S | S | S | 4 |
| ORA $r$ | 1 | 0 | 1 | 1 | 0 | S | S | S | 4 |
| CMP r | 1 | 0 | 1 | 1 | 1 | S | S | S | 4 |
| ADD M | 1 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 7 |
| ADC M | 1 | 0 | 0 | 0 | 1 | 1 | 1 | 0 | 7 |
| SUB M | 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 7 |
| SBB M | 1 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 7 |
| NDA M | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 7 |
| XRA M | 1 | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 7 |
| ORA M | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 7 |
| CMP M | 1 | 0 | 1 | 1 | 1 | 1 | 1 | 0 | 7 |
| ADI | 1 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 7 |
| ACI | 1 | 1 | 0 | 0 | 1 | 1 | 1 | 0 | 7 |
| SUI | 1 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 7 |
| SBI | 1 | 1 | 0 | 1 | 1 | 1 | 1 | 0 | 7 |
| NDI | 1 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 7 |
| XRI | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 7 |
| ORI | 1 | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 7 |
| CPI | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 7 |
| RLC | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 4 |
| RRC | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 4 |
| RAL | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 1 | 4 |
| RAR | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 4 |
| JMP | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 1 | 10 |
| JC | 1 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 10 |
| JNC | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 10 |
| JZ | 1 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 10 |
| JNZ | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 10 |
| JP | 1 | 1 | 1 | 1 | 0 | 0 | 1 | 0 | 10 |
| JM | 1 | 1 | 1 | 1 | 1 | 0 | 1 | 0 | 10 |
| JPE | 1 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 10 |
| JPO | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 10 |


| MNEMONIC | $\mathrm{D}_{7}$ | $\mathrm{D}_{6}$ | $\mathrm{D}_{5}$ | $\mathrm{D}_{4}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{1}$ | $\mathrm{D}_{0}$ | Number of Time Periods |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DCX B | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 5 |
| DXC D | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 5 |
| DCX H | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 5 |
| DCX SP | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 1 | 5 |
| CMA | 0 | 0 | 1 | 0 | 1 | 1 | 1 | 1 | 4 |
| STC | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 4 |
| CMC | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 4 |
| DAA | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 1 | 4 |
| SHLD | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 17 |
| LHLD | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 17 |
| EI | 1 | 1 | 1 | 1 | 1 | 0 | 1 | 1 | 4 |
| DI | 1 | 1 | 1 | 1 | 0 | 0 | 1 | 1 | 4 |
| NOP | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 4 |

## APPENDIX C

HEXADECIMAL PROGRAM TAPE FORMAT

The hexadecimal tape format used by the Intecolor ${ }^{\circledR} 8001$ system is a modified memory image, blocked into discrete records. Each record contains record length, record type, memory address, and checksum information in addition to data. A frame by frame description is as follows:

Frame 0

Frames l,2
(0-9,A-F)

Frames 3 to 6

Frames 7,8

Frames 9 to 9+2* (Record
Length) - l

Record Mark, Signals the start of a record. The ASCII character colon (":" HEX 3A) is used as the record mark.

Record Length. Two ASCII characters representing a hexadecimal number in the range 0 to 'FF'H (O to 255). This is the count of actual data bytes in the record type or checksum. A record length of 0 indicates end of file.

Load Address. Four ASCII characters that represent the initial memory location where the data following will be loaded. The first data byte is stored in the location pointed to by the load address, succeeding data bytes are loaded into ascending addresses.

Record Type. Two ASCII characters. Currently all records are type 0, this field is reserved for future expansion.

Data. Each 8 bit memory word is represented by two frames containing the ASCII characters ( 0 to 9, A to F) to represent a hexadecimal value 0 to 'FF'H (0 to 255).

| Frames 9+2* (Record Length) to <br> 9+2* (Record Length) +1 | Checksum. The checksum is the negative of the sum of all 8 bit bytes in the record since the record mark (":") evaluated modulus 256. That is, if you add together all the 8 bit bytes, ignoring all carries out of an 8-bit sum, then add the checksum, the result is zero. |
| :---: | :---: |

Example: If memory locations 1 through 3 contain 53F8EC, the format of the hex file produced when these locations are punched is:
: 0300010053F8ECC5

Note: This format is also known as the Intel format.

16 conemantions between blue \& RED ( 5 EACH) MulTIPLY By 4 Different green states (ir inch)


64 COMPUCOLR COLORS? BLUE LEFT, RED RIGHT, GREEN TOP


|  | +1 | $\div 4$ |
| :--- | :--- | :--- |
| 0 | 1 | .25 |
| 1 | 2 | .5 |
| 2 | 3 | .75 |
| 3 | 4 | 1 |
| 4 | 5 | 1.25 |
| 5 | 6 | 1.5 |
| 6 | 7 | 1.75 |
| 7 | 8 | 2 |
| 8 | 9 | 2.25 |
| 9 | $A$ | 2.5 |
| $A$ | $B$ | 2.75 |
| $B$ | $C$ | 3. |
| $C$ | $D$ | 3.25 |
| $D$ | $E$ | 3.5 |
| $E$ | $F$ | 3.75 |
| $F$ | 16 | 4 |

 Not AND DIVIDEBY 2 01100101


[^0]:    Signal Ground

[^1]:    NOTE: THE TERMINAL ACCEPTS ALL 80 TO 日f HEX CODES FROM THE KEYBOARD AMD REASSIGN8 THEM
    FQ TO FF HEX WHEN IN THE PLOT MODE, UNLESS THE OPTIONAL KEYS ARE INSTALLED. THEREFORE TO TO HEX WHEN IN THE PLOT MODE, UNLESS THE OPTIONAL KEYS ARE INS
    WITHOUT THE FUNCTION KEYS THE KEYBOARD CAN PLOT IN A RANGE OF TO IT5.
    $\qquad$

[^2]:    - Stresses beyond those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions beyond those indicated in the "Recommended Operating Conditions" section of this specification is not implied. Exposure to absolute-maximum-rated conditions for extended periods may affect device reliability.
    NOTE 1: Under absolute maximum ratings voltage values are with respect to the normally most negative supplyvoltage, $V_{B B}$ (substrate). Throughout the remainder of this data sheet, voltage values are with respect to $V_{S S}$ unless otherwise noted.

[^3]:    - Two possible cycle times (11/17) indicate instruction cycles dependent on condition flags.
    ${ }^{\dagger}$ All flags (C, Z, S, P, C1) affected.
    Only carry flag affected.

[^4]:    $\ddagger$ Only carry flag affected.
    \$All flags except carry affected.

[^5]:    - Two possible cycles times (11/17) indicate instruction cycles dependent on condition flags. ${ }^{\dagger}$ All flags (C, Z, S, P, C1) affected.
    + Only carry flag affected.

[^6]:    *Two possible cycles times ( $11 / 17$ ) indicate instruction cycles dependent on condition flags.
    $\dagger$ All flags (C, Z, S, P, C 1) affected.
    $\because$ Only carry flag affected.

[^7]:    Condition bits affected: None

